



VT6508

8 RMII PORTS OF 10/100BASE-T/TX ETHERNET SWITCH CONTROLLER

REVISION D DATASHEET (Preliminary)

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VIA Technologies, Inc.



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FIGURES AND TABLES

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REVERSION HISTORY

Reversion	Date	Reason for change	Ву
V0.01	11/23/1999	First release version	Murphy Chen



FEATURES

- Single chip 8 ports 10/100Mbps Ethernet switch controller
 - Highly integrated single chip shared memory switch engine
 - Supports 8 RMII (Reduced Media Independent Interface) ports
 - Non-blocking layer 2 switch, 148,810 packets/sec on each 100Mbps Ethernet port
- Media Access Control (MAC)
 - Dual 64-byte FIFO's per port for receiving and transmitting
- Auto-sensing 10/100Mbps media speed, full/half -duplex mode, and flow-control capability
- Two switching mechanisms
 - Supports 'store and forward' switching with filtering CRC-bad packets
 - Supports 'cut through' switching subject to long packets of length over 64 bytes
- Packet buffering
 - Glueless 32-bit interface to SSRAM as a packet buffer pool
 - 64 packet buffers for 32Kx32 SSRAM, 149 packet buffers for 64Kx32 SSRAM
 - 1536 bytes for each packet buffer
- Storing control data (forwarding table and packet link entries) in external SSRAM
 - Buffer status is stored in the internal free-buffer bit map registers
 - Shared forwarding table of 2K entries (or 4K slots) to support multiple (up to 4K) Mac addresses per port
- Efficient address recognition, self learning, and auto-aging mechanism
 - Two-slot hashing algorithm to prevent hash collision
 - Two optional hashing algorithms, CRC-map and direct-map
- Advanced congestion control mechanism
 - IEEE 802.3X compliant flow control for full duplex ports
 - Backpressure for half duplex ports
 - Drop control for full duplex ports without flow control capability
 - Incorporating with the output private buffering scheme to prevent HOL (head of line) blocking
- Supports port mirroring (Sniffer feature)
- By-pass VLAN packets
- Support port-based trunking
 - Support two individual trunk groups, each of 2 member ports
 - Load balance according to DMAC address and source port number
- Support chip initialization through EEPROM or by strapping only
 - Supports I²C EEPROM interface for customized configuration
 - Supports LED serial-out in the strapping-only initialization mode
- 50MHz internal reference clock rate
- 83~100MHz SSRAM clock rate, typically 83MHz
- Single +3.3V supply, 0.3µm TSMC CMOS technology
- 128-pin PQFP package





BLOCK DIAGRAM

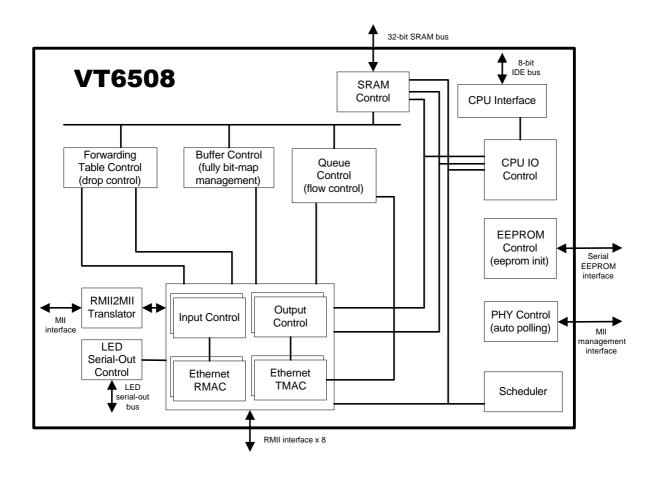


Figure 1: Function Block Diagram of VT6508. (Note that some interface signals are only available in VT6509 of 208-pin PQFP package.)

99/12/09



PINOUT DIAGRAM

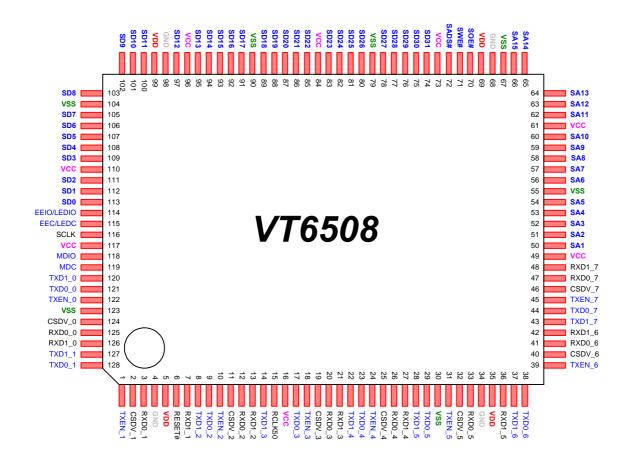


Figure 2: Pinout Diagram of VT6508.

99/12/09



PIN DESCRIPTIONS

No.	Name	Type	Description			
	RMII interface					
See Fig. 2	CSDV[7:0]	I	Carries sense and data valid from port 7 to port 0:			
See Fig. 2	RXD0[7:0]	I	Receive data zero from port 7 to port 0:			
See Fig. 2	RXD1[7:0]	I	Receive data one from port 7 to port 0:			
See Fig. 2	TXEN[7:0]	O	Transmit enable from port 7 to port 0:			
See Fig. 2	TXD0[7:0]	O	Transmit data zero from port 7 to port 0:			
See Fig. 2	TXD1[7:0]	O	Transmit data one from port 7 to port 0:			
			SRAM Interface			
See Fig. 2	SA[15:1]	О	SRAM Address Bus:			
			15-bit SDRAM data bus. These signals connect directly to the address input of the SDRAM devices.			
See Fig. 2	SOE	О	Output Enable			
See Fig. 2	SWE	О	SRAM Write Enable			
See Fig. 2	SADS	О	Synchronous Processor Address Status			
See Fig. 2	SD[31:0]	I/O	SRAM Data:			
			32-bit SRAM data bus. These signals connect directly to the data input/output pins of the SRAM devices.			
		•	Miscellaneous Interface			
See Fig. 2	EEC/LEDC	0	Serial EEPROM Interface Clock (if strapping SD25 is default high)			
			The on-board EEPROM device address must be 1010 001 XXXXXXXX.			
			Serial LED Serial-Out Clock (if strapping SD25 is low)			
			The LEDC is 1MHz and has a burst per 50ms.			
See Fig. 2	EEIO/LEDIO	I/O	Serial EEPROM Interface Data (if strapping SD25 is default high)			
			Serial LED Serial-Out Data Output (if strapping SD25 is low)			
See Fig. 2	MDC	О	Management Interface (MI) Clock Output			
See Fig. 2	MDIO	I/O	Management Interface (MI) Data I/O			
See Fig. 2	RCLK50	I	50MHz Main Reference Clock			
See Fig. 2	SCLK	I	SRAM Reference Clock			
			The suggested clock rate is 83MHz or more high for non-blocking requirement.			
See Fig. 2	RESET	I	SYSTEM RESET			

Power Sup	ply & Ground		
See Fig. 2	VDD	P	Positive 3.3V supply to the core digital logic.
See Fig. 2	VCC	P	Positive 3.3V supply to all I/O pads.



See Fig. 2	GND	G	Ground supply to the core digital logic.
See Fig. 2	VSS	G	Ground supply to all I/O pads.

DEFINITION OF VT6508B STRAPPING PINS

SD[31:0]: *Note that the default strapping bit value is 1

SD[31:0]:	*Note that the default strapping bit value is $ 1 $.
Pin	Description
SD[0]	Broadcast or Drop BPDU Packets: BCAST_BPDU == 1' b1 => broadcast BPDU packet if CPU_FWD_CFG[1] == 0 (default) BCAST_BPDU == 1' b0 => drop BPDU packet if CPU_FWD_CFG[1] == 0
	Note that BPDU packets will be forwarded to CPU port if CPU_FWD_CFG[1] == 1 without regard to the value of BCAST_BPDU.
SD[1]	LED Display Combination of Link Activity Status with RX/TX Event: LED_COMB [0] == 1' b1 => combined (default) LED_COMB [0] == 1' b0 => not combined
	Note that the LED bit sequence [10:1] is used originally to display link activity status. By setting combination with RX/TX event, it will indicate both the link activity status and RX/TX events so that required LED number can be reduced.
SD[2]	Enable Drop Control for Private Buffer Reservation: DROP_CONTROL_EN [0] == 1' b1 => enabled (default) DROP_CONTROL_EN [1] == 1' b0 => disabled
	Note that drop control has a lower priority than flow control and backpressure. If it is disabled, all TMACs will not make the drop window signals to Forwarding Control. If it is enabled and the flow control/backpressure mechanism is not enabled, all TMACs will make the drop window signals to Forwarding Control. If it is enabled and the flow control mechanism is also enabled, but the full-duplex party has no flow control capability, TMAC will make the drop window signal to Forwarding Control.
SD[3]	SRAM TYPE: SRAM_TYPE == 1' b1 => 64Kx32 (default) SRAM_TYPE == 1' b0 => 32Kx32
	size=128KB: FREEMCNT=64, XON_THRED=43, PRIVATE_BUF_SIZE=1 V_FREE=10, MIN_XOFF_CONST_THRED=15 size=256KB: FREEMCNT=149, XON_THRED=98, PRIVATE_BUF_SIZE=4 V_FREE=10, MIN_XOFF_CONST_THRED=28
SD[5:4]	CHIP_INIT[1:0] == 2' b11 => Chip Initialization via strapping only (default) CHIP_INIT[1:0] == 2' b10 => Chip Initialization via EEPROM CHIP_INIT[1:0] == 2' b01 => Chip Initialization via CPU CHIP_INIT[1:0] == 2' b00 => Chip Initialization via EEPROM in speedup mode In none speedup mode, EEC = 78.125K Hz, LED period = 50ms. In speedup mode for testing, EEC = 2.778MHz, LED period = 1ms without one-second flash.



SD[7:6]	Number of Tries of Excessive Collisions Before Dropping:
	RETRY_EXCE_COLL [1:0] == 2' b11 => 3 tries (default)
	RETRY_EXCE_COLL [1:0] == 2' b10 => 2 tries
	RETRY_EXCE_COLL [1:0] == 2' b01 => retry forever
	RETRY_EXCE_COLL [1:0] == 2' b00 => 1 try
	Internally, retry_no = RETRY_EXCE_COLL, or infinite if it is zero.
	By default, after 3 tries (i.e. 3 turns of (excessive) 16 collisions), the outgoing packet will
CD[45:0]	be dropped.
SD[15:8]	(reserved with default value pull-high)
SD[17,16]	Trunking Mode:
	SD16 == 1, port (0,1) no trunking (default)
	0, port (0,1) in trunking mode,
	SD17 == 1, port (6,7) no trunking (default)
	0, port (6,7) in trunking mode.
	To simplify layout, trunk group (6,7) can connect with neighboring (right-side)
	VT6508's trunk group (0,1) in on-board manner. Note that on-board trunk ports do not need auto-polling and they must be full-duplex &
	100Mbps ports & forced flow control enable mode.
SD[18]	Enable Backpressure:
OD[10]	BACKPRESSURE_EN [0] == 1' b1 => enabled (default)
	BACKPRESSURE_EN [1] == 1' b0 => disabled
	Briori Resource_ent[1] == 1 55 => disabled
	Note that if it is disabled, all TMAC of half-duplex ports will ignore the flow control
	XON/XOFF signals from Queue Control.
SD[20:19]	
' '	RMII_PORT_FC_EN[1:0] == 2' b11 => for RMII ports, set self flow control ability
	register bit and detect party's flow control ability (default)
	RMII_PORT_FC_EN[1:0] == 2' b10 => for RMII ports, disable self flow control ability
	register bit and assume (ignore auto-polling) party has no flow control capability
	RMII_PORT_FC_EN[1:0] == 2' b01 => for RMII ports, enable self flow control ability
	register bit and assume (ignore auto-polling) party has flow control capability
	RMII_PORT_FC_EN[1:0] == 2' b00 => same as 2' b11.
	While enabled (RMII_PORT_FC_EN[1:0] == 2' b11 or 2' b00), set self each PHY
	device's flow control ability register bit PHY_ANAR_4.10 as 1, and auto polling
	PHY_ANLPAR_5.10 to check whether the party has the flow control capability.
	While forced disabling (RMII_PORT_FC_EN[1:0] == 2' b10), set self PHY devices'
	flow control ability register bit PHY_ANAR_4.10 as 0 (no such capability), and does not
	need auto polling PHY_ANLPAR_5.10.
	While forced disabling (RMII_PORT_FC_EN[1:0] == 2' b01), set self PHY devices' flow control ability register bit PHY_ANAR_4.10 as 1 (with flow control capability), and
	does not need auto polling PHY ANLPAR 5.10.
	For forced disabling, turn off flow control enable bits of all RMII ports.
	For (forced) enabling, turn on flow control enable bits of all RMII ports.
SD[23:21]	
	LED_MASK[0] == 1' b1 => enable group 0 data out (default)
	LED_MASK[1] == 1' b1 => enable group 1 data out (default)
	LED_MASK[2] == 1' b1 => enable group 2 data out (default)
	LED Groups 3,4,5 serial out is always enabled. The time period of LED sequence is
	50ms.



SD[24]	Hash Algorithm Selection:
	HASH_ALG_SEL == 1 => CRC-map with scramble (default)
	HASH_ALG_SEL == 0 => Direct-map without scramble
	HASH_ALG_SEL == 1 will make the register HASH_ALG[1:0]=0 that is helpful to the
	X-Stream test.
	HASH_ALG_SEL == 0 will make the register HASH_ALG[1:0]=3 that can get a good result of 4094 MAC addresses in Address Handling test.
SD[25]	No Swap EEPROM and LED output pins:
	NO_SWAP_EEPROM_LED == 1 => The pins 192/193 are used as
	EEIO/EEC. (default)
	NO_SWAP_EEPROM_LED == 0 => The pins 192/193 are used as
	LEDIO/LEDC.
	In 128-pin package, we can bond this pin to VSS to select LED output from the original EEPROM output pins.
SD[26]	Enable Aging:
	EN_AGING == 1 => enable aging function (default)
	EN_AGING == 0 => disable aging function
CD[27]	Disable Cut Through Enables
SD[27]	Disable Cut Through Enable: DIS_CUT_THRU == 1 => disable cut-through feature
	(default, i.e. store & forward)
	DIS_CUT_THRU == 0 => enable cut-through feature
SD[28]	Disable Aggressive Backoff Algorithm:
05[20]	DIS AGGRES BACKOFF == 1' b1 => disabled (default)
	DIS_AGGRES_BACKOFF == 1' b0 => enabled
	DIG_NOCINED_B/NONCIN == 1 50 => chabled
	Note that if this feature is enabled, the backoff algorithm for all Ethernet ports is in MBA
	mode. By default, the backoff algorithm bits MBA=0 and OFSET=1.
SD[29]	Backpreasure Collision Point Configuration:
	COLL_POINT_CFG == 1 => 32th byte (default)
	COLL_POINT_CFG == 0 => 2nd byte
SD[30]	Disable TEST mode:
	DIS_TEST_MODE == 1 => disabled (default)
	DIS_TEST_MODE == 1 => enabled
SD[31]	Disable Latchup Mode:
	DIS_LATCH_UP_MODE == 1' b1 => disable Latchup mode (default)
	DIS_LATCH_UP_MODE == 1' b0 => enable Latchup mode, all output only IO PADs
	are in tri-state after reset.



SECTION I FUNCTIONAL DESCRIPTIONS

1 GENERAL DESCRIPTION

The VT6508 is a low-cost switch engine chip implementation of an 8 ports 10/100Mbps Ethernet switch system for IEEE 802.3 and IEEE 802.3u networks. Each port can be either auto-sensing or manually selected via EEPROM configuration to run at 10Mbps or 100Mbps speed rate, full or half duplex mode.

The VT6508 supports RMII (reduce MII) port interface. There are eight independent MACs within the VT6508 chip. The MAC controller controls the receiving, transmitting, and deferring of each individual port, and the MAC controller also provides framing, FCS checking, error handling, status indication and flow control function. It has wire-speed performance with forwarding rate of 148,810 packets/sec on each 100Mbps Ethernet port. The VT6508 can be configured via EEPROM or strapping only.

2 THE VIA ETHER SWITCH ARCHITECTURE

The VT6508 switch engine uses the shared memory architecture. In order to improve the packet latency, VT6508 provides two methods for packet switching, cut-through and store-and-forwarding.

2.1 Switch initialization procedures

- Step_1. Read strapping signal.
- Step_2 Write Reg 0619H with WR_DATA = 8'h01 to trigger SRAM_Ctrl module autotest.
- Step_3. If initialization by Strapping, then go to Step_4.

 If initialization by EEPROM, eeprom_ctrl module will download register data from EEPROM (by internal register bus IOW, CS_*, IO_ADDR, and IO_WR_DATA).

 If EEPROM data download error, then go to step6.
- Step_4 After SRAM auto-test completed, write Reg 0304H with WR_DATA = 8'h01 to trigger Forward Ctrl module auto aging.
- Step_5 Write Reg 0404H with WR_DATA = 8'h08 to trigger PHY_Ctrl module phy_auto-polling.
- Step_6 Write Reg 000aH with WR_DATA = 8'h01 to trigger LED_Ctrl module to display LED.
- Step_7 Chip initialization done.



2.2 Packet Switching Flow

- 1. After initialization, the input control of each enabled port will pre-allocate one packet buffer from the free buffer pool.
- 2. When the Receive MAC (RMAC) receives a packet data from the PHY device via the RMII network interface, it packs the data into 16-bit words then passes it to input control. If RMAC detects any error, it also notifies input control to stop forwarding process.
- 3. Input control extracts the destination MAC address from incoming data, passes it to forwarding table control for forwarding decision. In the mean while, it packs 16-bit words into 64-bit quad-words, and saves it to an input FIFO before storing the packet data to packet buffers in SRAM.
- 4. If the switch is configured to "store and forward" mode, input control queues the packet to the output queue of the destination port after input control is informed by RMAC that this is a good packet and it stores all packet data to SRAM. If the switch is configured to "cut-through" mode, the input control queues the packet to the output queue of the destination port after the first 64-byte packet data is stored in SRAM to prevent output FIFO underrun.
- 5. After the whole packet is received and FCS is correct, input control pass the source MAC address of the packet to forwarding table control for address learning.
- 6. Output control of the outbound port dequeues the packet from itself output queue, and fetches packet data from SRAM and saves it into output FIFO. Then it notifies the Transmit MAC (TMAC) to transmit the read-to-outgoing packet.
- 7. TMAC grabs 16-bit word at a time from output control, adds preamble and SFD to the beginning of the packet, then send them out. Proper deferring is done for collision to conform to 802.3 standard.
- 8. After the packet is successfully transmitted, TMAC notifies output control of the successful transmission. Output control then returns the packet to the free buffer pool.

2.3 PACKET BUFFERS AND FORWARDING TABLE

VT6508 provides a 32-bit SRAM interface for packet buffering and maintaining address table and per-port output link lists. Each packet buffer is a 1536-byte contiguous memory block in SRAM, and it also contains to a 8-byte link node data structure in every buffer's tail. A link node corresponds uniquely to a packet buffer that could belong an output queue, the free buffer pool, incoming packet buffers of an input control, or outgoing packet buffers of an output control.

Initially, each input port control will request one packet buffer from its private buffer pool. Each time when a packet buffer is consumed by an incoming packet, the input port control will request another packet buffer to prepare for next packet.

There are two data structures in SRAM memory layout, shown in Figure 3. One is the forwarding table that locates at the lowest 32K-byte address space and contains 2048 entries (buckets), each of 2 slots. The structure of an 64-bit forwarding table slot is shown in Figure 4, composed of:



(1) bits 39..0: MAC Tag

The MAC Tag is used to record the partial content of MAC address for lookup identification. The mapping of Tag and MAC address is as follows:

Tag[39:32] = MAC[15:8]

Tag[31:24] = MAC[23:16]

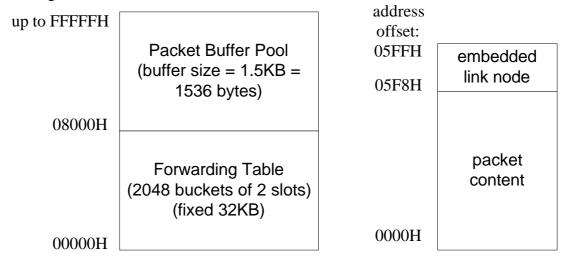


Figure 3. SRAM memory layout.

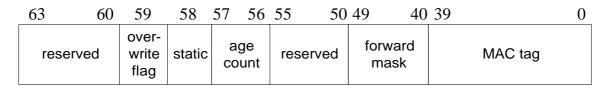


Figure 4. Data structure of forwarding table slot.

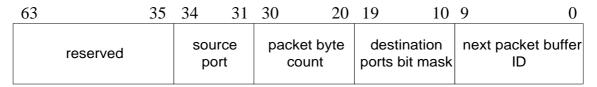


Figure 5. Data structure of embedded link node.

Tag[23:16] = MAC[31:24]

Tag[15:8] = MAC[39:32]

Tag[7:0] = MAC[47:40]

(2) bits 49..40: Forwarding Port Mask



- (3) This is the bit mask of the destination ports corresponding to the forwarding table slot. The bits 49 and 48, corresponding to CPU port and MII port respectively, should be zero in VT6508.
- (4) bits 57..56: Age Count
- (5) If this field is equal to zero, it means this slot is an empty one; otherwise, it is valid. When the auto-aging is enabled (by default), every valid slot' age count will be decreased by one per 100 seconds.
- (6) bits 58: Static Flag
- (7) This flag is used to program a static MAC slot, that will not be overwritten by dynamic learning mechanism, with specified forwarding port mask for special purpose.
- (8) bits 59: Over-Write Flag

During learning, if both slots are used, the 3rd address will replace one of the slots according to current value of the concatenated over-write order pointer and their static bits regardless of ages of the two slots.

Let the concatenated over-write pointer ow[1:0] = (over-write pointer of slot 1, over-write pointer of slot 0).

Let $static[1:0] = \{static bit of slot 1, static bit of slot 0\}.$

The possible mapping of ow[1:0] and static[1:0] are as follows:

	static[1:0]			
ow[1:0]	(0,0)	(0,1)	(1,1)	(1,0)
(0,0)	0	1	X	0
(0,1)	1	1	X	0
(1,1)	0	1	X	0
(1,0)	1	1	X	0

The internal operation rules for the over-write pointer are as follows:

- (1) When a slot is over-written, its over-write pointer is toggled.
- (2) When one of the slots is a static entry, the static entry must always be placed in slot 0, and slot 1 is always the entry over-written. Over-write pointer has no effect what so ever here.
- (3) When both slots are static, no entry can be written to either slot by the learning process.
- (4) When update an entry, value of the over-write pointer must not be changed unless indicated otherwise, i.e. over-writing an existing slot.

Another data structures in SRAM memory layout is the packet buffer pool, located at the upper address space. Every packet buffer is 1.5K bytes with a corresponding embedded link node in the tailing 8 bytes. The embedded link nodes are used to construct output queues, one queue for each output port. Its structure is shown in Figure 5, composed of:

- (1) bits 9..0: Next Packet Buffer ID
- (2) bits 19..10: Destination Ports Bit Mask



(3) bits 30..20: Packet Byte Count

(4) bits 34..31: Source Port ID

The number of packet buffers is determined by the configured SRAM size, listed as follows:

SRAM Size	128KB	256KB	512KB	1MB
maximum # of packet buffers	64 buffers	149 buffers	320 buffers	661 buffers

2.3 RMII INTERFACE

The VT6508 can directly connect to an 8 port RMII PHY device through the reduced MII (RMII) interface to construct a small-sized system board. The signals of RMII interface are described as follows:

Name	Type	Description
CRSDV	I	Carrier sense and Data
		valid
RXD[0-1]	I	Receive data bit 0 to 1,
		data rate with 50MHz
TXEN	О	Transmit Enable
TXD[0-1]	О	Transmit Data bit 0 to
		1

2.3 MANAGEMENT INTERFACE AND AUTO NEGOTIATION

The VT6508 communicates with the external 10/100M PHY and access the PHY register through MDC, MDIO. The signals of management interface are described as follows:

Name	Type	Description
MDC	О	management interface clock
MDIO	IO	management interface data signal

Auto-polling state machine:

step1. Auto-polling idle

- if auto-polling_command="on" then go to step2



- step2. Check the chip internal register, phyint_setup_reg, to setup the pause capability (Reg4.10) and phy auto-negotiation enable (Reg0.12)
 - if (this port do not need phy setup) and (phyint_setup_reg == done), then go to step 10(done)
 - else if (this port do not need phy setup) and (phyint_setup_reg != done), then config ioctl, set phyint_setup_reg == done, go to step10(done)
 - else if (phyint_setup_reg != done),
 then go to step3(phy fc setup)
 - else if (this port do not need auto-poll, then go to step 10(done)
 - else go to step 5(read phy link status)

step3. Phy flow_control ability setup

- if the phy_pause_ability(reg4.10) != port_phy_fc_set, then write reg4.10 = port_phy_fc_set and go to step4 (phy auto-negotiation ability setup)
- else go to step4 (phy auto-negotiation ability setup)

step4. PHY auto-negotiation ability setup

- if (Reg0.12)="1"(on), then modify chip internal register for phyint_setup_reg = done; go to step5
- else if (Reg0.12)="0"(off) or phy_pause_ability have been modified (in step3), then write PHY auto-negotiation_enable bit (Reg0.12) ="1" (on), restart_auto-negotiation bit (Reg0.9) = "1" and modify chip internal register for phyint_setup_reg = done; go to step5

step5. Read PHY link status bit (Reg1.2) 1st time

- if link is down (Reg1.2)="0", then go to step6
- else if link is up (Reg1.2)="1" and chip_internal_reg for link status is "down", then go to step7
- else if link is up (Reg1.2)="1" and chip_internal_reg for link status is "up", then go to step10.

step6. Read PHY link status bit (Reg1.2) 2st time

- if link is up (Reg1.2)="1", then go to step7
- else if link is down (Reg1.2)="0", then go to step10
 *if chip_internal_reg for link status is "down"=> generate link_change interrupt and disable io_port

step7. Read PHY Auto-negotiation Complete bit (Reg1.5)

- if auto-negotiation is completed (Reg1.5)="1", then go to step8.
- else if auto-negotiation is not completed (Reg1.5)="0", then go to step10.



*if chip_internal_reg for link status is "down"=> generate link_change interrupt and disable io_port

Step8. Read PHY ANAR(Reg4) and ANLPAR(Reg5)

- if speed_arbit_bit ="0"(speed from PHY proprietary reg), then go to step9
- if speed_arbit_bit ="1"(speed from ANLPAR), then go to step10
 - * generate link_change interrupt and enable io_port

Step9. Read PHY proprietary Reg for RMII speed, then go to step10

* generate link_change interrupt and enable io_port

Step10. Done

- if auto-polling_command = "off", then go to step1, and stay idle
- else go to step2, for next port.

Duplex mode decision rule:

```
RMII_DUPLEX_MODE = ~( ( ANAR.8 = 1 and ANLPAR.8 =1:100BASE-TX Full ) or ( ( ANAR.7 =0 or ANLPAR.7 =0) and ANAR.6 =1 and ANLPAR.6=1))
```

Note that value 1 denotes half duplex, 0 denotes full duplex.

Speed decision rule:

```
According to ANLPAR,
- if ~((ANLPAR.9=1 or ANLPAR.8=1 or ANLPAR.7=1)
then speed = 10MHz
```

2.4 FLOW CONTROL

In VT6508, there are three policies of resource management for different objectives:

- (1) flow control for full-duplex ports
 - Objective: slow down the input traffic without incurring any drops.
- (2) backpressure for half-duplex ports
 - Objective: slow down the input traffic with intended collision and the Ethernet backoff mechanism.
- (3) drop control for private buffer reservation
- Objective: shape the input traffic to be compliant to the buffer reservation discipline.

In the following, we call "congestion control" as aggregation of the above three policies.

The application conditions of resource management policies are as follows:

(1) flow control for full-duplex ports



Enabled only when the flow control feature is enabled and the full-duplex port's party has flow control capability.

Note that as both the flow control feature and the drop control feature are enabled, the drop control feature is not turned on if the full-duplex port's party has flow control capability.

(2) backpressure for half-duplex ports

Enabled only when the flow control feature is enabled and the port is in half-duplex mode. Note that as both the flow control feature and the drop control feature are enabled, the drop control feature is not turned on if the port is in half-duplex mode.

(3) drop control for private buffer reservation

Enabled when the flow control feature and the drop control feature are enabled, but the full-duplex port's party has not flow control capability.

Also enabled when the flow control feature is disabled and the drop control feature is enabled.

XON/XOFF Window and Drop Window

The flow control and backpressure are based on XON/XOFF window. Within XON window, the source port is in non-congestion status so that any incoming packets are forwarded normally. If an incoming packet violates the constraint of congestion control, the source port will leave the XON window and enter the XOFF window after this enqueue operation is served. Within XOFF window, the source port is in congestion status so that any incoming packet will trigger a congestion-control operation that depends on the used policy. The drop control mechanism, residing in Forwarding Control and enabled by DROP_EN = 1, is based on every ports' congestion windows that are defined as (1) any one is in XOFF window, (2) a destination port is in congestion status if its reservation buffer count == 0.

The XON/XOFF windows and the DROP ON/OFF windows are maintained by TMAC, according to port status and input signals from Queue Control. Figure 6 shows an example XOFF window status (named as XOFF_WINDOW[0]) of port 0, where XOFF[0] and XON[0] is an event generated by queue control to notify this source port should enter XOFF state, and XON[0] are events generated by queue control to notify this source port should enter or leave XOFF state, respectively. If the source port is within XOFF window as an XOFF event arrives, an flow control frame with pause time = 0xffff will be sent out. If the source port is within XON window, the arriving XON event will be ignored.



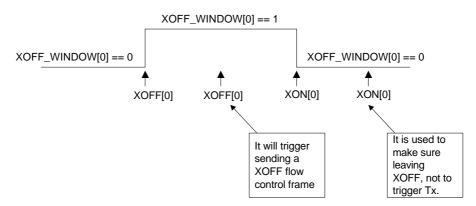


Figure 6. XON/XOFF Window Concept.

Congestion-Control Operations

The starvation control threshold is 9 in VT6508. Let Φ be the Free Memory Count. Let Ω be the Virtual Free Buffer Count, where $\Omega = \Phi$ - 9. In congestion control, decisions are based on Ω . Within XOFF window, the corresponding congestion-control operations are as follows:

(1) flow control operation

As a unicast or broadcast packet, enqueued by Input Control, violates the congestion control constraints, Queue Control will assert a trigger to its source port's TMAC to send out a flow control frame with pause time = FFFFH. There are two congestion conditions: (1) $\Phi \le \max\{\Psi,28\}$ and $R_k = 0$, (2) any port is in XOFF window and $R_k = 0$.

Note that even a source port is within XOFF window, an XOFF triggering signal from Queue Control will also make the TMAC to send out a flow control frame of pause time = FFFFH. Although it costs more bandwidth to perform the flow control operation, this duplicated transmission can guarantee the XOFF flow control frame could be received by the party successfully.

(2) backpressure for half-duplex ports

As a non-local incoming packet is received, Input Control will assert a signal of "Non_Local" to notify its TMAC sub-module about the status. TMAC will collide this packet if it is a half-duplex port and it is within XOFF window.

(3) drop control for private buffer reservation

As an incoming packet, coming from a port within DROP_ON window and destining to a port within CONGESTION window, is detected by Forwarding Control in lookup operation, Forwarding Control will unmask its port mask to make it to be dropped by Input Control. Note that the DROP_ON window indicates the time duration in which the drop feature to the specific port is enabled, that is determined by TMAC. Note that the CONGESTION window of port k is turned on when any port is in XOFF window and $R_k = 0$, that is determined by Queue Control.



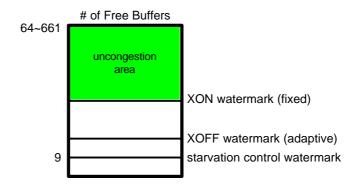
XON/XOFF Thresholds Selection

In VT6508, the XON threshold is fixed and programmable. Basically, the XON threshold is equal to half of total number of buffers -10.

The XOFF threshold is adaptive and has two objectives:

- (1) Prevent buffers are exhausted by some hungry port for unbalanced traffic condition. So, XOFF threshold >= total # of buffers necessary to be reserved.
- (2) Prevent packet drop because of too late response to buffer insufficience condition for balanced traffic condition. In VT6508, we assume the flow control response time to each port is 2 buffers.





Q[k]: queue length of output port k

 R_{max} : maximum reserved buffer number for each port

R[k]: reserved buffer number of port k, $R[k] = (Q[k] \ge R_{\text{max}} ? 0 : Q[k] - R_{\text{max}})$

 Ψ : # of buffers necessary to be reserved, $\Psi = \sum_{k=0}^{9} R[k]$

 Φ : free memory count

C: reserved memory count for virtual free space, C = 10 (default)

 Ω : virtual free buffer count, $\Omega = (\Phi \le C ? 0 : \Phi - C)$

Θ: XON threshold

When to enter XOFF window \Rightarrow ($\Omega = 0$)

or $(\Omega \le \max{\{\Psi,28\}} \text{ and } R[k] = 0)$

or (any port is XOFF and R[k] = 0)

When to leave XOFF window for any input ports $\Rightarrow \Omega > \Theta$

C: number of reserved buffers to prevent touching buffer starvation threshold (Virtual Free Space)

9 < C < D, default C = 10

D: offset of # buffers to tolerate incoming packets before source parties stop (Minimum XOFF Constant Threshold)

Default D = 2x9 + 9 + 1 = 28 for SRAM size of 256KB

£ XFree memory count

 \mathfrak{L} **Z**# of buffers necessary to be reserved for private policy.

XOFF condition:

 $((\pounds X_i \bigcirc A) + C)$ and $(R_k = 0)$ or $(\pounds X \bigcirc A)$

XON condition:



 $(\ (\pounds\ X\!\!\!> XON_Thred)\ XON_Thred = (1/2)\ x\ (MaxBufferNo + MaxXOFF - 20)$



SRAM size	# of buffers	Max XOFF	Min XOFF	XON	Rmax
128KB	64	D+C = 29	D+C = 29	43	1
		D = 19	D = 15		
		C = 10	C = 10		
256KB	149	4x9+C = 46	D+C = 38	98	4
			D = 28		
		C = 10	C = 10		
512KB	320	6x9+C = 76	D+C = 57	198	6
		C = 20	D = 37		
			C = 20		
1MB	661	8x9+C=102	D+C = 67	381	8
		C = 30	D = 37		
			C = 30		

2.5 Broadcast Storm Filtering

In VT6508, the feature of broadcast storm filtering is based on the above congestion control, and subjects to incoming broadcast packets. When an incoming broadcast packet is first enqueued by the Input Control of the source port, Queue Control will check if any one of its destination output ports is in XOFF window. If yes, it will trigger flow control to source port. For example, if a broadcast packet comes in from port 0, it will be forwarded to port1, port2, port3, port4, ... port8 in turn. If any one of ports 1~8 is congested, it will trigger flow control to source port 0, right on it is received and enqueued by Input Control of port 0.

If any port's drop window is turned on, Forwarding Control will drop the incoming broadcast packets to restrict the resource allocation.

2.6 SERIAL EEPROM INTERFACE AND CONFIGURATION COMMANDS

If the chip initialization mode is "EEPROM-init" (by strapping), VT6508 will read the external EEPROM device through EEC, EEIO during the system initialization. The signals of EEPROM interface are described as follows:

Name	Type	Description
EEC	0	serial EEPROM
		interface clock
EEIO	IO	serial EEPROM
		interface data signal

There are two types of EEPROM commands with the following format:

1. EOF Command: this must be the last command in EEPROM.



 $EE_DATA[7:0] = 8'b0000_0000$

- 2. Download Command:
 - EE_DATA[7:6] = 2'b01 → eeprom download data for a specific address
 - EE_DATA[5:0] → sequentially download data length from the following address
 - The next 2 EE_DATA will be Register base and offset address

Note that there must be an 8-bit check sum following the EOF command to make the sum of all commands and this check-sum byte is equal to zero bits 0000-0000.

EEPROM memory layout:

EEPROM Address	Data[7:0]
00H	- 1st Command[1:0](2 b01) + I^{st} _Data_Length[5:0] (1 <= n <= 63)
00H+1	- CS_Module[7:0]
00H+2	- IO_Write Start Address
<i>00H</i> +3 ∼	- IO_Write_DATA[7:0] *n
00H+2+1 st _Data_Length(n)	
00H+3+1 st _Data_Length(n)	- 2 nd Command
	- Ending Command[7:0]=8'b0000_0000 → End_of_file
	Check_Sum

Note that EEPROM device id and address must be 7'b1010_001.

2.7 TRUNKING

There are two trunk groups in VT6508, one is port (0,1), another is port (6,7). Each one can be enabled individually. The trunking load balance algorithm is as follows:

Assume DMAC=12'h000000000002, source_port_id = 4'b000<u>1</u>.

Assume the above value is 12'b000000011010 after lookup for DMAC = 12'b000000000000.



Let $dest_bit = DMAC_forwarding_entry_address[4]$ xor source_port_id[0]. If the DMAC refers to a trunk port of the trunk group (0,1), the selected outgoing trunk port is port $0 + dest_bit$.

If the DMAC refers to a trunk port of the trunk group (6,7), the selected outgoing trunk port is port 6 + dest_bit.



3 THE VT6508 SRAM ADDRESS MAPPING TABLE

128-pin 32bit mode (1 Bank of 32Kx32, 64Kx32):

Internal	6508 Pad(64bit mode)	6508 Package	SRAM Address Pin	SRAM chip
64bit		Address Pin		select
linear				
Address				
IA15	SA15 = IA14	SA15 =	A15 = IA14	
		IA14	N/A for 32kx32	
IA14	SA14 = IA13	SA14 =	A14 = IA13	
		IA13		
IA13	SA13 = IA12	SA13 =	A13 = IA12	
		IA12		
IA12	SA12 = IA11	SA12 =	A12 = IA11	
		IA11		
IA11	SA11 = IA10	SA11 =	A11 = IA10	
		IA10		
IA10	SA10 = IA9	SA10 = IA9	A10 = IA9	
IA9	SA9 = IA8	SA9 = IA8	A9 = IA8	
IA8	SA8 = IA7	SA8 = IA7	A8 = IA7	
IA7	SA7 = IA6	SA7 = IA6	A7 = IA6	
IA6	SA6 = IA5	SA6 = IA5	A6 = IA5	
IA5	SA5 = IA4	SA5 = IA4	A5 = IA4	
IA4	SA4 = IA3	SA4 = IA3	A4 = IA3	
IA3	SA3 = IA2	SA3 = IA2	A3 = IA2	
IA2	SA2 = IA1	SA2 = IA1	A2 = IA1	
IA1	SA1 = IA0	SA1 = IA0	A1 = IA0	
IA0	SA0 = 0	N/A	A0 = 0(directly tie	
	(internally)		to ground)	



SECTION II REGISTER MAP

1. REGISTERS TABLE

By using EEPROM initialization mode, all internal registers can be accessed.

	Register Description	Name	Bits	Default	R/W
(base/of				Value	
fset)	Chip Configuration				
	Chip Configuration by strapping:	CHIP_CFG	[63:0]	by	R/W
30 0711	CHIP_CFG[0]: BCAST_BPDU	J0, 0	[00.0]	strapping	, • •
	CHIP_CFG[1]: LED_COMB [0]			'' '	
	CHIP_CFG[2]: DROP_CONTROL_EN				
	CHIP_CFG[3]: SRAM_TYPE[1]				
	CHIP_CFG[5:4]: CHIP_INIT_MODE[1:0]				
	CHIP_CFG[7:6] RETRY_EXCE_COLL [1:0] CHIP_CFG[15:8]:RMII_SPEED_CFG[7:0]				
	CHIP_CFG[15.8].RMIII_SFEED_CFG[7.8]				
	CHIP_CFG[17] TRUNK_MODE[2]				
	CHIP_CFG[18]: EN_BACKPRESSURE [0]				
	CHIP_CFG[20:19]:RMII_PORT_FC_EN				
	CHIP_CFG[23:21]:LED_CFG[2:0]				
	CHIP_CFG[24]: HASH_ALG_SEL				
	CHIP_CFG[25]: NO_SWAP_EEPROM_LED				
	CHIP_CFG[26]: EN_AGING CHIP_CFG[27]: DIS_CUT_THRU				
	CHIP_CFG[28]: DIS_AGGRES_BACKOFF				
	CHIP_CFG[29]: COLL_POINT				
	CHIP_CFG[30]: DIS_TEST_MODE				
	CHIP_CFG[31]: DIS_LATCH_UP_MODE				
	CHIP_CFG[39:32]: MONITOR_PORT[7:0]				
	CHIP_CFG[41:40]: SNIFFER_PORT[1:0]				
	CHIP_CFG[43]: TRUNK_MODE[1]				
	CHIP_CFG[45]: TRUNK_MODE[3] CHIP_CFG[42]: RMII_SPEED_CFG[9]				
	CHIP_CFG[44]: RMII_SPEED_CFG[8]				
	CHIP_CFG[46]: SRAM_NUM				
	CHIP_CFG[47]: SRAM_TYPE[0]				
	CHIP_CFG[48]: DIS_DROP_STARV				
	CHIP_CFG[51:49]: MII_PORT_CFG				
	CHIP_CFG[52]: SRAM_32BIT_MODE				
	CHIP_CFG[53] DIS_SHORT_PREAM				
	CHIP_CFG[54]: PHY_ADDRESS_SEQ CHIP_CFG[55]: RMII_SPEED_ARB				
	CHIP_CFG[55]: RMII_SPEED_ARB CHIP_CFG[56]: DIS_9_PORT_LED_MODE				
	CHIP_CFG[62:57]: (reserved with value 1's)				
	CHIP_CFG[63]: DIS_INIT_DBG_MODE				



0.411	Froble Normal Operation of LED Cariel Out	EN LED OUT	[0]	10	D/M
0AH	Enable Normal Operation of LED Serial-Out	EN_LED_OUT	[0]	0	R/W
	The normal operation (periodic data sending) of				
	LED Serial-Out can be triggered by (1)				
	EEPROM-init complete signal from EEPROM-init				
	module, or (2) software writes value 1 to this				
4011	register.	0011 0055 05055			5 0 4 /
10H	CPU Soft Reset for the whole switch chip reset	CPU_SOFT_RESET	[0]	1	R/W
	For Read, 0: soft reset in progress 1: soft reset done				
	For Write, any value will trigger the whole chip				
	reset. The soft reset is similar to power-on reset				
	for the switch chip, except that it is asserted by				
	writing any value to this register. The CPU soft				
	reset has to take 16 RCLK50 cycles, i.e. 320ns.				
	Queue control				
00-01H	Congestion Control Xon Threshold register	XON_THRED	[9:0]	by strapping	R/W
02H	Output Private Buffer Size	PRIVATE_BUF_SIZE	[4:0]	by	R/W
		0115115 15		strapping	5 0 4 /
03H	Selected Output Queue ID	QUEUE_ID	[3:0]	0	R/W
	Selected Output Queue Length	OUEUE_LENGTH	[9:0]	0	R/W
06-07H	Selected Output Queue Head Pointer	OUEUE_HEAD	[9:0]	3FF (null)	R/O
08-09H	Virtual Free Space	V_FREE_SPACE	[9:0]	by SRAM size	R/W
0A-0BH	Minimum XOFF Constant Threshold	MIN_XOFF_CONST	[9:0]	by SRAM	R/W
		_THRED	_	size	
	Buffer control				
00-01H	Free Memory Block Count	FREEMCNT	[9:0]	by SRAM size	R/W
02H	Memory Allocation Bit Mask Byte ID	MASK_ID	[6:0]	0	R/W
03H	Memory Allocation Bit Mask	BUFFER_MASK	[7:0]	0	R/O
0300H	Forwarding table control				
00H	sniffer port id	SNIFFER_PID	[3:0]	by strapping	R/W
01-02H	monitor ports bit mask	MONITOR_PM	[8:0]	3	R/W
03H	monitor ports bit mask	MONITOR_MODE	[1:0]	0	R/W
	0 => to monitor input&output traffic (default)				
	1 => to monitor input traffic only				
	2 => to monitor output traffic only				
04H	auto-aging configuration	EN_AUTO_AGE	[0]	0	R/W
	0: disable auto-aging (default)				
	1: enable auto-aging Note: Initialization Control should enable				
	auto-aging for strapping-only				
	initialization.				
05H	inter-aging time (in millisecond)	INTER_AGE_TIME	[7:0]	50	R/W
	The velid value is 1, 255				
	The valid value is 1255. The default inter-aging time value is 50ms for an				
	entry by default strapping so that the forwarding				
	table of 4K 2K entries (4K slots) can be aged				
	completely in about 5 minutes.				
ı	Journal of the land of the lan	I	I	I	I



06-07H	aging index	AGING_INDEX	[10:0]	0	R/W
00 0111	06H: AGING_INDEX[7:0]	7.0	[10.0]		"
	07H: AGING_INDEX[10:8]				
	for CPU to trigger aging on specific entry				
08H	aging status	AGING_STATUS	[0]	0	R/W
	0: idle or complete				
	1: in progress (no matter how the process is				
	triggered by auto-init or by CPU)				
	for CPU to detect status of aging process				
09H	hash algorithm	HASH_ALG	[1:0]	0 if	R/W
	HASH_ALG[0]: 0, crc11, 1, direct map			strapping	
	HASH_ALG[1]: 0, scramble, 1, no scramble			SD[24]=1	
				,	
	The default setting 2b' 00 is optimized specially				
	for Smartbit test, each combination may just be			3 if	
	as good as others.			strapping	
				SD[24]=0	
	HASH_ALG[1:0]=2b' 11: direct map without				
	scramble, hash key = MAC[10:0]				
	HASH_ALG[1:0]=2b' 01: direct map with				
	scramble, hash key = MAC[7:0, 16:18]				
	HASH_ALG[1:0]=2b' 10: CRC map & no				
	scramble, hash key = CRC11{MAC[47:0]}				
	HASH_ALG[1:0]=2b' 00: CRC map with				
	scramble, hash key =				
	CRC11{MAC[47:24,15:0,16:23]}				
	The forwarding table entry format is as follows.				
	[58] static				
	[57:56] age count				
	[49:40] forwarding portmask				
	[39:0] tag				
	Regardless of the hash algorithm, the tag will				
	always use MAC[47:8]. The mapping is as				
	follows.				
	Tag[39:32] \longleftrightarrow MAC[15:8]				
	Tag[31:24] \longleftrightarrow MAC[23:16]				
	Tag[23:16] ←→ MAC[31:24]				
	Tag[15:8] \longleftrightarrow MAC[39:32]				
0.4	Tag[7:0] ←→ MAC[47:40]	D1440 02 :	F 4 =		D # * *
0A-	destination MAC object	DMAC_OBJ	[47:0]	0	R/W
0FH	0AH: DMAC[47:40],				
	0BH: DMAC[39:32],				
	0FH: DMAC[7:0].				
	A write to this register will make the				
	corresponding CRC-map lower 11-bit hash key				
	and upper 37-bit tag shown in HASH_CONTENT.				



10	calculation status of DMAC's hash key	HASH_CAL_STATU	[0]	0	R/W
	The HASK_CAL_STATUS register indicates if the content in HASH_KEY registers are valid. 1 for calculation in progress, 0 for calculation complete. The CRC-map result of DMAC_OBJ is shown in	S			
	this register, 11-bit hash key. The corresponding				
	tag is always DMAC[47:9].	114 011 1(5)	[40.0]		D 044
11-12H	hash key of DMAC_OBJ	HASH_KEY HASH_RESULT	[10:0]	0	R/W
	The CRC-map result of DMAC_OBJ is shown in this register, 11-bit hash key. The corresponding tag is always DMAC[47:9].				
	port mask for lookup-miss counter 1: to be accounted into LOOKUP_MISS counter 0: not to be accounted LOOKUP_MISS_MASK [7:0] at 13H. LOOKUP_MISS_MASK [9:8] at 14H.	LOOKUP_MISS_MA SK	[9:0]	1FF	R/W
	lookup-miss counter	LOOKUP_MISS	[15:0]	0	R/W
	15H: LOOKUP_MISS [7:0] 16H: LOOKUP_MISS [15:8] This counter is cleared automatically after CPU read.				
17-18H	hash-conflict counter	HASH_CONFLICT_C NT	[15:0]	0	R/W
	17H: HASH_CONFLICT_CNT [7:0] 18H: HASH_CONFLICT_CNT [15:8] This counter is cleared automatically after CPU read.				
19-1AH	destination port mask for lookup-miss unicast packets	UCAST_DPM	[9:0]	01FF	R/W
	19H: UCAST_DPM [7:0] for ports 70 1AH: UCAST_DPM [9:8] for cpu port & port 8 for unicast packet lookup miss broadcast control				
1B-1CH	destination port mask for lookup-miss multicast packets		[9:0]	01FF	R/W
	1BH: MCAST_DPM [7:0] 1CH: MCAST_DPM [9:8] for multicast packet lookup miss broadcast control				
1DH	selection of CPU_PM usage	CPU_PM_USER_SE L	[0]	0	R/W
	select cpu port to use CPU_PM for lookup select MII port to use CPU_PM for lookup				



			I	1-	I I
	Destination Ports Bit Mask for Port 8 (CPU port) 0: no affection on port 8's incoming packets that will be forwarded with lookup not 0: port 8's next incoming packets will be forwarded by CPU_PM, rather than lookup. CPU_PM will be set as '0' after processing exact ONE port 8's incoming packet.	СРО_РМ	[8:0]	0	R/W
	Note that this function can facilitate the				
	implement of some special system features, e.g.				
0011	production testing utility, boot diagnosis utility.	ODEOLAL OTI	[7.0]	0	DAM
	Reserved Special Internal Control bit 0 — enable early cpu port broadcast filtering bit 3:1 — unused	SPECIAL_CTL	[7:0]	0	R/W
	bit 4 - DROP_CONGESTION_WHEN_STARV bit 6:5 - SWITCH_MAC_MATCH_MODE bit 7 - LOOPUP_OPTION				
	Note that only the bit 0 is used for normal operation setting. Except the bit 0 is used for the				
	normal operation setting, the other bits are used only in Verilog simulation.				
21H	CPU port related forwarding configuration	CPU_FWD_CFG	[2 3:0]	0	R/W
	bit 0 – enable forwarding broadcast packets with DMAC=0xfffffffffff to CPU (default = 0 : disable) bit 1 – enable forwarding spanning-tree packets to CPU (default = 0 : disable) bit 2 – enable forwarding unicast packets with DMAC = switch MAC base to CPU (default = 0 : disable) bit 3 – enable forwarding multicast packets with DMAC = 01-80-C2-00-00-02~0F to CPU (default = 0 : disable forwarding to CPU port, i.e. drop these packets)	61 6_1 WB_61 6	[2 0.0]		
	Note that the USER_PM in VT3061 is changed as 10-bit signal from 10 output port enabling bits.				
	The three register bits are used to enable/disable forwarding the above three types of frames to the CPU port. If CPU_FWD_CFG[1]=1, i.e. forwarding spanning-tree packets to CPU is disabled, these SPT packets will be broadcasted, just like the case for Ethernet Hub.				
22-23H	Drop On	DROP_ON	[8:0]	0	R/O
	Value 1: the corresponding port's drop feature is enabled. Any incoming packets destining to any one port in congestion will be dropped by Forwarding Control by setting port mask = 0. Value 0: the port's drop feature is disabled. This signal is from TMAC.				



24-25H	Congestion Window On	CONGEST_ON	[9:0]	0	R/O
	Value 1: the port is in congestion status.				
	Value 0: the port is not in congestion status. This signal is from Queue Control.				
26-27H	XOFF Window	XOFF_WINDOW	[9:0]	0	R/O
20-2711	XOTT WINDOW	NOTT_WINDOW	[3.0]	0	14/0
	Value 1: the port is in XOFF window.				
	Value 0: the port is in XON window.				
	This signal is from TMAC.				
30H	spanning tree state for PORT 0	PORT0_STP_STATE	[1:0]	3	R/W
	2' b00 - blocking state				
	2' b01 – listening state				
	2' b10 - learning state				
	2' b11 – forwarding state (default)				
	The default value of all STP states is "forwarding				
	state".				
	The famous discussion in some Fabruary to and in				
	The forwarding operation in each Ethernet port is controlled by its associated spanning tree state.				
	In blocking or listening state, the incoming				
	packets will not trigger any DMAC lookup				
	operation and SMAC learning operation.				
	In learning state, the incoming packets will not				
	trigger DMAC lookup operation, but the SMAC				
	learning operation will be triggered for CRC-OK				
	packets.				
	Only in forwarding state, an incoming packet will				
	trigger DMAC lookup operation while the first 24				
	bytes are received, and it will trigger the SMAC				
	learning operation while the whole packet is received with good CRC.				
	For the 802.1d spanning tree algorithm, a				
	blocked port for loop avoidance should enter the				
	blocking state so that any incoming packets are				
	filtered without forward. A normal port that does				
	not cause any loop should be in the forwarding				
	state.	2027/ 072 07/77			- n.,
31H 32H	spanning tree state for PORT 1	PORT1_STP_STATE		3	R/W R/W
33H	spanning tree state for PORT 2 spanning tree state for PORT 3	PORT2_STP_STATE PORT3_STP_STATE		_	R/W
34H	spanning tree state for PORT 4	PORT4_STP_STATE		3	R/W
35H	spanning tree state for PORT 5	PORT5_STP_STATE		3	R/W
36H	spanning tree state for PORT 6	PORT6_STP_STATE		3	R/W
37H	spanning tree state for PORT 7	PORT7_STP_STATE		3	R/W
38H	spanning tree state for PORT 8	PORT8_STP_STATE	[1:0]	3	R/W
0400H	PHY control				
00H	CPU command port ID	CMD_PORTID	[3:0]	0	R/W
	Normally, it is the working port ID for CPU direct				
	read/write. It also shows the working port ID for				
	auto-polling debugging mode.			<u></u>	
01H	CPU command PHY REG number	PHY_REG_ADDR	[4:0]	0	R/W



02-03H	CPU command R/W data	PHYDATA	[15:0]		R/W
	The last data word read from PHY is stored in this				
	register for CPU read.				
	Note that the last data word written to PHY is				
	stored in an internal register that is also located in				
	this address for CPU write, but not readable.				
	CPU command register	PHYCMD	[3:0]		W/O
	PHYCMD[3]: enable phy auto-polling PHYCMD[2]: disable phy auto-polling				R/W
	PHYCMD[1]: phy read command.				
	PHYCMD[0]: phy write command.				
	When auto-polling mechanism is enabled, CPU				
	can not read/write PHY devices, even though the				
	PHY (auto-polling) Debug Mode is enabled.				
	PHY status register	PHYSTS	[2:0]	0	R/O
	Bit-0: 0=cmd_idle; 1=cmd_busy		1		
	Bit-1: 0=cmd_incomplete; 1=cmd_complete Bit-2: 0=autopoll_off; 1=autopoll_on				
	ibit-z. 0=autopoii_oii, i=autopoii_oii		1		
	Note that the "cmd" means the command from		1		
	CPU interface, not from auto-polling control.				
06-07H	Link status change	LINK_STATUS_CHA	[8:0]	0	R/O
	1' b1=link status changed	NGE			
	1' b0=link status without changed				
	(cleared after cpu read) Link status	LINIZ CTATLIC	[0.0]	0	D/O
08-09H	1' b0=link down (default)	LINK_STATUS	[8:0]	0	R/O R/W
	1' b1=link up				17,77
0AH	PHY debug mode enable	PHY_DBG_EN	[0]	0	R/W
	PHY_DBG_EN: enable debug auto-polling state		-		
	Machine for validation				
	Value 1: enabled				
0BH	Value 0: disabled	TDIC CTED DOLL	[0]	0	R/W
ЛРЦ	Trigger Auto-Polling Single Step	TRIG_STEP_POLL	[0]	U	IK/VV
	Write value 1 to trigger a single step, then this bit				
	is self-cleared after command complete.				
	This operation may takes about 30 s.				
	PHY Device Address	PHY_ADDR	[4:0]	0	R/W
	EEPROM control				
00H	EEPROM word address	EEWDADDR	[7:0]		R/W
	For a 256-byte FEDDOM dovice, an 8 bit data		1		
	For a 256-byte EEPROM device, an 8-bit data object is identified with this register. For a 512-		1		
	byte EEPROM device, an 8-bit data object is		1		
	identified with this register plus EEDEVADDR[1],		1		
	vice versa.				
01H	EEPROM data	EEDATA	[7:0]		R/W
	The last data bute read from EEDROM is stored		1		
	The last data byte read from EEPROM is stored in this register for CPU read.		1		
	Note that the last data byte written to EEPROM is		1		
	stored in an internal register that is also located in		1		
	this address for CPU write, but not readable.		ĺ		1



02H	EEPROM device address	EEDEVADDR	[7:0]		R/W
0211	bit 7-4 : device type id (default as1010)	LLDLVADDIA	[,, .0]		17,77
	bit 3-1 : device id				
	bit 0 : r/w command →				
	value 0: write; value 1: read				
03H	EEPROM status register	EESTS	[3:0]		R/O
0011	EESTS[0]: value 0: idle, 1: busy (r/w or auto-init		[0.0]		1.0
	in progress)				
	EESTS[1]: value 1: r/w complete without error				
	(cleared by register read)				
	EESTS[2]: value 1: r/w ack error				
	EESTS[3]: value 0: no init error, 1: init err				
0600H					
	CPU interface	IDOCTO	[5.0]	0	DAA
00H	interrupt status/clear register	IRQSTS	[5:0]	0	R/W
	bit 0: PHY command complete interrupt		[6:0]		
	bit 1: EEPROM command interrupt				
	bit 2: PHY auto-polling link status change				
	notification				
	bit 3: EEPROM initialization complete interrupt				
	bit 4: interrupt indication for CPU IO port				
	receiving an incoming packet from some				
	Ethernet port, that is ready to be received				
	by CPU via IDE.				
	bit 5: interrupt indication for CPU IO port finishing				
	the transmission of an outgoing packet that				
	was sent by CPU via IDE.				
	bit 6: interrupt indication for re-autonegotiation				
	caused by port speed mismatch				
	Nieta that conition of the than a surrounding or a sister.				
	Note that writing 1 to the corresponding register bits will clear that interrupts.				
01H	interrupt mask register	IRQMASK	[5:0]	0	R/W
5111	bit 0: PHY command complete interrupt mask	II (GIVI) (OI)	[6:0]	ľ	13/ ۷۷
	bit 1: EEPROM command interrupt mask		[[0.0]		
	bit 2: PHY auto-polling link status change				
	notification mask				
	bit 3: EEPROM initialization complete interrupt				
	mask				
	bit 4: interrupt mask for CPU IO port				
	receiving an incoming packet from some				
	Ethernet port, that is ready to be received				
	by CPU via IDE.				
	bit 5: interrupt mask for CPU IO port finishing				
	the transmission of an outgoing packet that				
	was sent by CPU via IDE.				
	bit 6: interrupt indication for re-autonegotiation				
	caused by port speed mismatch				
	Note: value 0 means "masked", value 1 means				
	"not masked". The default is "masked" so that				
	software has to poll IRQSTS to check if there are				
	any pending interrupts.				
	lany bending interrupts.				



0011	Inches address as white	CDAMADDD	[40:0]		DAM
02H- 04H	sram address register	SRAMADDR	[16:0]		R/W
0411	The data object addressed is in unit of 64 bits. The maximum allowable SRAM size is 1MB. For SRAM direct access, the CPU has to (1) set				
	SRAMADDR & SRAMDATA, (2) issue read/write command by SRAMCMD, and (3) check				
	command status by SRAMSTS.	001110171	100.01		5 044
05H- 0CH	sram data register	SRAMDATA	[63:0]		R/W
0DH	sram command register 2' b00 : nop 2' b01 : read 2' b10 : write	SRAMCMD	[1:0]		R/W
0EH	sram status register 2' b01 : read/write command done 2' b10 : busy (read/write in progress) 2' b00 : idle	SRAMSTS	[1:0]	0	R/O
0FH	Interrupt Mode INTRQ_MODE == 0 => High Active INTRQ_MODE == 1 => Low Active (default)	INTRQ_MODE	[0]	1	R/W
10H	Write packet command 3' b100: end of frame with the remaining data size = 2 byte 3' b101: end of frame with the remaining data size = 1 byte 3' b000: idle 3' b001: start of frame for the next write 3' b010: middle of frame for the next write 3' b011: abort the unfinished packet write CPU should write this command register before repeatedly writing 8 bit packet data to CPU Interface Control via the 8-bit IDE bus. The cpu interface with external is 8bit bus. However the interface between cpuif and cpuioctl is 16bit data bus. The WR_PKT_CMD will be passed to cpuioctl so that it will know all the 16bit are valid or only 8 bit are valid data. For cpu to send packet, the steps are as follows: write WR_PKT_CMD = 3'b001 to notify start of frame, then write first 2 byte data via IDE. write WR_PKT_CMD = 3'b010 to notify middle of frame, then write even number of data via IDE. Leave last two or one byte data. write WR_PKT_CMD = 3'b100 or 3'b101 to notify end of frame with remaining data 2 byte or 1 byte. If it is 2 byte, write the last two byte data first, and then write another byte for dummy padding.		[2:0]		W/O R/W
11H	Packet Abort Writing 1' b1 to this register will drop an incoming packet ready to be read by CPU. Reading this register will clear this register as 1' b0.	PKT_ABORT	[0]		R/W



12H	Packet Data Register	PKT_DATA	[7:0]		R/W
	Write this register (need to cooperate with WR_PKT_CMD at 0610H and WR_PKT_STATUS at 1910H) to send an incoming-to-switch packet to CPU Input Control. Read this register (need to cooperate with PKT_BYTE_CNT at 1900-1901H and RD_PKT_STATUS at 1902H) to retrieve an outgoing-from-switch packet from CPU Output Control.				
	Note that any read or write to this register will not cause auto-increment of the Address Register to support the continuous packet-byte read/write commands.				
	The cpu interface with external is 8bit bus. However the interface between cpuif and cpuioctl is 16bit data bus. Reading packet data will be served by cpuioctl so that it will know all the 16bit are valid or only 8 bit are valid data. For cpu to read packet, please read one more dummy byte				
	after reading the last byte if the packet length is				
13H	even.	SWITCH_MAC	[7:0]	00H	R/W
14H	bits [47:40] of switch MAC address [47:0] bits [39:32] of switch MAC address [47:0]	SWITCH_MAC	[7:0]	40H	R/W
15H	bits [31:24] of switch MAC address [47:0]	SWITCH_MAC	[7:0]	63H	R/W
16H	bits [23:16] of switch MAC address [47:0]	SWITCH_MAC	[7:0]	80H	R/W
17H	bits [15:8] of switch MAC address [47:0]	SWITCH_MAC	[7:0]	00H	R/W
18H	bits [7:0] of switch MAC address [47:0]	SWITCH MAC	[7:0]	00H	R/W
19H	trigger/check SRAM auto test	SRAM_AUTO_TEST	[2:0]	0	R/W
	Write to this register will trigger the SRAM autotest function. The status of execution of the SRAM autotest function is reported in this register, defined as follows: bit 2: $1 \rightarrow$ in progress, $0 \rightarrow$ complete or idle bit 1: $1 \rightarrow$ error in single r/w connectivity, $0 \rightarrow$ no error in single r/w connectivity bit 0: $1 \rightarrow$ error in burst r/w connectivity, $0 \rightarrow$ no error in burst r/w connectivity		r1		
20H	Revision ID	REVISION_ID	[7:0]	0	R/O
1000H	MAC & I/O Control Module of Port 0		_		
00H	configurable preamble bytes	PREAM_CFG	[2:0]	7	R/W
01H	configurable frame gap in di bits for 1st interval	IFG_CFG	[5:0]	32	R/W



0011	Dockoff configuration	DOLLCEO	[4.0]	E' 54000	D/M/
02H	Backoff configuration	BOFFCFG	[4:0]		R/W
	bit 0: CAP mode → DEC's capture effect			0	
	(default: CAP=0)				
	bit 1: MBA mode → HP's capture effect				
	(default: MBA=0)				
	bit 2: EEFAST mode → drop the 2 nd collided				
	packet for testing purpose, accelerate the drop				
	event (default: EEFAST=0)				
	bit 3: CRANDOM mode → use another random				
	algorithm (default: CRANDOM=0)				
	bit 4: OFSET → parameter for backoff timer,				
	For OFSET=1, the TMAC will follow the 802.3				
	standard backoff algorithm. For OFSET=0, the				
	TMAC will select the backoff time for the 1 st and				
	2 nd collision as that of the 3 rd collision, i.e. the				
	possible the backoff time for the 1 st and 2 nd				
	collision is ranged from 0 to 7 in unit of slot time.				
	(default: OFSET=1)				
	For the number of excessive retries, please see				
0011	the strapping pin SD[52:51].	1440050	[0.0]		D // / /
03H	MAC media type configuration	MACCFG	[3:0]	0	R/W
	bit 0: SPD_10M → value 0: 100Mbps,				
	value 1: 10Mbps				
	bit 1: HALF_DPX → 1: half duplex, 0: full duplex				
	bit 2: RCV_FC_DIS →				
	1: disable receive flow control frame				
	0: enable receive flow control frame				
	bit 3: XMT_FC_DIS →				
	1: disable send flow control frame for				
	full-duplex mode				
	0: enable send flow control frame				
	U. enable send now control frame				
	Niete dele le configura d'action d'alle le conte				
	Note: this is configured automatically by auto-				
	polling control or CPU if auto-polling is				
	disabled to that port. There is no forced				
	mode if auto-polling is enabled to that port.				
04H	IO port enable	IO_CFG	[2:0]	0	R/W
	bit 0: input port enable				
	0: input disabled, (drop incoming packets)				
	1: input enabled				
	bit 1: output port enable				
	0: output disabled (default)				
	1: output enabled				
	bit 2: output port dequeue hold				
	1: dequeue disabled				
	0: dequeue enabled (default)				
	Note that the broadcast domain is determined by				
	the CPUIO_CFG[1] and 9 bits of all Ethernet				
	ports' IO_CFG[1]. If IO_CFG[1]==0, this				
	Ethernet port must be not within the broadcast				
	domain.				
	<u> </u>		1	1	



	Granted 1 st free buffer Addr in Input Control bit [10:0]: buffer address (ID*3 + 40h) bit [11]: buffer address valid (default: invalid) bit [12]: buffer free (available for the new incoming packet)	FREE_BUF_ADDR1	[12:0]	0	R/O
07-08H	Granted 2 nd free buffer Addr in Input Control bit [10:0]: buffer address (ID*3 + 40h) bit [11]: buffer address valid (default: invalid) bit [12]: buffer free (available for the new incoming packet)	FREE_BUF_ADDR2	[12:0]	0	R/O
09-0AH	Event mask for counters EVENT_MASK[0]: rx CRC good packets for Counter 1 EVENT_MASK[1]: rx flow control frames for Counter 1	EVENT_MASK	[14:0]	7FFF	R/W
	EVENT_MASK[2]: rx a runt packet for Counter 2 EVENT_MASK[3]: rx an over-size packet for Counter 2 EVENT_MASK[4]: rx a regular-size but CRC error packet for Counter 2				
	EVENT_MASK[5]: FIFO over-run without buffer starvation for Counter 3 EVENT_MASK[6]: FIFO over-run with buffer starvation for Counter 3 EVENT_MASK[7]: intended drop local packets for Counter 3 EVENT_MASK[8]: drop broadcast packets due to congestion for Counter 3 EVENT_MASK[9]: drop unicast packets due to congestion for Counter 3				
	EVENT_MASK[10]: tx packets excluding flow control packets for counter 4 EVENT_MASK[11]: tx flow control packets for counter 4				
	- EVENT_MASK[12]: FIFO under-run for Counter 5 EVENT_MASK[13]: collision for Counter 5 EVENT_MASK[14]: first 16-bit packet data late (can not catch up preamble) for counter 5				
10H- 11H	received good packet count (Counter 1)	RCV_GOOD_PKT	[15:0]	0	R/O
12H- 13H	It is cleared automatically after CPU read. received bad packet count (Counter 2)	RCV_BAD_PKT	[15:0]	0	R/O
14H- 15H	It is cleared automatically after CPU read. drop packet counter(Counter 3)	DROP_PKT	[15:0]	0	R/O
	It is cleared automatically after CPU read.				



16H- 17H	sent good packet count(Counter 4)	XMT_GOOD_PKT	[15:0]	0	R/O
.,	It is cleared automatically after CPU read.				
18H- 19H	sent bad packet counter(Counter 5)	XMT_BAD_PKT	[15:0]	0	R/O
	It is cleared automatically after CPU read.				
1AH- 1BH	sent abort packet counter(Counter 6) (abort due to excessive transmission fails)	XMT_ABORT_PKT	[15:0]	0	R/O
	It is cleared automatically after CPU read.				
1C- 1DH	dequeued 1 st buffer Addr in Output Control bit [10:0]: buffer address (ID*3 + 40h) bit [11]: buffer address valid (default: invalid)	DEQUEUED_BUF_A DDR1	[11:0]	0	R/O
1E- 1FH	dequeued 2 nd buffer Addr in Output Control bit [10:0]: buffer address (ID*3 + 40h) bit [11]: buffer address valid (default: invalid)	DEQUEUED_BUF_A DDR2	[11:0]	0	R/O
20H	Internal state of Input Control (reserved for Weipin's definition) (only valid for port 0)	INPUT_CTL_STATE	[7:0]	0	R/O
21H	Internal state of Output Control (reserved for Weipin's definition) (only valid for port 0)	OUTPUT_CTL_STA TE	[7:0]	0	R/O
1100H	MAC & I/O Control Module of Port 1	same as Port 0			
1200H	MAC & I/O Control Module of Port 2	same as Port 0			
1300H	MAC & I/O Control Module of Port 3	same as Port 0			
1400H	MAC & I/O Control Module of Port 4	same as Port 0			
1500H	MAC & I/O Control Module of Port 5	same as Port 0			
1600H	MAC & I/O Control Module of Port 6	same as Port 0			
1700H	MAC & I/O Control Module of Port 7	same as Port 0			
1800H	MAC & I/O Control Module of Port 8	same as Port 0			
1900H	CPU I/O Control Module				
00H	CPU packet read byte count register bits [7:0] CPU can check the incoming (ready to be received by CPU via IDE) packet length via the 11-bit register PKT_BYTE_CNT [10:0] before starting to read it.	PKT_BYTE_CNT	[7:0]	0	R/O
01H	CPU packet read byte count register bits [10:8]	PKT_BYTE_CNT	[10:8]	0	R/O
02H	CPU packet read status register 2' b00: idle or packet read (by CPU) in progress 2' b01: packet read (by CPU) complete without error 2' b10: packet read (by CPU) with error (CPU needs to read the same packet again)	RD_PKT_STATUS	[1:0]	0	R/O
03H	Packet source port ID CPU can check the incoming packet's source port ID via the 3-bit register PKT_SRC_PORT before starting to read it. It is useful to the spanning tree algorithm.	PKT_SRC_PORT	[3:0]	0	R/O



T	ı			
		[2:0]	0	R/W
bit 0 : input port enable → 1: input enable, 0: input				
disable				
bit 1 : output port enable → 1: output enable, 0:				
output disable				
bit 2 : output port dequeue hold → 1: dequeue				
disabled, 0: dequeue enabled				
ports' IO_CFG[1]. If CPUIO_CFG[1]==0, the				
CPU port must be not within the broadcast				
domain.				
Granted (at most one) free buffer Addr in	FREE_BUF_ADDR	[11:0]	0	R/O
Input Control				
Bit [10:0]: buffer address (ID*3 + 40h)				
dequeued buffer Addr in Output Control		[11:0]	0	R/O
bit [10:0]: buffer address (ID*3 + 40h)	DDR			
bit [11]: buffer address valid (default: invalid)				
CPU packet write status register	WR_PKT_STATUS	[2:0]	0	R/O
bits [1:0]: packet write status				
2' b00: idle or packet write in progress]		
2' b01: CPU sent packet successfully				
·				
(CPU needs to re-write the]		
]		
]		
, , ,				
,]		
displaying XON/XOFF status in		1	1	
XOFF_WINDOW[9] of Forwarding Control.				
	disable bit 1 : output port enable → 1: output enable, 0: output disable bit 2 : output port dequeue hold → 1: dequeue disabled, 0: dequeue enabled Note that the broadcast domain is determined by the CPUIO_CFG[1] and 9 bits of all Ethernet ports' IO_CFG[1]. If CPUIO_CFG[1]==0, the CPU port must be not within the broadcast domain. Granted (at most one) free buffer Addr in Input Control Bit [10:0]: buffer address (ID*3 + 40h) Bit [11]: buffer address valid dequeued buffer Addr in Output Control bit [10:0]: buffer address (ID*3 + 40h) bit [11]: buffer address valid (default: invalid) CPU packet write status register bits [1:0] : packet write status 2' b00: idle or packet write in progress 2' b01: CPU sent packet successfully (CPU needs to re-write the packet again) Note that CPU packets have highest priority and will not be dropped by Forwarding Control. bit 2: CPU Input Control is ready for CPU to write packets (It can be ready only after setting CPUIO_CFG[0] = 1.) 0: not ready (default) 1: ready	bit 0 : input port enable → 1: input enable, 0: input disable bit 1 : output port enable → 1: output enable, 0: output disable bit 2 : output port dequeue hold → 1: dequeue disabled, 0: dequeue enabled Note that the broadcast domain is determined by the CPUIO_CFG[1] and 9 bits of all Ethernet ports' IO_CFG[1]. If CPUIO_CFG[1]==0, the CPU port must be not within the broadcast domain. Granted (at most one) free buffer Addr in Input Control Bit [10:0]: buffer address (ID*3 + 40h) Bit [11]: buffer address valid dequeued buffer Addr in Output Control bit [10:0]: buffer address valid (default: invalid) CPU packet write status register bits [1:0] : packet write status 2' b00: idle or packet write in progress 2' b01: CPU sent packet unsuccessfully (CPU needs to re-write the packet again) Note that CPU packets have highest priority and will not be dropped by Forwarding Control. bit 2: CPU Input Control is ready for CPU to write packets (It can be ready only after setting CPUIO_CFG[0] = 1.) 0: not ready (default) 1: ready Note that if CPU IO is in XOFF window, it will be not ready with WR_PKT_STATUS[2]==1. CPU IO always uses flow control mechanism with	bit 0 : input port enable → 1: input enable, 0: input disable bit 1 : output port enable → 1: output enable, 0: output disable bit 2 : output port dequeue hold → 1: dequeue disabled, 0: dequeue enabled Note that the broadcast domain is determined by the CPUIO_CFG[1] and 9 bits of all Ethernet ports' IO_CFG[1]. If CPUIO_CFG[1]==0, the CPU port must be not within the broadcast domain. Granted (at most one) free buffer Addr in Input Control Bit [10:0]: buffer address (ID*3 + 40h) Bit [11]: buffer address valid dequeued buffer Addr in Output Control bit [10:0]: buffer address (ID*3 + 40h) bit [11]: buffer address (ID*3 + 40h) bit [11]: buffer address valid (default: invalid) CPU packet write status register bits [1:0] : packet write status 2' b00: idle or packet write in progress 2' b01: CPU sent packet successfully (CPU needs to re-write the packet again) Note that CPU packets have highest priority and will not be dropped by Forwarding Control. bit 2: CPU Input Control is ready for CPU to write packets (It can be ready only after setting CPUIO_CFG[0] = 1.) 0: not ready (default) 1: ready Note that if CPU IO is in XOFF window, it will be not ready with WR_PKT_STATUS[2]==1. CPU IO always uses flow control mechanism with	bit 0 : input port enable → 1 : input enable, 0 : input disable bit 1 : output port enable → 1 : output enable, 0 : output disable bit 2 : output port dequeue hold → 1 : dequeue disabled, 0 : dequeue enabled Note that the broadcast domain is determined by the CPUIO_CFG[1] and 9 bits of all Ethernet ports' IO_CFG[1]. If CPUIO_CFG[1]==0, the CPU port must be not within the broadcast domain. Granted (at most one) free buffer Addr in Input Control Bit [10:0]: buffer address (ID*3 + 40h) Bit [11]: buffer address valid dequeued buffer Addr in Output Control bit [10:0]: buffer address (ID*3 + 40h) bit [11]: buffer address valid (default: invalid) CPU packet write status register bits [1:0] : packet write status 2' b00: idle or packet write in progress 2' b01: CPU sent packet successfully 2' b10: CPU sent packet unsuccessfully (CPU needs to re-write the packet again) Note that CPU packets have highest priority and will not be dropped by Forwarding Control. bit 2: CPU Input Control is ready for CPU to write packets (It can be ready only after setting CPUIO_CFG[0] = 1.) 0: not ready (default) 1: ready Note that if CPU IO is in XOFF window, it will be not ready with WR_PKT_STATUS[2]==1. CPU IO always uses flow control mechanism with



SECTION III ELECTRICAL SPECIFICATIONS ABSOLUTE MAXIMUM RATINGS

Parameter	Min	Max	Unit
Ambient operating temperature	0	70	oC.
Case temperature	0	100	оС
Storage temperature	-55	125	оС
Input voltage	-0.5	5.5	Volts
Output voltage ($V_{CC} = 3.1 - 3.6V$)	-0.5	$V_{CC} + 0.5$	Volts

Note: Stress above the conditions listed may cause permanent damage to the device. Functional operation of this device should be restricted to the conditions described under operating conditions.

DC CHARACTERISTICS

TA-0-70°C, V_{CC}=3.3V+/-5%, GND=0V

Symbol	Parameter	Min	Max	Unit	Condition
$V_{ m IL}$	Input low voltage	-0.50	0.8	V	
V _{IH}	Input high voltage	2.0	V _{CC} +0.5	V	
V _{OL}	Output low voltage	-	0.45	V	I _{OL} =4.0mA
V _{OH}	Output high voltage	2.4	-	V	I _{OH} =-1.0mA
${ m I}_{ m IL}$	Input leakage current	-	+/-10	uA	$0 < V_{IN} < V_{CC}$
I_{OZ}	Tristate leakage current	-	+/-20	uA	$0.45 < V_{OUT} < V_{CC}$
I_{CC}	Power supply current	-	TBD	mA	



AC CHARACTERISTICS

AC timing specifications provided are based on external zero-pf capacitance load. Min/Max cases are based on the following table:

Parameter	Min	Max	Unit
3.3V power (Vcc)	3.135	3.465	V
Temperature	0	95	оС

- SRAM interface Timing Characteristics

	SETUP	HOLD	MIN	MAX	UNIT	
SA output delay			2	7		
SD input	2.5	1.5				
SD output delay			2	7		
SADS			2	7		
SCS0			2	7		
SWE			2	7		

- RMII Interface Timing Characteristics

THE INVESTMENT THANKS CHARACTERS					
Parameter	min	type	max	unit	condition
RCLK50 cycle time			20	ns	
RXD CRS_DV setup time	4	-	-	ns	to RCLK50 rising
					edge
RXD CRS_DV hold time	2	-	-	ns	to RCLK50 rising
					edge
TXD TX_EN output	3	-	12	ns	to RCLK50 rising
delay					edge

- Management Interface (MI) Timing Characteristics

Parameter	min	typ	max	unit	condition
MDC cycle time	-	400	-	ns	
MDC high time	180	200	220	ns	
MDC low time	180	200	220	ns	
MDIO setup time (source by PHY)	30	-	-	ns	to MDC rising edge
MDIO hold time (source by PHY)	0	-	-	ns	to MDC rising edge



MDIO output	200	-	300	ns	to MDC rising
delay (source by					edge
vt3061)					

- EEPROM Interface Timing Characteristics

Parameter	min	typ	max	unit	condition
EEC clock		0	78.12	-	kHz
frequency					
Clock high time	-	6.4	-	μs	
Clock low time	-	6.4	-	μs	
Start Condition setup time		6.4	-	-	μs
Start Condition hold time	6.4	-	-	μs	
Stop Condition setup time		6.4	-	-	μs
Stop Condition hold time	6.4	-	-	μs	
Read Data In setup time	0	-	-		to EEC rising edge
Read Data In hold time	0	-	-		to EEC falling edge
EEIO Data out delay	2.6	-	3.0	μs	to EEC falling edge
Write Cycle time	-	11.4	-	ms	



PACKAGE MECHANICAL SPECIFICATIONS

