H8S/2655 Series

Hitachi Single-Chip Microcomputer

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Preface

Hitachi's H8S Series of single-chip microcomputers comprises new series which offer the high performance and low power consumption of the existing H8 Series, which is widely used for machine control, etc., together with significantly greater ease of use,

This initial series—the H8S/2655 Series—offers CPU object-level compatibility with the H8/300H Series, H8/300 Series, and H8/300L Series within the H8 Series.

Series	Features
H8S/2655	Upward-compatible with the H8/300H Series and H8/300 Series; twice the performance at the same frequency; multiply-and-accumulate instructions
H8/300H	16-Mbyte linear address space; upward-compatible with the H8/300 Series; concise instruction set; powerful word-size and longword-size arithmetic instructions
H8/300	64-kbyte address space; general register system; concise instruction set; powerful bit manipulation instructions
H8/300L	Same CPU as the H8/300 Series; consumer application oriented peripheral functions; low voltage, low power consumption

Intended Readership

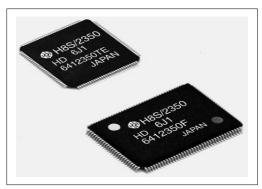
This Overview is intended for readers who require a basic understanding of microcomputers, or are looking for information on the features and functions of the H8S/2655 Series. Readers undertaking system design using these products, or requiring more detailed information on their use, should refer to the H8S/2655 Hardware Manual and H8S/2600 Series Programming Manual.

Related Documents

Contents	Document Title and No.
On H8S/2655 hardware	H8S/2655 Hardware Manual ADE-602-094
On H8S/2655 Series execution instructions	H8S/2600 Series, H8S/2000 Series Programming Manual ADE-602-083

Section 1 H8S/2655 Series Features

1.1 H8S/2655 Series Functions



H8S/2655 Series microcomputers are designed for faster instruction execution, using a realtime control oriented CPU with an internal 32bit architecture, and can run programs based on the C high-level language efficiently. As well as large-capacity ROM and RAM, these microcomputers include on-chip the peripheral functions needed for control systems. These features simplify the implementation of sophisticated, high-performance systems.

High-Performance H8S/2600 CPU

- General-register architecture
 - Sixteen 16-bit general registers (also usable as sixteen 8-bit registers or eight 32-bit registers)
- High-speed operation suitable for realtime control
 - 20 MHz maximum operating frequency (20 MHz oscillation frequency)
 - High-speed arithmetic operations

8/16/32-bit register-register add/subtract: 50 ns

 16×16 -bit register-register multiply: 200 ns

 $16 \times 16 + 42$ -bit multiply-and-accumulate: 200 ns

32 ÷ 16-bit register-register divide: 1000 ns

- Instruction set suitable for high-speed operation
 - Sixty-nine basic instructions
 - 8/16/32-bit move/arithmetic and logic instructions
 - Unsigned/signed multiply and divide instructions
 - Multiply-and-accumulate instruction
 - Powerful bit-manipulation instructions
- Two CPU operating modes
 - Normal mode: H8/300 Series compatible, maximum 64-kbyte address space
 - Advanced mode: Maximum 16-Mbyte address space

On-Chip Byte PROM (Mask ROM)

• 64 kbytes or 128 kbytes

On-Chip 4-kbyte High-Speed Static RAM

On-Chip Bus Controller

- Address space divided into 8 areas, with bus specifications settable independently for each area
- Chip select output possible for each area
- Selection of 8-bit or 16-bit access space for each area
- 2-state or 3-state access space can be designated for each area
- · Number of program wait states can be set for each area
- Burst ROM directly connectable
- Maximum 8-Mbyte DRAM or PSRAM directly connectable (or use of interval timer possible)
- External bus release function

DMA Controller (DMAC)

- · Selection of short address mode or full address mode
- Four channels in short address mode, two channels in full address mode
- Transfer possible in repeat mode, block transfer mode, etc.
- Single address mode transfer possible
- · Can be activated by internal interrupt

Data Transfer Controller (DTC)

- Activated by internal interrupt or software
- Multiple transfers or multiple types of transfer possible for one activation source
- Transfer possible in repeat mode, block transfer mode, etc.
- Request can be sent to CPU for interrupt that activated DTC

16-Bit Timer-Pulse Unit (TPU)

- Six-channel 16-bit timer on-chip
- Pulse I/O processing capability for up to 16 pins'
- Automatic 2-phase encoder count capability

Programmable Pulse Generator (PPG)

- Maximum 16-bit pulse output possible with TPU as time base
- Output trigger selectable in 4-bit groups
- Non-overlap margin can be set
- Direct output or inverse output setting possible

Two On-Chip 8-Bit Timer Channels

- 8-bit up-counter (external event count capability)
- Two time constant registers
- Two-channel connection possible

On-Chip Watchdog Timer (WDT)

• Watchdog timer or interval timer selectable

Three On-Chip Serial Communication Interface (SCI) Channels

- Asynchronous mode or synchronous mode selectable
- Multiprocessor communication function
- Smart card interface function

On-Chip A/D Converter

- Resolution: 10 bits
- Input: 8 channels
- High-speed conversion: 2.3 µs minimum conversion time (at 20 MHz operation)
- Select or group mode, and single or scan mode selectable
- · Sample and hold circuit
- A/D conversion can be activated by external trigger or timer trigger

On-Chip D/A Converter

Resolution: 8 bitsOutput: 2 channels

Thirteen I/O Ports

• 87 I/O pins, 8 input-only pins

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On-Chip Interrupt Controller

- Nine external interrupt pins (NMI, IRQ₀ to IRQ₇)
- 52 internal interrupt sources
- Selection of four interrupt control modes

Power-Down State

- Medium-speed mode
- Sleep mode
- Module stop mode
- Software standby mode
- Hardware standby mode

Seven MCU Operating Modes

	CPU			Extern	nal Data Bus
Mode	Operating Mode	Description	On-Chip ROM	Initial Value	Maximum Value
1	Normal	On-chip ROM disabled expansion mode	Disabled	8 bits	16 bits
2		On-chip ROM enabled expansion mode	Enabled	8 bits	16 bits
3	-	Single-chip mode	Enabled	_	_
4	Advanced	On-chip ROM disabled expansion mode	Disabled	16 bits	16 bits
5	-	On-chip ROM disabled expansion mode	Disabled	8 bits	16 bits
6	_	On-chip ROM enabled expansion mode	Enabled	8 bits	16 bits
7	-	Single-chip mode	Enabled	_	

On-Chip Clock Pulse Generator

• Built-in duty correction circuit

Packages

- 120-pin plastic TQFP (TFP-120)
- 128-pin plastic QFP (FP-128)

Product Lineup

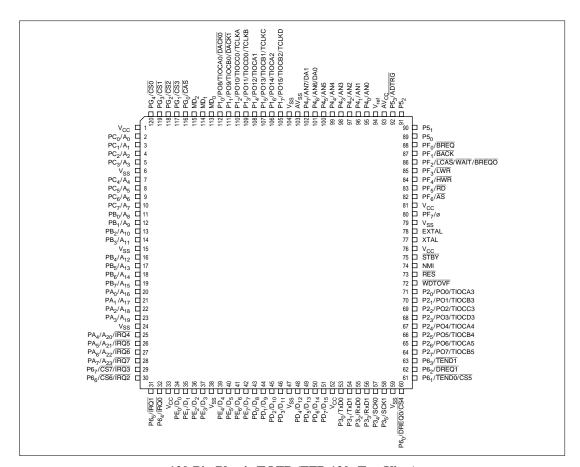
Model

Mask ROM Version	ZTAT™ Version	ROM/RAM (Bytes)	Packages
HD6432655	HD6472655	128 k/4 k	TFP-120 FP-128
HD6432653	_	64 k/4 k	TFP-120 FP-128

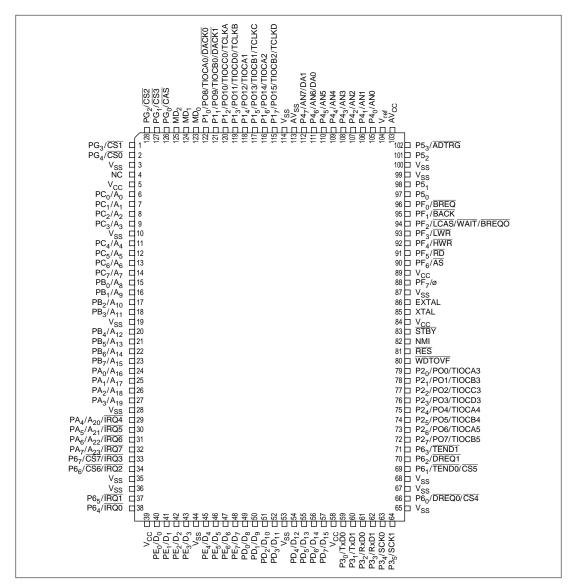
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1.2 Pin Description

Pin Arrangement



120-Pin Plastic TQFP (TFP-120: Top View)



128-Pin Plastic QFP (FP-128: Top View)

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Pin Functions

Туре	Symbol	I/O	Name a	nd Function	1		
Power	V _{CC}	Input	Power supply: All V_{CC} pins should be connected to the system power supply.			ld be connected to the	
	V _{SS}	Input	Ground: All V_{SS} pins should be connected to the system power supply (0 V).				
Clock	XTAL	Input	Pulse Ge	Connects to a crystal oscillator. See section 20, Clock Pulse Generator, for typical connection diagrams for a crystal oscillator and external clock input.			
	EXTAL	Input	input an Generate	Connects to a crystal oscillator. The EXTAL pin can also input an external clock. See section 20, Clock Pulse Generator, for typical connection diagrams for a crystal oscillator and external clock input.			
	Ø	Output	System clock: Supplies the system clock to an external device.				
Operating mode control	MD ₂ to MD ₀	Input	The relation	Mode pins: These pins set the operating mode. The relation between the settings of pins MD_2 to MD_0 and the operating mode is shown below. These pins should not be changed while the H8S/2655 Series is operating.			
			MD2	MD1	MD0	Operating Mode	
			0	0	0	_	
					1	Mode 1	
				1	0	Mode 2	
					1	Mode 3	
			1	0	0	Mode 4	
					1	Mode 5	
				1	0	Mode 6	
					1	Mode 7	
System control	RES	Input	Reset input: When this pin is driven low, the chip is reset.				
	STBY	Input	Standby: When this pin is driven low, a transition is made to hardware standby mode.				
	BREQ	Input	Bus request: Used by an external bus master to issue a bus request to the H8S/2655 Series.				
	BREQO	Output	Bus request output: The external bus request signal used when an internal bus master accesses external space in the external bus-released state.				

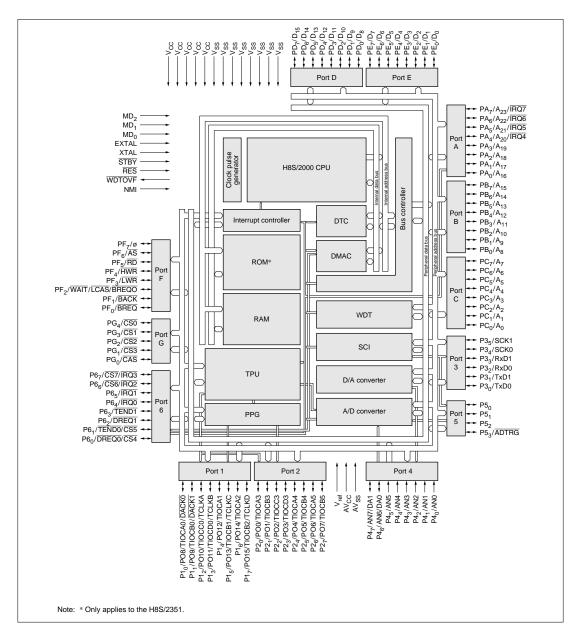
Туре	Symbol	I/O	Name and Function
	BACK	Output	Bus request acknowledge: Indicates that the bus has been released to an external bus master.
Interrupts	NMI	Input	Nonmaskable interrupt: Requests a nonmaskable interrupt.
	IRQ ₇ to IRQ ₀	Input	Interrupt request 7 to 0: These pins request a maskable interrupt.
Address bus	A ₂₃ to A ₀	Output	Address bus: These pins output an address.
Data bus	D ₁₅ to D ₀	I/O	Data bus: These pins constitute a bidirectional data bus.
Bus control	CS ₇ to CS ₀	Output	Chip select: Signals for selecting areas 7 to 0.
	AS	Output	Address strobe: When this pin is low, it indicates that address output on the address bus is enabled.
	RD	Output	Read: When this pin is low, it indicates that the external address space can be read.
	HWR	Output	High write/write enable/upper write enable: A strobe signal that writes to external space and indicates that the upper half (D ₁₅ to D ₈) of the data bus is enabled. The 2CAS type DRAM write enable signal. The 2WE type DRAM upper write enable signal.
	LWR	Output	Low write/lower column address strobe/lower write enable: A strobe signal that writes to external space and indicates that the lower half (D_7 to D_0) of the data bus is enabled. The 2CAS type DRAM column address strobe signal. The 2WE type DRAM lower write enable signal.
	CAS/OE	Output	Upper column address strobe/column address strobe/output enable/refresh: The 2CAS type DRAM upper column address strobe signal. The 2WE type DRAM column address strobe signal. The PSRAM output enable signal.
	WAIT Input		Wait: Requests insertion of a wait state in the bus cycle when accessing external 3-state address space.
DMA controller (DMAC)	DREQ ₁ , DREQ ₀ Input		DMA request 1 and 0: These pins request DMAC activation.
	TEND ₁ , TEND	Output	DMA transfer end 1 and 0: These pins indicate the end of DMAC data transfer.
	DACK ₁ , DACK ₀	Output	DMA transfer acknowledge 1 and 0: These are the DMAC single address transfer acknowledge pins.
16-bit timer- pulse unit (TPU	TCLKD to TCLKA	Input	Clock input D to A: These pins input an external clock.

Туре	Symbol	I/O	Name and Function
	TIOCA0, TIOCB0, TIOCC0, TIOCD0	I/O	Input capture/output compare match A0 to D0: The TGR0A to TGR0D input capture input or output compare output, or PWM output pins.
	TIOCA1, TIOCB1	I/O	Input capture/output compare match A1 and B1: The TGR1A and TGR1B input capture input or output compare output, or PWM output pins.
	TIOCA2, TIOCB2	I/O	Input capture/output compare match A2 and B2: The TGR2A and TGR2B input capture input or output compare output, or PWM output pins.
	TIOCA3, TIOCB3, TIOCC3, TIOCD3	I/O	Input capture/output compare match A3 to D3: The TGR3A to TGR3D input capture input or output compare output, or PWM output pins.
	TIOCA4, TIOCB4	I/O	Input capture/output compare match A4 and B4: The TGR4A and TGR4B input capture input or output compare output, or PWM output pins.
	TIOCA5, TIOCB5	I/O	Input capture/output compare match A5 and B5: The TGR5A and TGR5B input capture input or output compare output, or PWM output pins.
Programmable pulse generator (PPG)	PO ₁₅ to PO ₀	Output	Pulse output 15 to 0: Pulse output pins.
8-bit timer	TMO ₀ , TMO ₁	Output	Compare match output: The compare match output pins.
	TMCI ₀ , TMCI ₁	Input	Counter external clock input: Input pins for the external clock input to the counter.
	TMRI ₀ , TMRI ₁	Input	Counter external reset input: The counter reset input pins.
Watchdog timer (WDT)	WDTOVF	Output	Watchdog timer overflows: The counter overflows signal output pin in watchdog timer mode.
Serial communication	TxD_2 , TxD_1 , TxD_0	Output	Transmit data (channel 0, 1, 2): Data output pins.
interface (SCI) Smart Card interface	RxD_2 , RxD_1 , RxD_0	Input	Receive data (channel 0, 1, 2): Data input pins.
	SCK ₂ , SCK ₁ , SCK ₀	I/O	Serial clock (channel 0, 1, 2): Clock I/O pins.
A/D converter	AN ₇ to AN ₀	Input	Analog 7 to 0: Analog input pins.
	ADTRG	Input	A/D conversion external trigger input: Pin for input of an external trigger to start A/D conversion.
D/A converter	DA ₁ , DA ₀	Output	Analog output: D/A converter analog output pins.
·			

Туре	Symbol	I/O	Name and Function		
A/D converter and D/A converters	AV _{CC}	Input	This is the power supply pin for the A/D converter and D/A converter. When the A/D converter and D/A converter are not used, this pin should be connected to the system power supply (+5 V).		
	AV _{SS}	Input	This is the ground pin for the A/D converter and D/A converter. When the A/D converter and D/A converter are not used, this pin should be connected to the system power supply (0 V).		
	V _{ref}	Input	This is the reference voltage input pin for the A/D converter and D/A converter. When the A/D converter and D/A converter are not used, this pin should be connected to the system power supply (+5 V).		
I/O ports	P1 ₇ to P1 ₀	I/O	Port 1: An 8-bit I/O port. Input or output can be designated for each bit by means of the port 1 data direction register (P1DDR).		
	P2 ₇ to P2 ₀	I/O	Port 2: An 8-bit I/O port. Input or output can be designated for each bit by means of the port 2 data direction register (P2DDR).		
	P3 ₅ to P3 ₀	I/O	Port 3: A 6-bit I/O port. Input or output can be designated for each bit by means of the port 3 data direction register (P3DDR).		
	P4 ₇ to P4 ₀	Input	Port 4: An 8-bit input port.		
	P5 ₃ to P5 ₀	I/O	Port 5: A 4-bit I/O port. Input or output can be designated for each bit by means of the port 5 data direction register (P5DDR).		
	P6 ₇ to P6 ₀	I/O	Port 6: An 8-bit I/O port. Input or output can be designated for each bit by means of the port 6 data direction register (P6DDR).		
I/O ports	PA ₇ to PA ₀	I/O	Port A: An 8-bit I/O port. Input or output can be designated for each bit by means of the port A data direction register (PADDR).		
	PB ₇ to PB ₀	I/O	Port B: An 8-bit I/O port. Input or output can be designated for each bit by means of the port B data direction register (PBDDR).		
	PC ₇ to PC ₀	I/O	Port C: An 8-bit I/O port. Input or output can be designated for each bit by means of the port C data direction register (PCDDR).		
	PD ₇ to PD ₀	I/O	Port D: An 8-bit I/O port. Input or output can be designated for each bit by means of the port D data direction register (PDDDR).		
	_				

Type	Symbol	I/O	Name and Function
	PE ₇ to PE ₀	I/O	Port E: An 8-bit I/O port. Input or output can be designated for each bit by means of the port E data direction register (PEDDR).
	PF ₇ to PF ₀	I/O	Port F: An 8-bit I/O port. Input or output can be designated for each bit by means of the port F data direction register (PFDDR).
	PG ₄ to PG ₀	I/O	Port G: A 5-bit port. Input or output can be designated for each bit by means of the port G data direction register (PGDDR).

1.3 Block Diagram



Block Diagram

Section 2 CPU

2.1 Features

The H8S/2600 CPU is a high-speed central processing unit with an internal 32-bit architecture, and is upward compatible with the H8/300 and H8/300H CPUs.

The H8S/2600 CPU has sixteen 16-bit general registers, can address a 16-Mbyte (architecturally 4-Gbyte) linear access space, and is ideal for realtime control.

- Upward-compatible with H8/300 and H8/300H CPUs
 - Can execute H8/300 and H8/300H object programs
- General-register architecture
 - Sixteen 16-bit general registers (also usable as sixteen 8-bit registers or eight 32-bit registers)
- Sixty-nine basic instructions
 - 8/16/32-bit arithmetic and logic instructions
 - Multiply and divide instructions
 - Powerful bit-manipulation instructions
 - Multiply-and-accumulate instruction
- Eight addressing modes
 - Register direct (Rn)
 - Register indirect (@ERn)
 - Register indirect with displacement (@(d:16,ERn) or @(d:32,ERn))
 - Register indirect with post-increment or pre-decrement (@ERn+ or @-ERn)
 - Absolute address (@aa:8, @aa:16, @aa:24, or @aa:32)
 - Immediate (#xx:8, #xx:16, or #xx:32)
 - Program-counter relative (@(d:8,PC) or @(d:16,PC))
 - Memory indirect (@@aa:8)
- 16-Mbyte access space
 - Program: 16 Mbytes
 - Data: 16 Mbytes (architecturally 4 Gbytes)
- High-speed operation
 - All frequently-used instructions execute in one or two states
 - Maximum clock frequency: 20 MHz
 - 8/16-32-bit register-register add/subtract: 50 ns
 - 8 × 8-bit register-register multiply: 150 ns

- 16 ÷ 8-bit register-register divide: 600 ns
- 16×16 -bit register-register multiply: 200 ns
- 32 ÷ 16-bit register-register divide: 1000 ns
- Two CPU operating modes
 - Normal mode/advanced mode
- Low-power state
 - Transition to power-down state by SLEEP instruction
 - CPU clock speed selectable

Differences from H8/300 CPU

In comparison with the H8/300 CPU, the H8S/2600 CPU has the following enhancements.

- More general registers and control registers
 - Eight 16-bit registers, one 8-bit and two 32-bit control registers added
- · Expanded address space
 - Normal mode supports the same 64-kbyte address space as the H8/300 CPU
 - Advanced mode supports a maximum 16-Mbyte address space
- · Enhanced addressing
 - For effective use of the 16-Mbyte address space
- Enhanced instructions
 - Addressing modes of bit-manipulation instructions enhanced
 - Signed multiply and divide instructions added
 - Multiply-and-accumulate instruction added
 - Two-bit shift instructions added
 - Instructions for saving and restoring multiple registers added
 - Test-and-set instruction added
- Higher speed
 - Basic instructions execute twice as fast

Differences from H8/300H CPU

In comparison with the H8/300H CPU, the H8S/2600 CPU has the following enhancements.

- Additional control register
 - One 8-bit and two 32-bit control registers added
- Enhanced instructions

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- Addressing modes of bit-manipulation instructions enhanced
- Multiply-and-accumulate instruction added
- Two-bit shift instructions added
- Instructions for saving and restoring multiple registers added
- Test-and-set instruction added
- Higher speed
 - Basic instructions execute twice as fast

2.2 Register Configuration

The H8S/2600 CPU has general registers and control registers.

The eight 32-bit general registers all have identical functions and can be used as either address registers or data registers. The control registers are the 24-bit program counter (PC), 8-bit extend register (EXR), 8-bit condition code register (CCR), and 64-bit multiply-and-accumulate register (MAC).

CPU Internal Register Configuration

	15	0	7	0 7 0
ER0	E0		R0H	R0L
ER1	E1		R1H	R1L
ER2	E2		R2H	R2L
ER3	E3		R3H	R3L
ER4	E4		R4H	R4L
ER5	E5		R5H	R5L
ER6	E6		R6H	R6L
ER7 (SP)	E7		R7H	R7L
			CC	7 6 5 4 3 2 1 0
	63			41 32
MAC	(Sign ex			MACH
	31	IVIA	ACL	0
Legend SP: PC: EXR: T: I2 to I0: CCR: I: UI:	Stack pointer Program counter Extend register Trace bit Interrupt mask bits Condition code register Interrupt mask bit User bit/interrupt mask bit	H: U: N: Z: V: C: MAC:	Half-carry flag User bit Negative flag Zero flag Overflow flag Carry flag Multiply-and-accu	mulate register

General Registers

The CPU has eight 32-bit general registers. These general registers are all functionally alike and can be used as either address registers or data registers.

When a general register is used as a data register, it can be accessed as a 32-bit, 16-bit, or 8-bit register.

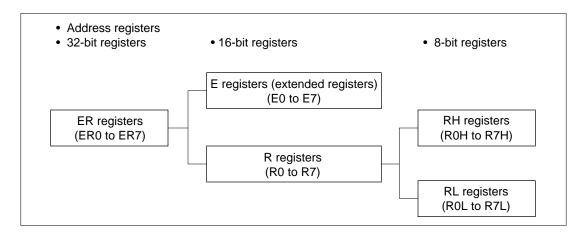
When the general registers are used as 32-bit registers or address registers, they are designated by the letters ER (ER0 to ER7).

The ER registers divide into 16-bit general registers designated by the letters E (E0 to E7) and R (R0 to R7). These registers are functionally equivalent, providing a maximum sixteen 16-bit registers. The E registers (E0 to E7) are also referred to as extended registers.

The R registers divide into 8-bit general registers designated by the letters RH (R0H to R7H) and RL (R0L to R7L). These registers are functionally equivalent, providing a maximum sixteen 8-bit registers.

The figure below illustrates the usage of the general registers. The usage of each register can be selected independently.

Usage of General Registers



Control Registers

The control registers are the 24-bit program counter (PC), 8-bit extend register (EXR), 8-bit condition code register (CCR), and 64-bit multiply-and-accumulate register (MAC).

Program Counter (PC): This 24-bit counter indicates the address of the next instruction the CPU will execute. The length of all CPU instructions is 2 bytes (one word) or a multiple of 2 bytes, so the least significant PC bit is ignored. When an instruction is fetched, the least significant PC bit is regarded as 0.

Extend Register (EXR): This 8-bit register comprises a trace bit (T) and interrupt mask bits (I2 to I0).

- Bit 7—Trace Bit (T)
 - Specifies whether or not trace mode is set. When this bit is cleared to 0, instructions are executed sequentially. When set to 1, trace exception handling is started each time an instruction is executed.
- Bits 6 to 3—Reserved
- Bits 2 to 0—Interrupt Mask Bits (I2 to I0)
 These bits specify the interrupt request mask level (0 to 7). See section 2.9, Interrupts, for details.

EXR can be manipulated by the LDC, STC, ANDC, ORC, and XORC instructions. Except in the case of STC, interrupts (including NMI) are not accepted for 3 states after the instruction is executed.

Condition Code Register (CCR): This 8-bit register contains internal CPU status information, including the interrupt mask bit (I), and the half-carry (H), negative (N), zero (Z), overflow (V), and carry (C) flags.

- Bit 7—Interrupt Mask Bit (I)
 - Masks interrupts other than NMI when set to 1. NMI is accepted regardless of the I bit setting. The I bit is set to 1 at the start of an exception-handling sequence. See section 2.9, Interrupts for details.
- Bit 6—User Bit or Interrupt Mask Bit (UI)
 Can be written or read by software using the LDC, STC, ANDC, ORC, and XORC instructions. This bit can also be used as an interrupt mask bit. See section 2.9, Interrupts, for
- Bit 5—Half-Carry Flag (H)

details.

- When the ADD.B, ADDX.B, SUB.B, SUBX.B, CMP.B, or NEG.B instruction is executed, this flag is set to 1 if there is a carry or borrow at bit 3, and cleared to 0 otherwise. When the ADD.W, SUB.W, CMP.W, or NEG.W instruction is executed, the H flag is set to 1 if there is a carry or borrow at bit 11, and cleared to 0 otherwise. When the ADD.L, SUB.L, CMP.L, or NEG.L instruction is executed, the H flag is set to 1 if there is a carry or borrow at bit 27, and cleared to 0 otherwise.
- Bit 4—User Bit (U)

Can be written and read by software using the LDC, STC, ANDC, ORC, and XORC instructions.

• Bit 3—Negative Flag (N)

Stores the value of the most significant bit (sign bit) of data.

• Bit 2—Zero Flag (Z)

Set to 1 to indicate zero data, and cleared to 0 to indicate non-zero data.

• Bit 1—Overflow Flag (V)

Set to 1 when an arithmetic overflow occurs, and cleared to 0 at other times.

• Bit 0—Carry Flag (C)

Set to 1 when a carry occurs, and cleared to 0 otherwise. Used by:

- Add instructions, to indicate a carry
- Subtract instructions, to indicate a borrow
- Shift and rotate instructions, to store the value shifted out of the end bit

The carry flag is also used as a bit accumulator by bit-manipulation instructions.

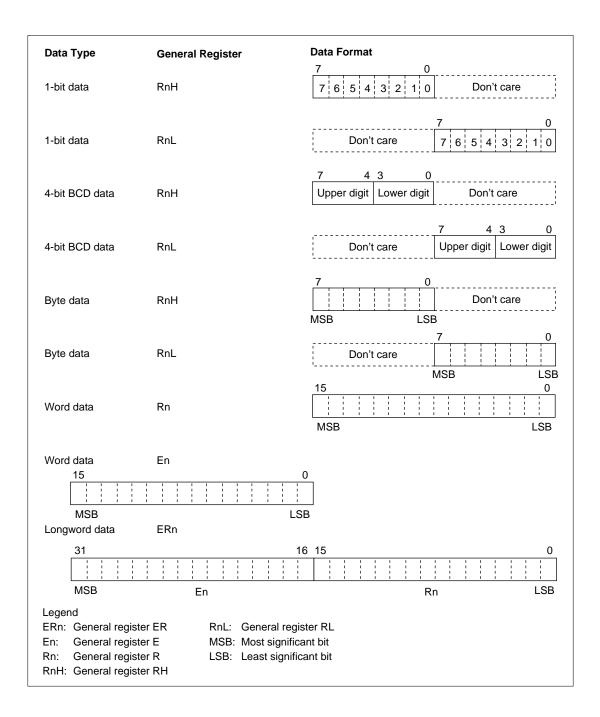
2.3 Data Formats

The CPU can process 1-bit, 4-bit (BCD), 8-bit (byte), 16-bit (word), and 32-bit (longword) data.

Bit-manipulation instructions operate on 1-bit data by accessing bit n (n = 0, 1, 2, ..., 7) of byte operand data.

The DAA and DAS decimal-adjust instructions treat byte data as two digits of 4-bit BCD data.

General Register Data Formats



Memory Data Formats

Data Type	Data Format
	Address
1-bit data	7 0 Address L 7 6 5 4 3 2 1 0
Byte data	Address L MSB LSB
Word data	Address 2M MSB LSB
Longword data	Address 2N MSB Address 2N + 1 Address 2N + 2 Address 2N + 3 LSB

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2.4 Addressing Modes

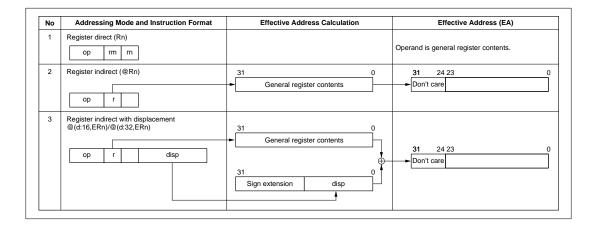
The H8S/2600 CPU supports eight addressing modes.

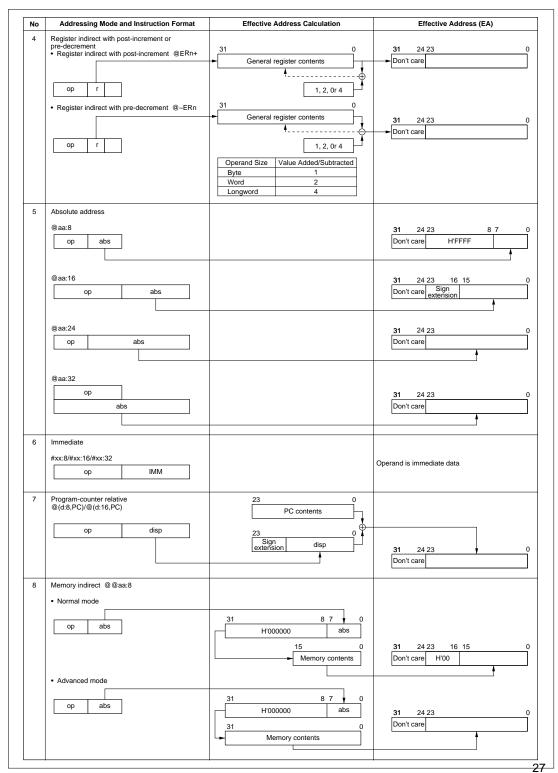
Addressing Modes

No.	Addressing Mode	Symbol
1	Register direct	Rn
2	Register indirect	@ERn
3	Register indirect with displacement	@(d:16,ERn)/@(d:32,ERn)
4	Register indirect with post-increment	@Ern+
	Register indirect with pre-decrement	@-ERn
5	Absolute address	@aa:8/@aa:16/@aa:24/@aa:32
6	Immediate	#xx:8/#xx:16/#xx:32
7	Program-counter relative	@(d:8,PC)/@(d:16,PC)
8	Memory indirect	@ @aa:8

Effective Address (EA) Calculation

In normal mode, the upper 8 bits of the effective address are ignored in order to generate a 16-bit effective address.





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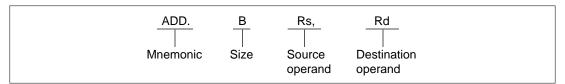
2.5 Instruction Set

The H8S/2600 CPU has 69 types of instructions.

Features

- Upward-compatible at object level with H8/300H and H8/300 CPUs.
- General register architecture
- 8/16/32-bit transfer instructions and arithmetic and logic instructions
 - Byte (B), word (W), and longword (L) formats for transfer instructions and basic arithmetic and logic instructions
- Unsigned and signed multiply and divide instructions
- Multiply-and-accumulate instruction
- Powerful bit-manipulation instructions
- Instructions for saving and restoring multiple registers

The ADD instruction format is shown below as an example.



Assembler Format

Instruction Set Table

1. Data transfer instructions

			Мс	ode	A e/In	str	uct	sir ior	١Ľ	enç	gth		Co	onc	litic	on	Со	de	No. Sta *1	of tes
Mner	monic	Operand Size	xx#	Rn	@ERn	@(d,ERn)	@-ERn/@ERn+	@aa	@(d,PC)	@ @aa	I	Operation	ı	н	N	Z	٧	С	Normal	Advanced
MOV	MOV.B #xx:8, Rd	В	2									#xx:8ÆRd8	_	_	\(\)	\(\)	0	_	1	
	MOV.B Rs,Rd	В		2								Rs8ÆRd8	_	_	\	\	0	_	1	
	MOV.B @ERs, Rd	В			2							@ERsÆRd8	_	_	\	\	0	_	2	2
	MOV.B @(d:16 ,ERs),R d	В				4						@(d:16,ERs)ÆRd8			\Q	\Q	0		3	}
	MOV.B @(d:32 ,ERs),R d	В				8						@(d:32,ERs)ÆRd8	_		\rightarrow	\rightarrow	0		5	,
	MOV.B @ERs+ ,Rd	В					2					@ERsÆRd8,ERs32+1Æ ERs32	_	_	\(\)	\(\)	0	_	3	}
	MOV.B @aa:8, Rd	В						2				@aa:8ÆRd8	_	_	\(\)	\(\)	0	_	2	<u>)</u>
	MOV.B @aa:16 ,Rd	В						4				@aa:16ÆRd8	_	_	◊	◊	0	_	3	3
	MOV.B @aa:32 ,Rd	В						6				@aa:32ÆRd8	_	_	\(\)	\(\)	0	_	2	ļ
	MOV.B Rs,@E Rd	В			2							Rs8Æ@ERd	_	_	\rightarrow	\rightarrow	0	_	2)
	MOV.B Rs,@(d :16,ER d)	В				4						Rd8Æ@(d:16,ERd)			\rightarrow	\rightarrow	0		3	}
	MOV.B Rs,@(d :32,ER d)					8						Rd8Æ@(d:32,ERd)			\Q	◊	0		5	5

			Мо	ode	A e/In	dd str (B	res uct	ior	ng n L	enç	gth		Co	ond	litic	on	Со	de	No. Sta *1	. of ites
Mner	monic	Operand Size	xx#	Rn	@ERn	@(d,ERn)	@-ERn/@ERn+	Фаа	@(d,PC)	@ @aa	ı	Operation	ı	Н	N	Z	v	С	Normal	Advanced
	MOV.B Rs,@- ERd	В					2					ERd32- 1ÆERd32,Rs8Æ@ERd	_		\rightarrow	\rightarrow	0			3
	MOV.B Rs,@a a:8	В						2				Rs8Æ@aa:8			<	<	0		12	2
	MOV.B Rs,@a a:16	В						4				Rs8Æ@aa:16	_		\Q	\Q	0	_	``	3
	MOV.B Rs,@a a:32	В						6				Rs8Æ@aa:32			<	<	0		2	4
	MOV.W #xx:16, Rd	W	4									#xx:16ÆRd16	_		\Q	\Q	0	_	2	2
	MOV.W Rs,Rd	W		2								Rs16ÆRd16			\rightarrow	\rightarrow	0		•	1
	MOV.W @ERs, Rd	W			2							@ERsÆRd16	_		\(\)	\(\)	0	_	2	2
	MOV.W @(d:16 ,ERs),R d					4						@(d:16,ERs)ÆRd16			\rightarrow	\rightarrow	0		(3
	MOV.W @(d:32 ,ERs),R d					8						@(d:32,ERs)ÆRd16			\rightarrow	\rightarrow	0		ļ	5
	MOV.W @ERs+ ,Rd						2					ERsÆRd16,ERs32+2Æ @ERd32	_		\(\)	\(\)	0	_	(3
	MOV.W @aa:16 ,Rd							4				@aa:16ÆRd16	_	_	◊	◊	0	_	(3
	MOV.W @aa:32 ,Rd							6				@aa:32ÆRd16	_	_	\(\)	◊	0	_		4
	MOV.W Rs,@E Rd	W			2							Rs16Æ@ERd	_		◊	◊	0		2	2
	MOV.W Rs,@(d :16,ER d)					4						Rs16Æ@(d:16,ERd)			\Q	◊	0		,	3

			Mo	ode	A e/In	str	Syte	tior	ıĽ	enç	gth		Co	onc	litic	on	Со	de	No. Sta *1	
Mner	monic	Operand Size	xx#	Rn	@ERn	@(d,ERn)	@-ERn/@ERn+	@aa	@(d,PC)	@ @aa	I	Operation	ı	н	N	z	V	С	Normal	Advanced
	MOV.W Rs,@(d :32,ER d)					8						Rs16Æ@(d:32,ERd)			\Q	◊	0		5	5
	MOV.W Rs,@- ERd	W					2					ERd32– 2ÆERd32,Rs16Æ@ERd	_		◊	◊	0		3	8
	MOV.W Rs,@a a:16	W						4				Rs16Æ@aa:16	_	_	◊	◊	0		3	3
	MOV.W Rs,@a a:32	W						6				Rs16Æ@aa:32	_	_	◊	◊	0		2	1
	MOV.L #xx:32, Rd	L	6									#xx:32ÆRd32	_	_	◊	◊	0		3	3
	MOV.L ERs,E Rd	L		2								ERs32ÆERd32	_	_	◊	◊	0		1	l
	MOV.L @ERs, ERd	L			4							@ERsÆERd32	_		◊	◊	0		4	1
	MOV.L @(d:16 ,ERs),E Rd	L				6						@(d:16,ERs)ÆERd32	_		◊	◊	0		5	5
	MOV.L @(d:32 ,ERs),E Rd	L				10						@(d:32,ERs)ÆERd32			◊	◊	0		7	7
	MOV.L @ERs+ ,ERd	L					4					@ERsÆERd32,ERs32+4 Æ @ERs32			◊	◊	0		5	5
	MOV.L @aa:16 ,ERd							6				@aa:16ÆERd32			◊	◊	0		5	5
	MOV.L @aa:32 ,ERd	L						8				@aa:32ÆERd32	_	_	◊	◊	0		6	6

			Mo	ode	A e/In	dd str (B	res uct	ior	ig i Lo	enç	jth		Co	ond	litic	on	Co	de	No. Sta *1	
Mner	monic	Operand Size	xx#	Rn	@ERn	@(d,ERn)	@-ERn/@ERn+	@aa	@(q,PC)	@ @ aa	-	Operation	-	н	N	Z	>	C	Normal	Advanced
MOV	MOV.L ERs,@ ERd	L			4							ERs32Æ@ERd		_	\Diamond	\Diamond	0		4	ļ
	MOV.L ERs,@(d:16,E Rd)	L				6						ERs32Æ@(d:16,ERd)		_	<	<	0		5	5
	MOV.L ERs,@(d:32,E Rd)	L				10						ERs32Æ@(d:32,ERd)			\diamond	\diamond	0		7	7
	MOV.L ERs,@ -ERd	L					4					ERd32- 4ÆERd32,ERs32Æ@ER d	_	_	\Diamond	\Diamond	0		5	5
	MOV.L ERs,@ aa:16	L						6				ERs32Æ@aa:16		_	\rightarrow	\rightarrow	0		5	
	MOV.L ERs,@ aa:32	L						8				ERs32Æ@aa:32		_	\rightarrow	\rightarrow	0		6	÷
POP	POP.W Rn	W									2	@SPÆRn16,SP+2ÆSP	_		\	\	0		3	3
	POP.L ERn	L									4	@SPÆERn32,SP+4ÆS P	_		\	\	0		5	5
PUSH	PUSH. W Rn	W									2	SP-2ÆSP,Rn16Æ@SP		_	\Diamond	\Diamond	0		3	3
	PUSH. L ERn	L									4	SP- 4ÆSP,ERn32Æ@SP	_	_	\Diamond	\Diamond	0		5	5
LDM	LDM @SP+,(ERm- ERn)	L									4	(@SPÆERn32,SP+4ÆS P) Repeated for each register restored							7/9, [1	
STM	STM (ERm- ERn), @-SP	L									4	(SP– 4ÆSP,ERn32Æ@SP) Repeated for each register saved							7/9. [1	
MOVFP E	MOVFP E @aa:16 ,Rd		ann	ot l	be	use	ed i	n th	ne H	185	6/26	655 Series							[2	2]
MOVTP E	MOVTP E Rs,@a a:16																		[2	2]

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2. Arithmetic instructions

			A	ddre				le/In 3yte		ıctio	on			Con	ditio	on C	ode	е	No. Stat	of es ^{*1}
Mner	monic	Operand Size	xx#	Rn	@ERn	@(d,ERn)	@-ERn/@ERn+	@aa	@(d,PC)	@ @ aa	1	Operation	ı	Н	Ν	Z	٧	С	Normal	Advanced
ADD	ADD.B #xx:8,R d	В	2									Rd8+#xx:8ÆRd8	_	\Diamond	\rightarrow	\Diamond	\rightarrow	\Q	,	1
	ADD.B Rs,Rd	В		2								Rd8+Rs8ÆRd8	_	\Diamond	\rightarrow	\Diamond	\rightarrow	\	,	1
	ADD.W #xx:16, Rd	W	4									Rd16+#xx:16ÆRd16	_	[3]	\rightarrow	\Diamond	\rightarrow	\Q	2	2
	ADD.W Rs,Rd	8		2								Rd16+Rs16ÆRd16	_	[3]	\Diamond	\Diamond	\Diamond	\Diamond		1
	ADD.L #xx:32, ERd	L	6									ERd32+#xx:32ÆERd32	_	[4]	\diamond	\diamond	\Diamond	\Diamond	•	3
	ADD.L ERs,E Rd	L		2								ERd32+ERs32ÆERd32	_	[4]	\Diamond	\Diamond	\Diamond	\rightarrow		1
ADDX	ADDX #xx:8,R d	В	2									Rd8+#xx:8+CÆRd8	_	\Diamond	\rightarrow	[5]	\rightarrow	\Q	,	1
	ADDX Rs,Rd	В		2								Rd8+Rs8+CÆRd8	_	\Diamond	\Diamond	[5]	\Diamond	\	,	1
ADDS	ADDS #1,ERd	L		2								ERd32+1ÆERd32	_						•	1
	ADDS #2,ERd	L		2								ERd32+2ÆERd32	_						•	1
	ADDS #4,ERd	L		2								ERd32+4ÆERd32							•	1
INC	INC.B Rd	В		2								Rd8+1ÆRd8	_		\(\)	\Diamond	\Diamond	_		1
	INC.W #1,Rd	W		2								Rd16+1ÆRd16	_	_	\	\	\rightarrow	_		1
	INC.W #2,Rd	W		2								Rd16+2ÆRd16	_		\(\)	\Diamond	\Diamond	_		1
	INC.L #1,ERd	L		2								ERd32+1ÆERd32			\rightarrow	\Diamond	\Diamond			1
	INC.L #2,ERd	L		2								ERd32+2ÆERd32			\rightarrow	\Diamond	\Diamond			1
DAA	DAA Rd	В		2								Rd8 decimal adjust Æ Rd8	_	*	\Diamond	\Diamond	*	\Diamond		1
SUB	SUB.B Rs,Rd	В		2								Rd8-Rs8ÆRd8	_	\Diamond	\Diamond	\Diamond	\Diamond	\	•	1

			A	ddre			th (I	le/In Byte		uctio	on			Con	diti	on (Code	е	No. Stat	of es ^{*1}
Mner	monic	Operand Size	xx#	Rn	@ERn	@(d,ERn)	@-ERn/@ERn+	@aa	@(d,PC)	@@aa	I	Operation	I	н	N	Z	٧	С	Normal	Advanced
	SUB.W #xx:16, Rd	W	4									Rd16-#xx:16ÆRd16		[3]	\Q	◊	\rightarrow	\rightarrow	2	2
	SUB.W Rs,Rd	W		2								Rd16-Rs16ÆRd16		[3]	\Diamond	◊	\	\	1	ı
	SUB.L #xx:32, ERd	L	6									ERd32-#xx:32ÆERd32		[4]	\Q	◊	\Q	\Q	(1)	3
	SUB.L ERs,E Rd	L		2								ERd32-ERs32ÆERd32		[4]	\Q	◊	<	<	1	l
SUBX	SUBX #xx:8,R d	В	2									Rd8-#xx:8-CÆRd8		◊	\rightarrow	[5]	\rightarrow	\rightarrow	1	I
	SUBX Rs,Rd	В		2								Rd8-Rs8-CÆRd8	_	◊	\	[5]	\	\	1	
SUBS	SUBS #1,ERd	L		2								ERd32-1ÆERd32							1	
	SUBS #2,ERd	L		2								ERd32-2ÆERd32							1	
	SUBS #4,ERd	L		2								ERd32-4ÆERd32							1	
DEC	DEC.B Rd	В		2								Rd8-1ÆRd8			\Diamond	\Q	\Diamond	_	1	ı
	DEC.W #1,Rd	8		2								Rd16-1ÆRd16			\Diamond	\Q	\Diamond	_	1	ı
	DEC.W #2,Rd	W		2								Rd16-2ÆRd16			\Diamond	◊	\Diamond		1	
	DEC.L #1,ERd	L		2								ERd32-1ÆERd32	_	_	\Diamond	\	\Diamond	_	1	
	DEC.L #2,ERd	L		2								ERd32-2ÆERd32	_	_	\Diamond	\	\Diamond	_	1	
DAS	DAS Rd	В		2								Rd8 decimal adjust Æ Rd8		*	\Diamond	\Q	*		1	
MULXU	.B Rs,Rd			2								Rd8¥Rs8ÆRd16 (unsigned multiplication)					_		(7)	3
	MULXU .W Rs,ERd			2								Rd16¥Rs16ÆERd32 (unsigned multiplication)							4	1
MULXS	MULXS .B Rs,Rd	В		4								Rd8¥Rs8ÆRd16 (signed multiplication)			\Q	◊	_	_	2	1

			A	ddre		_	Mod th (E			ıctio	on			Con	ditio	on C	Code	e	No. Stat	- 4
Mne	monic	Operand Size	xx#	Rn	@ERn	@(d,ERn)	@-ERn/@ERn+	@aa	@(d,PC)	@ @aa	1	Operation	ı	Н	N	Z	٧	С	Normal	Advanced
	MULXS .W Rs,ERd			4								Rd16¥Rs16ÆERd32 (signed multiplication)			\rightarrow	\Q			ţ	5

			A	ddre		_	th (E	e/In Byte		ıctio	on			Con	diti	on C	Code	9	No. of States ^{*1}
Mner	monic	Operand Size	xx#	Rn	@ERn	@(d,ERn)	@-ERn/@ERn+	gaa	@(q,PC)	ee @ @	1	Operation	-	н	N	Z	>	C	Normal Advanced
DIVXU	DIVXU. B Rs,Rd	В		2								Rd16÷Rs8ÆRd16 (RdH: remainder, RdL: quotient) (unsigned division)			[6]	[7]			12
	DIVXU. W Rs,ERd			2								ERd32÷Rs16ÆERd32 (Ed: remainder, Rd: quotient) (unsigned division)			[6]	[7]			20
DIVXS	divxs.B Rs,Rd	В		4								Rd16÷Rs8ÆRd16 (RdH: remainder, RdL: quotient) (signed division)		_	[8]	[7]		_	13
	DIVXS. W Rs,ERd			4								ERd32÷Rs16ÆERd32 (Ed: remainder, Rd: quotient) (signed division)	-	_	[8]	[7]			21
CMP	CMP.B #xx:8,R d	В	2									Rd8-#xx:8		\	◊	\rightarrow	\Diamond	◊	1
	CMP.B Rs,Rd	В		2								Rd8-Rs8		\	◊	◊	\Diamond	\	1
	CMP.W #xx:16, Rd	W	4									Rd16-#xx:16	_	[3]	◊	\rightarrow	\rightarrow	\(\)	2
	CMP.W Rs,Rd	W		2								Rd16-Rs16		[3]	◊	◊	\Diamond	\	1
	CMP.L #xx:32, ERd	L	6									ERd32-#xx:32	_	[4]	◊	◊	\diamond	\rightarrow	3
	CMP.L ERs,E Rd	L		2								ERd32-ERs32	_	[4]	◊	\rightarrow	\rightarrow	\(\)	1
NEG	NEG.B Rd	В		2								0-Rd8ÆRd8	_	\	◊	◊	\Diamond	\	1
	NEG.W Rd	W		2								0-Rd16ÆRd16		\	◊	◊	\rightarrow	\	1
	NEG.L ERd	L		2								0-ERd32ÆERd32		\Q	◊	◊	\Diamond	\	1
EXTU	EXTU. W Rd	W		2								0 Æ (<bits 15="" 8="" to=""> of Rd16)</bits>	_	_	0	\Q	0	_	1
	EXTU.L ERd	L		2								0 Æ (<bits 16="" 31="" to=""> of ERd32)</bits>	_	_	0	◊	0	_	1
EXTS	EXTS. W Rd	W		2								(<bit 7=""> of Rd16) Æ (<bits 15="" 8="" to=""> of Rd16)</bits></bit>		_	\rightarrow	◊	0	_	1
	EXTS.L ERd	L		2								(<bit 15=""> of ERd32) Æ (<bits 16="" 31="" to=""> of ERd32)</bits></bit>		_	◊	\rightarrow	0		1

			A	ddre		ng I engi				ıctio	on		,	Con	ditie	on (Code	е	No. Stat	of es ^{*1}
Mner	monic	Operand Size	xx#	Rn	@ERn	@(d,ERn)	@-ERn/@ERn+	@aa	@(d,PC)	@ @ aa	1	Operation	ı	Н	N	Z	٧	С	Normal	Advanced
TAS	TAS @ERd	В			4							@ERd-0 Æ CRR set, (1) Æ (<bit 7=""> of @ERd)</bit>			\(\)	◊	0		2	1
MAC	MAC @ERn+ ,@ERm +						4					@ERn¥@ERm+MACÆM AC (signed multiplication) ERn+2ÆERn,ERm+2ÆE Rm			 [9]	 [9]	 [9]		2	1
CLRMA C	CLRMA C										2	0ÆMACH,MACL	_	_		_	_		2	2
LDMAC	LDMAC ERs,M ACH	L		2								ERsÆMACH					_		2	2
	LDMAC ERs,M ACL	L		2								ERsÆMACL	_				_		2	2
STMAC	STMAC MACH, ERd	L		2								MACHÆERd			\(\)	\Q	\(\)		1	1
	STMAC MACL, ERd	L		2								MACLÆERd			\Q	\Q	◊		1	1

3. Logical instructions

			A	ddre		_	th (E	le/In Byte		uctio	on		(Con	diti	on (Code	е	No.	of es ^{*1}
Mner	monic	Operand Size	xx#	Rn	@ERn	@(d,ERn)	@-ERn/@ERn+	@aa	@(d,PC)	@ @ aa	I	Operation	ı	Н	N	z	v	С	Normal	Advanced
AND	AND.B #xx:8,R d	В	2									Rd8Ÿ#xx:8ÆRd8	_	_	◊	◊	0	_	1	
	AND.B Rs,Rd	В		2								Rd8ŸRs8ÆRd8	_	_	◊	◊	0	_	1	l
	AND.W #xx:16, Rd	W	4									Rd16Ÿ#xx:16ÆRd16	_	_	◊	◊	0	_	2	<u>?</u>
	AND.W Rs,Rd	W		2								Rd16ŸRs16ÆRd16	_	_	◊	◊	0	_	1	ĺ
	AND.L #xx:32, ERd	L	6									ERd32Ÿ#xx:32ÆERd32		_	◊	◊	0	_	3	}
	AND.L ERs,E Rd	L		4								ERd32ŸERs32ÆERd32	_	_	◊	◊	0	_	2	<u>,</u>
OR	OR.B #xx:8,R d	В	2									Rd8/#xx:8ÆRd8	_	_	◊	◊	0	_	1	
	OR.B Rs,Rd	В		2								Rd8/Rs8ÆRd8	_	_	◊	◊	0	_	1	ĺ
	OR.W #xx:16, Rd	W	4									Rd16/#xx:16ÆRd16		_	◊	◊	0	_	2	2
	OR.W Rs,Rd	W		2								Rd16/Rs16ÆRd16	_	_	◊	◊	0	_	1	ĺ
	OR.L #xx:32, ERd	L	6									ERd32/#xx:32ÆERd32	_		◊	◊	0	_	3	}
	OR.L ERs,E Rd	L		4								ERd32/ERs32/EERd32	_		◊	◊	0		2	<u>}</u>
XOR	XOR.B #xx:8,R d	В	2									Rd8≈#xx:8ÆRd8	_		◊	◊	0	_	1	
	XOR.B Rs,Rd	В		2								Rd8≈Rs8ÆRd8			\Q	◊	0		1	
	XOR.W #xx:16, Rd	W	4									Rd16≈#xx:16ÆRd16	_	_	◊	\rightarrow	0	_	2	2
	XOR.W Rs,Rd	W		2								Rd16≈Rs16ÆRd16	_	_	◊	\Diamond	0		1	

			A	ddre		ng l engt				ıctio	on			Con	ditio	on C	ode	е	No.	
Mner	monic	Operand Size	xx#	Rn	@ERn	@(d,ERn)	@-ERn/@ERn+	@aa	@(d,PC)	@ @aa	I	Operation	I	Н	N	Z	>	C	Normal	Advanced
	XOR.L #xx:32, ERd	L	6									ERd32≈#xx:32ÆERd32			\Diamond	\rightarrow	0		3	3
	XOR.L ERs,E Rd	L		4								ERd32≈ERs32ÆERd32			\Diamond	◊	0	_	2	2
NOT	NOT.B Rd	В		2								¬Rd8ÆRd8	_		\	\	0	_	1	1
	NOT.W Rd	W		2								¬Rd16ÆRd16			\rightarrow	\	0	_	1	1
	NOT.L ERd	L		2								¬Rd32ÆRd32			\rightarrow	\Q	0		1	1

4. Software instructions

			Ad	dre		_		de/l Byt		uct	ion					onc				No. of State s*1
	Mnemonic	Operand Size	xx#	Rn	@ERn	@(d,ERn)	@-ERn/@ERn+	@aa	@(d,PC)	@ @ aa	ı		Operation	ı	Н	N	z	٧	С	Normal Advanced
SHAL	SHAL.B Rd	В		2									-0			۰0	\Diamond	\Diamond	\Diamond	1
	SHAL.B #2,Rd	В		2								С	S MSB 			· (\)	\Diamond	\Diamond	\Diamond	1
	SHAL.W Rd	8		2												\diamond	\Diamond	\	\Diamond	1
	SHAL.W #2,Rd	8		2												\diamond	\Diamond	\	\Diamond	1
	SHAL.L ERd	L		2												\diamond	\Diamond	\	\Diamond	1
	SHAL.L #2,ERd	L		2												۰0	\Diamond	_	_	1
SHAR	SHAR.B Rd	В		2									-			۰0	\Diamond	0	\Diamond	1
	SHAR.B #2,Rd	В		2									MSB → LSB C			۰0	\Diamond	0	\Diamond	1
	SHAR.W Rd	W		2												۰0	\Diamond	0	\Diamond	1
	SHAR.W #2,Rd	W		2												۰0	\Diamond	0	\Diamond	1
	SHAR.L ERd	L		2												\	\Diamond	0	\Diamond	1
	SHAR.L #2,ERd	L		2												-◊	\Diamond	0	\Diamond	1
SHLL	SHLL.B Rd	В		2								إ	-0			۰0	♦	F		1
	SHLL.B #2,Rd	В		2								С	C MSB ← LSB			۰0	\Diamond	0	\Diamond	1
	SHLL.W Rd	W		2												۰0	\Diamond	0	\Diamond	1
	SHLL.W #2,Rd	W		2												۰0	\Diamond	0	\Diamond	1
	SHLL.L ERd	L		2												\	\Diamond	0	\Diamond	1
	SHLL.L #2,ERd	L		2									•			· (\)	\Diamond	_	\Diamond	1
SHLR	SHLR.B Rd	В		2								0	-			0	♦	0	\Diamond	1
	SHLR.B #2,Rd	В		2									MSB — ► LSB C			0	\	0	\Diamond	1
	SHLR.W Rd	W		2												0	\Diamond	0	\Diamond	1
	SHLR.W #2,Rd	W		2												0	\Diamond	0	\Diamond	1
	SHLR.L ERd	L		2												0	\Diamond	0	\Diamond	1
	SHLR.L #2,ERd	L		2											Е	0	\Diamond	0	\Diamond	1
ROTXL	ROTXL.B Rd	В		2													_		\	1
	ROTXL.B #2,Rd	В		2									C MSB ← LSB			_	_	_	\Diamond	1
	ROTXL.W Rd	W		2									<u> </u>			-◊	\Diamond	0	\Diamond	1
	ROTXL.W #2,Rd	W		2									[-◊	\	0	\Diamond	1
	ROTXL.L ERd	L		2									[.◊	\Diamond	0	\Diamond	1
	ROTXL.L #2,ERd	L		2									•			-◊	\Diamond	0	\Diamond	1

			Ad	dre				de/li Byte		uct	ion					ditio			No. of State s*1
	Mnemonic	Operand Size	#xx	Rn	@ERn	@(d,ERn)	@-ERn/@ERn+	@aa	@(d,PC)	@ @ aa	ı	Operation	I	Н	N	z	V	С	Normal Advanced
ROTXR	ROTXR.B Rd	В		2										H	۰0	\	0	\Diamond	1
	ROTXR.B #2,Rd	В		2								MSB → LSB C		Н	· (>	\Diamond	0	\Diamond	1
	ROTXR.W Rd	W		2											· \	\Diamond	0	\Diamond	1
	ROTXR.W #2,Rd	W		2											· \	\Diamond	0	\Diamond	1
	ROTXR.L ERd	L		2											\	\Diamond	0	\Diamond	1
	ROTXR.L #2,ERd	L		2											· (>	\	0	\Diamond	1
ROTL	ROTL.B Rd	В		2											\Diamond	\Diamond	0	\Diamond	1
	ROTL.B #2,Rd	В		2								C MSB ← LSB			\Diamond	\Diamond	0	\Diamond	1
	ROTL.W Rd	W		2											\Diamond	\Diamond	0	\Diamond	1
	ROTL.W #2,Rd	W		2											\Diamond	\Diamond	0	\Diamond	1
	ROTL.L ERd	L		2											\Diamond	\Diamond	0	\Diamond	1
	ROTL.L #2,ERd	L		2											\Diamond	\Diamond	0	\Diamond	1
ROTR	ROTR.B Rd	В		2											\Diamond	\Diamond	0	\Diamond	1
	ROTR.B #2,Rd	В		2								MSB → LSB C			\Diamond	\Diamond	0	\Diamond	1
	ROTR.W Rd	8		2											\Diamond	\Diamond	0	\Diamond	1
	ROTR.W #2,Rd	8		2											\Diamond	\Diamond	0	\Diamond	1
	ROTR.L ERd	L		2											\Diamond	\Diamond	0	\Diamond	1
	ROTR.L #2,ERd	L		2									\vdash		\Diamond	\Diamond	0	\Diamond	1

5. Bit manipulation instructions

			Α	ddre		_	th (E	le/In Byte		uctio	on			Con	diti	on (Code	е	No. Stat	of es ^{*1}
м	Inemonic	Operand Size	*xx#	Rn	@ERn	@(d,ERn)	@-ERn/@ERn+	@aa	@(q,PC)	@ @ aa	I	Operation	ı	H	N	z	>	С	Normal	Advanced
BSET	BSET #xx:3,Rd	В		2								(#xx:3 of Rd8)"1	_			_			1	1
	BSET #xx:3,@ERd	В			4							(#xx:3 of @ERd)"1	_		_	_			2	ţ
	BSET #xx:3,@aa:8	В						4				(#xx:3 of @aa:8)"1	_		_				4	ţ
	BSET #xx:3,@aa:16	В						6				(#xx:3 of @aa:16)"1	_						5	5
	BSET #xx:3,@aa:32	В						8				(#xx:3 of @aa:32)"1	_		_	_			6	3
	BSET Rn,Rd	В		2								(Rn8 of Rd8)"1	_			_			1	1
	BSET Rn,@ERd	В			4							(Rn8 of @ERd)"1	_	_	_	_			4	1
	BSET Rn,@aa:8	В						4				(Rn8 of @aa:8)"1	_		_	_			4	1
	BSET Rn,@aa:16	В						6				(Rn8 of @aa:16)"1	_		_	_			5	5
	BSET Rn,@aa:32	В						8				(Rn8 of @aa:32)"1	_	_	_	_	_	_	6	3
BCLR	BCLR #xx:3,Rd	В		2								(#xx:3 of Rd8)"0	_			_			1	
	BCLR #xx:3,@ERd	В			4							(#xx:3 of @ERd)"0	_	_	_	_	_	_	4	1
	BCLR #xx:3,@aa:8	В						4				(#xx:3 of @aa:8)"0	_	_	_	_	_	_	4	1
	BCLR #xx:3,@aa:16	В						6				(#xx:3 of @aa:16)"0	_		_	_			5	5
	BCLR #xx:3,@aa:32	В						8				(#xx:3 of @aa:32)"0	_		_	_			6	3
	BCLR Rn,Rd	В		2								(Rn8 of Rd8)"0	_		_				1	
	BCLR Rn,@ERd	В			4							(Rn8 of @ERd)"0	_			_			4	1
	BCLR Rn,@aa:8	В						4				(Rn8 of @aa:8)"0	_		_				4	1
	BCLR Rn,@aa:16	В						6				(Rn8 of @aa:16)"0	_	_	_	_	_	_	5	5
	BCLR Rn,@aa:32	В						8				(Rn8 of @aa:32)"0	_	_	_	_	_	_	6	6
BNOT	BNOT #xx:3,Rd	В		2								(#xx:3 of Rd8)" [¬(#xx:3 of Rd8)]	_		_	_			1	1
	BNOT #xx:3,@ERd	В			4							(#xx:3 of @ERd)" [¬(#xx:3 of @ERd)]	_		_				4	1

			A	ddre		ng I engt				ıctio	on		·	Con	ditio	on (Code	е	No. Stat	
M	nemonic	Operand Size	#xx	Rn	@ERn	@(d,ERn)	@-ERn/@ERn+	@aa	@(d,PC)	@ @ aa	1	Operation	ı	H	Ν	Z	٧	С	Normal	Advanced
	BNOT #xx:3,@aa:8	В						4				(#xx:3 of @aa:8)" [¬(#xx:3 of @aa:8)]	_			_	_		4	1
	BNOT #xx:3,@aa:16	В						6				(#xx:3 of @aa:16)" [¬(#xx:3 of @aa:16)]	_				_		5	5
	BNOT #xx:3,@aa:32	В						8				(#xx:3 of @aa:32)" [¬(#xx:3 of @aa:32)]	_			_	_		6	ó
	BNOT Rn,Rd	В		2								(Rn8 of Rd8)" [¬(Rn8 of Rd8)]	_			_	_		1	1
	BNOT Rn,@ERd	В			4							(Rn8 of @ERd)" [¬(Rn8 of @ERd)]	_			_	_		4	1
	BNOT Rn,@aa:8	В						4				(Rn8 of @aa:8)" [¬(Rn8 of @aa:8)]	_			_	_		4	1
	BNOT Rn,@aa:16	В						6				(Rn8 of @aa:16)" [¬(Rn8 of @aa:16)]	_			_	_		Ę	5
	BNOT Rn,@aa:32	В						8				(Rn8 of @aa:32)" [¬(Rn8 of @aa:32)]	_			_	_		6	6
BTST	BTST #xx:3,Rd	В		2								(#xx:3 of Rd8)ÆZ				\Diamond			,	1
	BTST #xx:3,@ERd	В			4							(#xx:3 of @ERd)ÆZ	_			\	_		3	3
	BTST #xx:3,@aa:8	В						4				(#xx:3 of @aa:8)ÆZ				\			3	3
	BTST #xx:3,@aa:16	В						6				(#xx:3 of @aa:16)ÆZ	_			\(\)			4	1
	BTST #xx:3,@aa:32	В						8				(#xx:3 of @aa:32)ÆZ	_			\rightarrow	_		5	5
	BTST Rn,Rd	В		2								(Rn8 of Rd8)ÆZ				\Diamond			,	1
	BTST Rn,@ERd	В			4							(Rn8 of @ERd)ÆZ	_			\Q	_		3	3
	BTST Rn,@aa:8	В						4				(Rn8 of @aa:8)ÆZ	_			\rightarrow	_		3	3
	BTST Rn,@aa:16	В						6				(Rn8 of @aa:16)ÆZ	_			\(\)	_		4	4
	BTST Rn,@aa:32	В						8				(Rn8 of @aa:32)ÆZ	_			\				5
BLD	BLD #xx:3,Rd	В		2								(#xx:3 of Rd8)ÆC	_	_	_	_	_	\Diamond	,	1
	BLD #xx:3,@ERd	В			4							(#xx:3 of @ERd)ÆC	_					\rightarrow	3	3
	BLD #xx:3,@aa:8	В						4				(#xx:3 of @aa:8)ÆC						\Diamond	3	3
	BLD #xx:3,@aa:16	В						6				(#xx:3 of @aa:16)ÆC	_	_	_	_	_	\Diamond		1
	BLD #xx:3,@aa:32	В						8				(#xx:3 of @aa:32)ÆC	_				_	\rightarrow	5	5

			A	ddre	essi Le	_	th (E			ıcti	on			Con	diti	on (Code	е	No. Stat	
M	nemonic	Operand Size	xx#	Rn	@ERn	@(d,ERn)	@-ERn/@ERn+	@aa	@(d,PC)	@ @ aa	I	Operation	I	Н	N	z	>	С	Normal	Advanced
BILD	BILD #xx:3,Rd	В		2								¬(#xx:3 of Rd8)ÆC	_			_		\Diamond	1	1
	BILD #xx:3,@ERd	В			4							¬(#xx:3 of @ERd)ÆC	_	_		_		\Diamond	3	3
	BILD #xx:3,@aa:8	В						4				¬(#xx:3 of @aa:8)ÆC	_	_				\	3	3
	BILD #xx:3,@aa:16	В						6				¬(#xx:3 of @aa:16)ÆC	_	_	_			\	4	ţ
	BILD #xx:3,@aa:32	В						8				¬(#xx:3 of @aa:32)ÆC	_	_		_		\	5	5
BST	BST #xx:3,Rd	В		2								CÆ(#xx:3 of Rd8)	_		_	_			1	
	BST #xx:3,@ERd	В			4							CÆ(#xx:3 of @ERd24)	_	_				_	4	ł
	BST #xx:3,@aa:8	В						4				CÆ(#xx:3 of @aa:8)	_	_				_	4	1
	BST #xx:3,@aa:16	В						6				CÆ(#xx:3 of @aa:16)	_	_	_	_		_	5	5
	BST #xx:3,@aa:32	В						8				CÆ(#xx:3 of @aa:32)	_	_	_	_		_	6	;
BIST	BIST #xx:3,Rd	В		2								¬CÆ(#xx:3 of Rd8)	_	_	_	_	_	_	1	1
	BIST #xx:3,@ERd	В			4							¬CÆ(#xx:3 of @ERd24)	_	_	_	_		_	4	1
	BIST #xx:3,@aa:8	В						4				¬CÆ(#xx:3 of @aa:8)	_	_		_		_	4	1
	BIST #xx:3,@aa:16	В						6				¬CÆ(#xx:3 of @aa:16)	_	_		_		_	5	5
	BIST #xx:3,@aa:32	В						8				¬CÆ(#xx:3 of @aa:32)	_	_		_		_	6	3
BAND	BAND #xx:3,Rd	В		2								CŸ(#xx:3 of Rd8)ÆC	_	_			_	\Diamond	1	1
	BAND #xx:3,@ERd	В			4							CŸ(#xx:3 of @ERd24)ÆC	_	_	_	_		\	3	3
	BAND #xx:3,@aa:8	В						4				CŸ(#xx:3 of @aa:8)ÆC	_	_				\	3	3
	BAND #xx:3,@aa:16	В						6				CŸ(#xx:3 of @aa:16)ÆC	_	_	_			\rightarrow	2	1
	BAND #xx:3,@aa:32	В						8				CŸ(#xx:3 of @aa:32)ÆC	_	_			_	\rightarrow	5	5
BIAND	BIAND #xx:3,Rd	В		2								CŸ [¬(#xx:3 of Rd8)]ÆC	_	_				\rightarrow		1
	BIAND #xx:3,@ERd	В			4							CŸ [¬(#xx:3 of @ERd24)]ÆC	_					\Q	3	3
	BIAND #xx:3,@aa:8	В						4				CŸ [¬(#xx:3 of @aa:8)]ÆC	_					\rightarrow	3	3

			A	ddr			Mod th (E			ıctio	on		·	Con	diti	on C	Code	9	No. Stat	- 4
!	Mnemonic	Operand Size	xx#	Rn	@ERn	@(d,ERn)	@-ERn/@ERn+	@aa	@(d,PC)	@ @aa	1	Operation	ı	н	N	Z	٧	С	Normal	Advanced
	BIAND #xx:3,@aa:16	В						6				CŸ [¬(#xx:3 of @aa:16)]ÆC	_	_	_	_		\Diamond	4	1
	BIAND #xx:3,@aa:32	В						8				CŸ [¬(#xx:3 of @aa:32)]ÆC						\Q	5	5

			A	ddre	essi Le		h (E	le/In Byte		ucti	on		(Con	diti	on (Code	е	No. Stat	
м	nemonic	Operand Size	xx#	Rn	@ERn	@(d,ERn)	@-ERn/@ERn+	@аа	@(d,PC)	@ @aa	ı	Operation	ı	н	N	z	٧	С	Normal	Advanced
BOR	BOR #xx:3,Rd	В		2								C/(#xx:3 of Rd8)ÆC	_	_		_	_	\Diamond	1	
	BOR #xx:3,@ERd	В			4							C/(#xx:3 of @ERd24)ÆC						\rightarrow	3	3
	BOR #xx:3,@aa:8	В						4				C/(#xx:3 of @aa:8)ÆC	_			_	_	\	3	3
	BOR #xx:3,@aa:16	В						6				C/(#xx:3 of @aa:16)ÆC	—	_		_		\Q	2	1
	BOR #xx:3,@aa:32	В						8				C/(#xx:3 of @aa:32)ÆC	—	_		_		\Q	5	5
BIOR	BIOR #xx:3,Rd	В		2								C/[¬(#xx:3 of Rd8)]ÆC						\Diamond	1	l
	BIOR #xx:3,@ERd	В			4							C/[¬(#xx:3 of @ERd24)]ÆC	_			_		\Diamond	3	3
	BIOR #xx:3,@aa:8	В						4				C/ [¬(#xx:3 of @aa:8)]ÆC	_	_	_	_		\	3	3
	BIOR #xx:3,@aa:16	В						6				C/[¬(#xx:3 of @aa:16)]ÆC		_				\	4	1
	BIOR #xx:3,@aa:32	В						8				C/[¬(#xx:3 of @aa:32)]ÆC	_	_		_		\Diamond	5	5
BXOR	BXOR #xx:3,Rd	В		2								C≈ (#xx:3 of Rd8)ÆC	_	_	_	_		\	1	
	BXOR #xx:3,@ERd	В			4							C≈ (#xx:3 of @ERd24)ÆC	_	_	_	_		\	3	3
	BXOR #xx:3,@aa:8	В						4				C≈ (#xx:3 of @aa:8)ÆC	_	_	_	_		\	3	3
	BXOR #xx:3,@aa:16	В						6				C≈ (#xx:3 of @aa:16)ÆC	_					\Diamond	2	1
	BXOR #xx:3,@aa:32	В						8				C≈ (#xx:3 of @aa:32)ÆC	_					\Diamond	5	5
BIXOR	BIXOR #xx:3,Rd	В		2								C≈ [¬(#xx:3 of Rd8)]ÆC	_					\Diamond	1	ı
	BIXOR #xx:3,@ERd	В			4							C≈ [¬(#xx:3 of @ERd24)]ÆC	_					\Diamond	3	3
	BIXOR #xx:3,@aa:8	В						4				C≈ [¬(#xx:3 of @aa:8)]ÆC	_					\Diamond	63	3
	BIXOR #xx:3,@aa:16	В						6				C≈ [¬(#xx:3 of @aa:16)]ÆC	_			_		\rightarrow	2	1
	BIXOR #xx:3,@aa:32	В						8				C≈ [¬(#xx:3 of @aa:32)]ÆC	_					\rightarrow	5	5

6. Branch instructions

			A	ddr		eng	th (I	le/In 3yte		uctio	on	Оре	eration	(Con	diti	on (Cod	le	No. Stat	of es ^{*1}
м	nemonic	Operand Size	xx#	Rn	@ERn	@(d,ERn)	@-ERn/@ERn+	@aa	@(q,PC)	@ @ aa	1		Branching Conditions	I	Н	N	z	٧	С	Normal	Advanced
Всс	BRA d:8(BT d:8)	_							2			if condition is true	Always		_				_	4	2
	BRA d:16(BT d:16)	_							4			then			_				_	,	3
	BRN d:8(BF d:8)	_							2			PC"PC+d	Never		_				_	2	2
	BRN d:16(BF d:16)	_							4			else next;			_					,	3
	BHI d:8	_							2				C/Z=0							2	2
	BHI d:16	_							4											,	3
	BLS d:8	_							2				C/Z=1		_					2	2
	BLS d:16	_							4						_				_	í	3
	BCC d:8(BHS d:8)	_							2				C=0		_				_	2	2
	BCC d:16(BHS d:16)	_							4						_					,	3
	BCS d:8(BLO d:8)	_							2				C=1		_					2	2
	BCS d:16(BLO d:16)	_							4						_				_	(3
	BNE d:8	_							2				Z=0		_				_	2	2
	BNE d:16	_							4						_		_		_	(3
	BEQ d:8	_							2				Z=1		_			_	_	2	2
	BEQ d:16	_							4						_				_	(3
	BVC d:8	_							2				V=0		_		_		_	2	2
	BVC d:16	_							4												3
	BVS d:8	_							2				V=1								2
	BVS d:16	_							4												3
	BPL d:8	_							2				N=0							_ 2	2
	BPL d:16	_							4						E						3
	BMI d:8	_							2				N=1							_ 2	2
	BMI d:16	_							4										_	(3
	BGE d:8	_							2				N≈V=0								2
	BGE d:16	_							4												3
	BLT d:8	_							2				N≈V=1							_ 2	2
	BLT d:16	_							4												3
4 8	BGT d:8	_							2				Z/(N≈V)=0						_	2	2

			A	ddre		_	Mod th (E			ıctio	on	Ope	eration	0	Con	diti	on (Cod	le	No. Stat	of es ^{*1}
Mı	nemonic	Operand Size	xx#	Rn	@ERn	@(d,ERn)	@-ERn/@ERn+	@aa	@(d,PC)	@ @ aa	I		Branching Conditions	ı	н	N	z	V	С	Normal	Advanced
	BGT d:16	_							4											``	3
	BLE d:8	_							2				Z/(N≈V)=1		_				_	. 4	2
	BLE d:16	_							4						_			_	_	(3
JMP	JMP @ERn	_			2							PC"ERn		_	_		_	_	_	2	2
	JMP @aa:24	_						4				PC"aa:24			_			_	_	(3
	JMP @@aa:8	_								2		PC"@aa:8	}	_						4	5
BSR	BSR d:8	_							2			PCÆ@- SP,PC"PC	C+d:8							3	4
	BSR d:16	_							4			PCÆ@- SP,PC"PC	C+d:16		_				_	4	5
JSR	JSR @ERn	_			2							PCÆ@-S	P,PC"ERn							3	4
	JSR @aa:24	_						4				PCÆ@-S	P,PC"aa:24		_			_	_	4	5
	JSR @@aa:8	_								2		PCÆ@- SP,PC¨@a	aa:8		_				_	4	6
RTS	RTS	_									2	PC"@SP+	-		_				_	4	5

7. System control instructions

			A	ddr	essi Le	eng	h (E	le/In Byte		ıctio	on		,	Con	diti	on C	Code	е	No. Stat	of es ^{*1}
М	nemonic	Operand Size	xx#	Rn	@ERn	@(d,ERn)	@-ERn/@ERn+	@aa	@(d,PC)	@ @ aa	I	Operation	I	Н	N	Z	v	С	Normal	Advanced
TRAPA	TRAPA #xx:2										2	PCÆ@- SP,CCRÆ@-SP, EXRÆ@- SP, <vector>ÆPC</vector>	1						7 [10]	8 [10]
RTE	RTE	_										EXR"@SP+,CCR"@ SP+, PC"@SP+	◊	◊	◊	\rightarrow	◊	\rightarrow	5 [10]
SLEEP	SLEEP	_										Transition to power- down state		_		_			2	2
LDC	LDC #xx:8,CCR	В	2									#xx:8ÆCCR	\Diamond	\Diamond	\Diamond	\Diamond	\Diamond	\Diamond	,	1
	LDC #xx:8,EXR	В	4									#xx:8ÆEXR	_		_		_		2	2
	LDC Rs,CCR	В		2								Rs8ÆCCR	\Diamond	\(\)	\(\)	\Diamond	\(\)	\Diamond	,	1
	LDC Rs,EXR	В		2								Rs8ÆEXR	_						1	1
	LDC @ERs,CCR	W			4							@ERsÆCCR	\	\(\)	\(\)	\	\(\)	\	3	3
	LDC @ERs,EXR	W			4							@ERsÆEXR	_	_	_	_	_	_	3	3
	LDC @(d:16,ERs),C CR	W				6						@(d:16,ERs)ÆCCR	◊	◊	◊	◊	◊	◊	4	4
	LDC @(d:16,ERs),E XR	W				6						@(d:16,ERs)ÆEXR		_	_	_	_	_	4	1
	LDC @(d:32,ERs),C CR	W				10						@(d:32,ERs)ÆCCR	\rightarrow	◊	◊	\(\)	◊	◊	(3
	LDC @(d:32,ERs),E XR	W				10						@(d:32,ERs)ÆEXR		_	_		_		6	6
	LDC @ERs+,CCR	W					4					@ERsÆCCR,ERs32 +2ÆERs32	\Diamond	◊	◊	\Diamond	\	\Diamond	4	4
	LDC @ERs+,EXR	W					4					@ERsÆEXR,ERs32 +2ÆERs32	_			_		_	4	1
	LDC @aa:16,CCR	W						6				@aa:16ÆCCR	\Q	◊	◊	\	◊	\Q	4	1
	LDC @aa:16,EXR	W						6				@aa:16ÆEXR							4	4
	LDC @aa:32,CCR	W						8				@aa:32ÆCCR	\(\)	◊	◊	\	◊	\	Ę	5
	LDC @aa:32,EXR	W						8				@aa:32ÆEXR	_					_	Ę	5
STC	STC CCR,Rd	В		2								CCRÆRd8							1	1

			Α	ddre	essi Le		th (E			ıcti	on			Con	diti	on C	ode	e	No. o	of es ^{*1}
М	Inemonic	Operand Size	xx#	Rn	@ERn	@(d,ERn)	@-ERn/@ERn+	@aa	@(d,PC)	@ @aa	I	Operation	I	н	N	Z	>	C	Normal	Advanced
	STC EXR,Rd	В		2								EXRÆRd8		_	_	_	_	_	1	
	STC CCR,@ERd	W			4							CCRÆ@ERd	_	_	_			_	3	}
	STC EXR,@ERd	W			4							EXRÆ@ERd		_				_	(3)	}
	STC CCR,@(d:16,E Rd)	W				6						CCRÆ@(d:16,ERd)							4	r
	STC EXR,@(d:16,E Rd)	W				6						EXRÆ@(d:16,ERd)							4	ļ
	STC CCR,@(d:32,E Rd)	W				10						CCRÆ@(d:32,ERd)		_	_			_	6	;
	STC EXR,@(d:32,E Rd)	W				10						EXRÆ@(d:32,ERd)		_	_			_	6	;
	STC CCR,@- ERd	W					4					ERd32- 2ÆERd32,CCRÆ@ ERd	_	_	_				4	ŀ
	STC EXR,@- ERd	W					4					ERd32- 2ÆERd32,EXRÆ@ ERd		_	_			_	4	-
	STC CCR,@aa:16	W						6				CCRÆ@aa:16	_		_			_	4	ŀ
	STC EXR,@aa:16	W						6				EXRÆ@aa:16	_	_	_			_	4	ŀ
	STC CCR,@aa:32	W						8				CCRÆ@aa:32	_	_	_			_	5	;
	STC EXR,@aa:32	W						8				EXRÆ@aa:32	_	_	_	_	_	_	5	,
ANDC	ANDC #xx:8,CCR	В	2									CCRŸ#xx:8ÆCCR	\(\)	\(\)	◊	\Diamond	\rightarrow	\Q	1	
	ANDC #xx:8,EXR	В	4									EXRŸ#xx:8ÆEXR	-	_	$\overline{}$			_	2	<u>, </u>
ORC	ORC #xx:8,CCR	В	2									CCR/#xx:8ÆCCR	\(\)	\(\)	◊	\Diamond	\rightarrow	\(\)	1	
	ORC #xx:8,EXR	В	4									EXR/#xx:8ÆEXR	-	_	$\overline{}$			_	2	<u>, </u>
XORC	XORC #xx:8,CCR	В	2									CCR≈#xx:8ÆCCR	\(\)	\(\)	◊	\rightarrow	\rightarrow	\rightarrow	1	
	XORC #xx:8,EXR	В	4									EXR≈#xx:8ÆEXR	_					_	2	<u>,</u>
NOP	NOP	_									2	PC"PC+2	_	_			_		1	

8. Program transfer instructions

			A	ddre	essi Le	_	Mod th (E			uctio	on		•	Con	ditio	on C	ode	е	No. o	of es ^{*1}
Mr	nemonic	Operand Size	хх#	Rn	@ERn	@(d,ERn)	@-ERn/@ERn+	@ aa	@(d,PC)	@ @ aa	ı	Operation	ı	н	N	Z	٧	С	Normal	Advanced
EEPMOV	EEPMOV.B	_									4	if R4L≠0 Repeat @ER5+Æ@ER6+ R5+1ÆR5 R6+1ÆR6 R4L-1ÆR4L Until R4L=0 else next;							4+2	n*2
	EEPMOV.W	_									4	if R4≠0 Repeat @ER5+Æ@ER6+ R5+1ÆR5 R6+1ÆR6 R4–1ÆR4 Until R4=0 else next;							4+2	n*2

Notes: *1 The number of states is the number of states required for execution when the instruction and its operands are located in on-chip memory.

- *2 n is the initial value of R4L or R4.
- [1] Seven states for saving or restoring two registers, nine states for three registers, or eleven states for four registers.
- [2] Cannot be used in the H8S/2655 Series.
- [3] Set to 1 when a carry or borrow occurs at bit 11; otherwise cleared to 0.
- [4] Set to 1 when a carry or borrow occurs at bit 27; otherwise cleared to 0.
- [5] Retains its previous value when the result is zero; otherwise cleared to 0.
- [6] Set to 1 when the divisor is negative; otherwise cleared to 0.
- [7] Set to 1 when the divisor is zero; otherwise cleared to 0.
- [8] Set to 1 when the quotient is negative; otherwise cleared to 0.
- [9] MAC instruction results are indicated in the flags when the STMAC instruction is executed.
- [10]One additional state is required for execution when EXR is valid.

Number of States Required for Execution

The number of states shown in the instruction set table is the number of states required for execution when the op code and operand data are located in a one-cycle area on which word access is possible, such as on-chip memory. When the op code or operand data is accessed from an on-chip supporting module or an external address, the number of states increases as shown in the table below.

		Access	Conditio	ns				
			On-Chip Support		External	Data Bus		
Cycle		On-Chip	Module		8-Bit Bus	S	16-Bit Bu	ıs
		Memory	8-Bit Bus	16-Bit Bus	2-State Access	3-State Access	2-State Access	3-State Access
Instruction fetch SI								
Branch address read	SJ		2n		4	6+2m		
Stack operation SK		1		n			2	3+m*
Byte data access SL			n	1	2	3+m		
Word data access	SM		2n		4	6+2m	1	
Internal operation SN		1	•	•	•	•	•	•

Note: * Refer to the hardware manual for the relevant product for details of MOVFPE and MOVTPE.

Legend

- m: Number of wait states inserted into external device access
- n: Number of wait states inserted into on-chip supporting module access. Refer to the hardware manual for the relevant product for the actual number.

Condition Code Notation

Symbol	Meaning		
‡	Changes according to the result of instruction execution		
*	Undetermined (no guaranteed value)		
0	Always cleared to 0		
1	Always set to 1		
_	Not affected by execution of the instruction		
Δ	Varies depending on conditions; see the notes		

Operation Notation

Rd	General register (destination)*		
Rs	General register (source)*		
Rn	General register*		
ERn	General register (32-bit register)		
MAC	Multiply-and-accumulate register (32-bit register)		
(EAd)	Destination operand		
(EAs)	Source operand		
EXR	Extend register		
CCR	Condition code register		
N	N (negative) flag of CCR		
Z	Z (zero) flag of CCR		
V	V (overflow) flag of CCR		
С	C (carry) flag of CCR		
PC	Program counter		
SP	Stack pointer		
#IMM	Immediate data		
disp	Displacement		
+	Addition		
_	Subtraction		
х	Multiplication		
÷	Division		
Ÿ	AND logical		
/	OR logical		
≈	Exclusive OR logical		
Æ	Transfer from left-hand operand to right-hand operand, or transition from left-hand state to right-hand state		
7	NOT (logical complement)		
() <>	Operand contents		
:8/:16/:24/:32	8-, 16-, 24-, or 32-bit length		

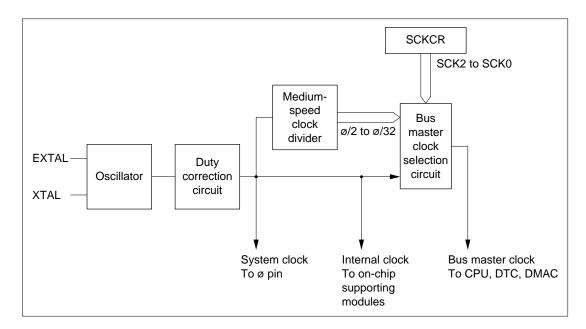
Note: * General registers include 8-bit registers (R0H to R7H, R0L to R7L), 16-bit registers (R0 to R7, E0 to E7), and 32-bit registers (ER0 to ER7).

2.6 Basic Bus Timing

The CPU operates on the basis of the system clock (Ø). One Ø clock cycle is called a state, and a bus cycle consists of one, two, or three states. Different access methods are used for on-chip memory, on-chip supporting modules, and external devices.

Basic Clock Timing

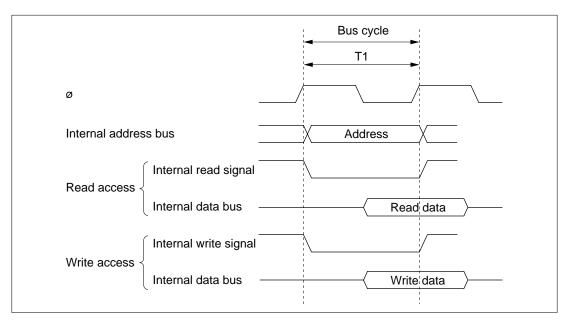
An external clock is input to the EXTAL pin, or a crystal oscillator is connected to the EXTAL pin, to generate the system clock (\emptyset) . An external clock or crystal oscillator of the same frequency as the \emptyset clock should be used.



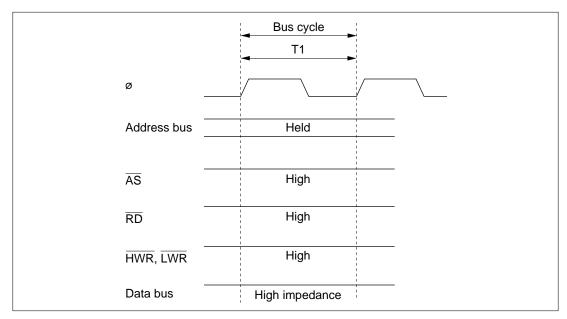
CPU Read/Write Cycles

The CPU operates on the basis of the system clock (\emptyset) . One \emptyset clock cycle is called a state, and a bus cycle consists of one, two, or three states. Different access methods are used for on-chip memory, on-chip supporting modules, and external devices. Access to the external address space can be controlled by the bus controller.

On-Chip Memory: On-chip memory is accessed in one state. The data bus is 16 bits wide, permitting both byte and word access.

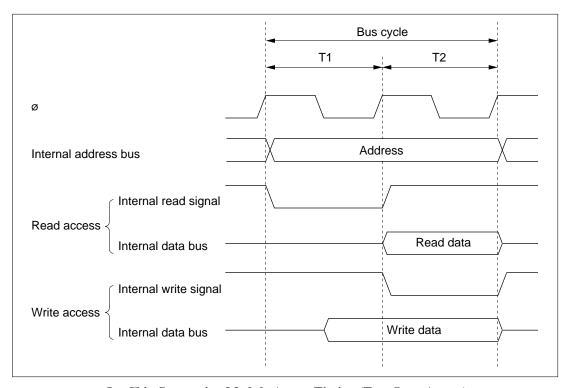


On-Chip Memory Access Cycle (One-State Access)

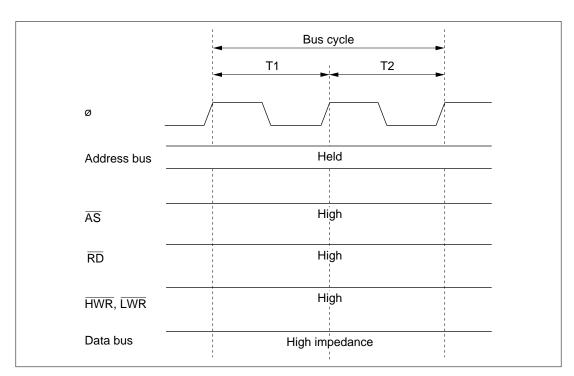


Pin States during On-Chip Memory Access

On-Chip Supporting Module: The on-chip supporting modules are accessed in two states. The data bus is 8 or 16 bits wide, depending on the internal I/O register being accessed.



On-Chip Supporting Module Access Timing (Two-State Access)



Pin States during On-Chip Supporting Module Access

External Address Space: The external address space is accessed via an 8-bit or 16-bit bus, and in two or three states. Wait state insertion is possible in the case of 3-state access. See the Bus Controller section for details.

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2.7 Processing States

The H8S/2600 CPU has five processing states: the reset state, program execution state, exception-handling state, bus-released state, and power-down state.

Reset State

State in which the CPU and all on-chip supporting modules are initialized and halted

Program Execution State

State in which the CPU executes the program sequentially

Exception-Handling State

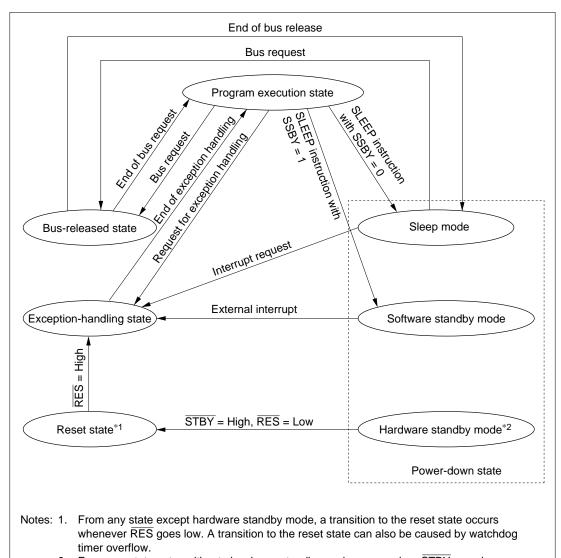
Transient state in which exception handling is executed as the result of an reset, interrupt, or trap instruction exception handling source

Bus-Released State

State in which the external bus is released in response to a bus request signal from a bus master other than the CPU

Power-Down State

State in which CPU operation is stopped, and power consumption is kept low (sleep mode, software standby mode, hardware standby mode). The power-down state also includes medium-speed mode and module stop mode.



2. From any state, a transition to hardware standby mode occurs when $\overline{\text{STBY}}$ goes low.

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2.8 Exception Handling

H8S/2600 CPU exception handling is initiated by a reset, a trap instruction, or an interrupt. A priority system is provided for exception handling, and simultaneously generated exceptions are handled in order of priority.

Exception Handling Types and Priorities

Priority	y Exception Type Start of Exception Handling		
High	Reset	After a low-to-high transition at the RES pin, or when the watchdog timer overflows	
	Trace	After instruction or exception handling execution when the trace (T) bit is 1	
	Interrupt	When an interrupt is generated, after instruction or exception handling execution	
Low	Trap instruction (TRAPA)	When a trap (TRAPA) instruction is executed	

Exception Handling Operation

Exception handling is started by any of the exception handling sources. Trap instruction exception handling is always accepted in the program execution state.

The operations in trap instruction and interrupt exception handling are as follows.

- (1) The program counter (PC), condition code register (CCR), and extended register (EXR) are saved on the stack.
- (2) The interrupt mask bit is updated, and the T bit is cleared to 0.
- (3) The vector address corresponding to the activation source is generated, and program execution is started from the address indicates by the contents of the vector address.

In reset exception handling, only operations (2) and (3) are performed.

Exception Vector Table

Vector Address*1

Exception Source		Vector Number	Normal Mode	Advanced Mode
Power-on reset		0	H'0000-H'0001	H'0000-H'0003
Manual reset		1	H'0002-H'0003	H'0004-H'0007
Reserved for system use		2	H'0004-H'0006	H'0008-H'000B
		3	H'0006-H'0007	H'000C-H'000F
		4	H'0008-H'0009	H'0010-H'0013
Trace		5	H'000A-H'000B	H'0014-H0017
Reserved for system use		6	H'000C-H'000D	H'0018-H001B
External interrupt NMI		7	H'000E-H'000F	H'001C-H'001F
Trap instruction (4 sources)		8	H'0010-H'0011	H'0020-H'0023
		9	H'0012-H'0013	H'0024-H'0027
		10	H'0014-H'0015	H'0028-H'002B
		11	H'0016-H'0017	H'002C-H'002F
Reserved for system use		12	H'0018-H'0019	H'0030-H'0033
		13	H'001A-H'001B	H'0034-H'0037
		14	H'001C-H'001D	H'0038-H'003B
		15	H'001E-H'001F	H'003C-H'003F
External interrupt	IRQ0	16	H'0020-H'0021	H'0040-H'0043
	IRQ1	17	H'0022-H'0023	H'0044-H'0047
	IRQ2	18	H'0024-H'0025	H'0048-H'004B
	IRQ3	19	H'0026-H'0027	H'004C-H'004F
	IRQ4	20	H'0028-H'0029	H'0050-H'0053
	IRQ5	21	H'002A-H'002B	H'0054-H'0057
	IRQ6	22	H'002C-H'002D	H'0058-H'005B
	IRQ7	23	H'002E-H'002F	H'005C-H'005F
Internal interrupt*2		24 to 91	H'0030-H'0031 to H'00B6-H'00B7	H'0060-H'0063 to H'016C-H'019F

Notes: 1. Lower 16 bits of address

2. See the Interrupt Exception Vector Table for the internal interrupt vector table.

2.9 Interrupts

This section describes the sru interrupt, one of the external interrupt sources.

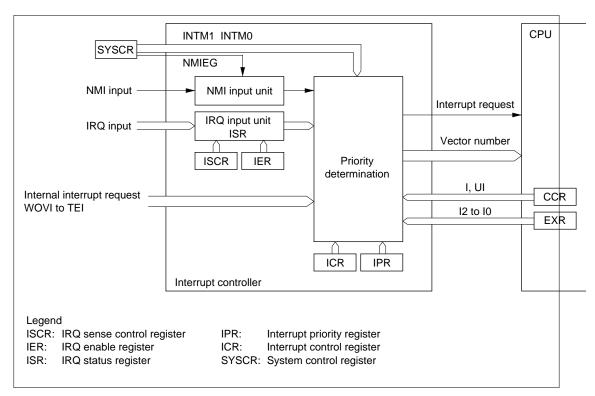
Interrupts are controlled by the interrupt controller. There are a total of 62 interrupt sources, comprising nine external interrupts from the external pins (NMI, IRQ_0 to IRQ_7), and 53 internal interrupts from on-chip supporting modules. A separate vector number is assigned to each interrupt.

Interrupt Control

Any of four interrupt control modes can be set by means of the INTM1 and INTM0 bits in the system control register (SYSCR).

The interrupt controller controls interrupts on the basis of the control mode set by the INTM1 and INTM0 bits, the interrupt priorities set by interrupt control register (ICR) and interrupt priority register (IPR), and the masking conditions set by the I and UI bits in CCR and bits I2 to I0 in EXR.

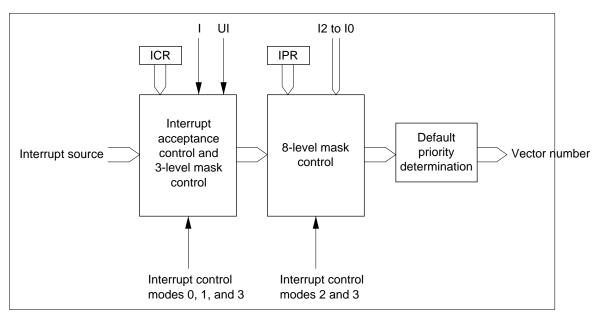
NMI is the highest-priority interrupt, and is always accepted.



Block Diagram of Interrupt Controller

Interrupt Control Modes

Interrupt Control		SCR	Priority Setting	Interrupt			
Mode	INTM1	INTM0	Registers	Mask Bits	Description		
0	0	0	ICR	1	Interrupt mask control is performed by the I bit.		
	_				Priority can be set with ICR.		
1		1	ICR	I, UI	3-level interrupt mask control is performed by the I and UI bits.		
					Priority can be set with ICR.		
2	1	0	IPR	I2 to I0	8-level interrupt mask control is performed by bits I2 to I0.		
					8 priority levels can be set with IPR.		
3	_	1	ICR, IPR	I, UI, I2 to I0	Control is performed by a combination of interrupt masking set by the I and UI bits and priority setting by ICR, based on 8-level interrupt mask control performed by bits I2 to I0 and 8-level priority setting by IPR.		



Block Diagram of Interrupt Control Operation

Interrupt Control Mode 0

Enabling and disabling of IRQ interrupts and on-chip supporting module interrupts can be set by means of the I bit in CCR. Control level setting can be performed with ICR. Interrupts are enabled when the I bit is cleared to 0, and disabled when set to 1.

Control level 1 interrupt sources have higher priority.

Interrupt Control Mode 1

Three-level masking can be implemented for IRQ interrupts and on-chip supporting module interrupts by means of the I and UI bits in CCR, and ICR.

- (1) Control level 0 interrupt requests are enabled when the I bit is cleared to 0, and disabled when set to 1.
- (2) Control level 1 interrupt requests are enabled when the I bit or UI bit is cleared to 0, and disabled when both the I bit and the UI bit are set to 1.

Interrupt Control Mode 2

Eight-level masking can be implemented for IRQ interrupts and on-chip supporting module interrupts by comparing the interrupt mask level bits (I2 to I0) in EXR and the IPR priority level.

Interrupt Control Mode 3

Control of IRQ interrupts and on-chip supporting module interrupts is performed by a combination of interrupt masking set by the I and UI bits and control level setting by ICR, based on 8-level interrupt mask control performed by comparing the interrupt mask level bits (I2 to I0) in EXR and the IPR priority level.

The following control is performed in addition to (1) and (2) in interrupt control mode 1:

(3) Eight-level priority control is performed for interrupt requests enabled in (1) and (2).

Interrupt Sources, Vector Addresses, and Interrupt Priorities

	Origin of		Vector	Address*			
Interrupt Source	Interrupt Source	Vector Number	Normal Mode	Advanced Mode	IPR	Priority	
NMI	External	7	H'000E	H'001C			High
IRQ_0	pin	16	H'0020	H'0040	ICRA7	PRA6-IPRA4	\uparrow
IRQ ₁		17	H'0022	H'0044	ICRA6	IPRA2-IPRA0	
IRQ ₂		18	H'0024	H'0048	ICRA5	IPRB6-IPRB4	
IRQ ₃		19	H'0026	H'004C			
IRQ ₄		20	H'0028	H'0050	ICRA4	IPRB2-IPRB0	
IRQ ₅		21	H'002A	H'0054			
IRQ ₆		22	H'002C	H'0058	ICRA3	IPRC6-IPRC4	
IRQ ₇		23	H'002E	H'005C			
SWDTEND (software activation interrupt end	DTC)	24	H'0030	H'0060	ICRA2	IPRC2-IPRC0	
WOVI (interval timer)	Watchdog timer	25	H'0032	H'0064	ICRA1	IPRD6-IPRD4	
CMI (compare-match)	Refresh controller	26	H'0034	H'0068	ICRA0	IPRD2-IPRD0	
ADI (A/D conversion end)	A/D	28	H'0038	H'0070	ICRB6	IPRE2-IPRE0	
TGI0A (TGR0A input capture/ compare-match)	TPU channel 0	32	H'0040	H'0080	ICRB5	PRF6–IPRF4	
TGI0B (TGR0B input capture/ compare-match)		33	H'0042	H'0084			
TGI0C (TGR0C input capture/ compare-match)		34	H'0044	H'0088			
TGI0D (TGR0D input capture/ compare-match)		35	H'0046	H'008C			
TCI0V (overflow 0)		36	H'0048	H'0090			Low

Note: *Lower 16 bits of the start address.

	Origin of		Vector	Address*			
	Interrupt						
Interrupt Source	Source	Number	Mode	d Mode	ICR	IPR	Priority

TGI1A (TGR1A input capture/ compare-match)	TPU channel 1	40	H'0050	H'00A0	ICRB4 IPF	RF2-IPRF0	High	
TGI1B (TGR1B input capture/ compare-match)		41	H'0052	H'00A4				
TCI1V (overflow 1)		42	H'0054	H'00A8				
TCI1U (underflow 1)		43	H'0056	H'00AC				
TGI2A (TGR2A input capture/ compare-match)	TPU channel 2	44	H'0058	H'00B0	ICRB3 IPF	RG6–IPRG4		
TGI2B (TGR2B input capture/ compare-match)		45	H'005A	H'00B4				
TCI2V (overflow 2)		46	H'005C	H'00B8				
TCI2U (underflow 2)		47	H'005E	H'00BC				
GI3A (TGR3A input capture/ compare-match)	TPU channel 3	48	H'0060	H'00C0	ICRB2 IPF	RG2–IPRG0		
TGI3B (TGR3B input capture/ compare-match)		49	H'0062	H'00C4				
TGI3C (TGR3C input capture/ compare-match)		50	H'0064	H'00C8				
TGI3D (TGR3D input capture/ compare-match)		51	H'0066	H'00CC				
TCI3V (overflow 3)		52	H'0068	H'00D0				
TGI4A (TGR4A input capture/ compare-match)	TPU channel 4	56	H'0070	H'00E0	ICRB1 IPF	RH6–IPRH4		
TGI4B (TGR4B input capture/ compare-match)		57	H'0072	H'00E4				
TCI4V (overflow 4)		58	H'0074	H'00E8				
TCI4U (underflow 4)		59	H'0076	H'00EC				
TGI5A (TGR5A input capture/ compare-match)	TPU channel 5	60	H'0078	H'00F0	ICRB0 IPF	RH2–IPRH0		
TGI5B (TGR5B input capture/ compare-match)		61	H'007A	H'00F4				
TCI5V (overflow 5)		62	H'007C	H'00F8				_
TCI5U (underflow 5)		63	H'007E	H'00FC			Low	69

Note: *Lower 16 bits of the start address

	Origin of		Vector /	Address*			
Interrupt Source	Interrupt Source	Vector Number		Advanced Mode	ICR	IPR	Prio rity

CMIA0 (compare-match A)	8-bit timer channel 0	64	H'0080	H'0100	ICRC7	IPRI6–IPRI4	High
CMIB0 (compare-match B)		65	H'0082	H'0104			
OVI0 (overflow 0)		66	H'0084	H'0108			
CMIA1 (compare-match A)	8-bit timer channel 0	68	H'0088	H'0110	ICRC6	IPRI2-IPRI0	
CMIB1 (compare-match B)		69	H'008A	H'0114			
OVI1 (overflow 0)		70	H'008C	H'0118			
DEND0A (channel 0/channel 0A transfer end)	DMAC	72	H'0090	H'0120	ICRC5	IPRJ6–IPRJ4	
DEND0B (channel 0B transfer end)		73	H'0092	H'0124			
DEND1A (channel 1/channel 1A transfer end)		74	H'0094	H'0128			
DEND1B (channel 1B transfer end		75	H'0096	H'012C		IPRJ2-IPRJ0	
ERI0 (receive error 0)	SCI	80	H'00A0	H'0140	ICRC4		
RXI0 (reception completed 0)	channel 0	81	H'00A2	H'0144			
TXI0 (transmit data empty 0)	,	82	H'00A4	H'0148			
TEI0 (transmission end 0)		83	H'00A6	H'014C		IPRK6-	
ERI1 (receive error 1)						IPRK4	
RXI1 (reception completed 1)	SCI channel 1	84 85		H'0150 H'0154	ICRC3	l	
TXI1 (transmit data empty 1)	•	86	H'00AC	H'0158			
TEI1 (transmission end 1)							
ERI2 (receive error 2)		87	H'00AE	H'015C		IPRK2-	
RXI2 (reception completed 2)	SCI channel 2	88		H'0160	ICRC2	IPPK0	
TXI2 (transmit data empty 2)		89		H'0164			
TEI2 (transmission end 2)		90	H'00B4	H'0168			
		91	H'00B6	H'016C			
							Low

Note: * Lower 16 bits of the start address

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2.10 Operating Modes

The H8S/2655 Series has seven operating modes. These modes enable the selection of initial settings for the CPU operating mode, enabling/disabling of on-chip ROM, and bus width, by setting the mode pins (MD_2 to MD_0).

Normal Modes (Modes 1 to 3)

Mode 1 (Expansion Mode with On-Chip ROM Disabled): The CPU can access a 64-kbyte address space in normal mode. The on-chip ROM is disabled, and 8-bit bus mode is set immediately after a reset.

Ports B and C function as an address bus, port D functions as a data bus, and part of port F carries bus control signals.

Mode 2 (Expansion Mode with On-Chip ROM Enabled): The CPU can access a 64-kbyte address space in normal mode. The on-chip ROM is enabled, and 8-bit bus mode is set. immediately after a reset.

Ports B and C function as input ports immediately after a reset. They can each be set to output addresses by setting the corresponding bits in the data direction register (DDR) to 1. Port D functions as a data bus, and part of port F carries bus control signals.

The amount of on-chip ROM that can be used is limited to 56 kbytes.

Mode 3 (Single-Chip Mode): The CPU can access a 64-kbyte address space in normal mode. The on-chip ROM is enabled, but external addresses cannot be accessed.

All I/O ports are available for use as input-output ports.

The amount of on-chip ROM that can be used is limited to 56 kbytes.

Advanced Modes (Modes 4 to 7)

Mode 4 (Expansion Mode with On-Chip ROM Disabled): The CPU can access a 16-Mbyte address space in advanced mode. The on-chip ROM is disabled.

Ports A, B and C function as an address bus, ports D and E function as a data bus, and part of port F carries bus control signals.

The initial bus mode after a reset is 16 bits, with 16-bit access to all areas. However, if 8-bit access is designated by the bus controller for all areas, the bus mode switches to 8 bits.

Mode 5 (Expansion Mode with On-Chip ROM Disabled): The CPU can access a 16-Mbyte address space in advanced mode. The on-chip ROM is disabled.

Ports A, B and C function as an address bus, ports D and E function as a data bus, and part of port F carries bus control signals.

The initial bus mode after a reset is 8 bits, with 8-bit access to all areas. However, if at least one area is designated for 16-bit access by the bus controller, the bus mode switches to 16 bits and port E becomes a data bus.

Mode 6 (Expansion Mode with On-Chip ROM Enabled): The CPU can access a 16-Mbyte address space in advanced mode. The on-chip ROM is enabled.

Ports A, B and C function as input ports immediately after a reset. They can each be set to output addresses by setting the corresponding bits in the data direction register (DDR) to 1. Port D functions as a data bus, and part of port F carries bus control signals.

The initial bus mode after a reset is 8 bits, with 8-bit access to all areas.

Mode 7 (**single-chip mode**): The CPU can access a 16-Mbyte address space in advanced mode. The on-chip ROM is enabled, but external addresses cannot be accessed.

All I/O ports are available for use as input-output ports.

Kinds of Operating Mode

MCU CPU					Externa Bus	al Data		
Operating Mode	MD2	MD1	MD0	Operating Mode	Description	On-Chip ROM	Initial Width	Max. Width
0	0	0	0	_	_	_	_	_
1	_		1	Normal	Expanded mode with on- chip ROM disabled	Disabled	8 bits	16 bits
2	_	1	0	_	Expanded mode with on- chip ROM enabled	Enabled	8 bits	16 bits
3	_		1	_	Single-chip mode	_	_	_
4	1	0	0	Advanced	Expanded mode with on- chip ROM disabled	Disabled	16 bits	16 bits
5	_		1	_			8 bits	16 bits
6	_	1	0	_	Expanded mode with on- chip ROM enabled	Enabled	8 bits	16 bits
7	_		1	_	Single-chip mode	_		_

2.11 Address Map

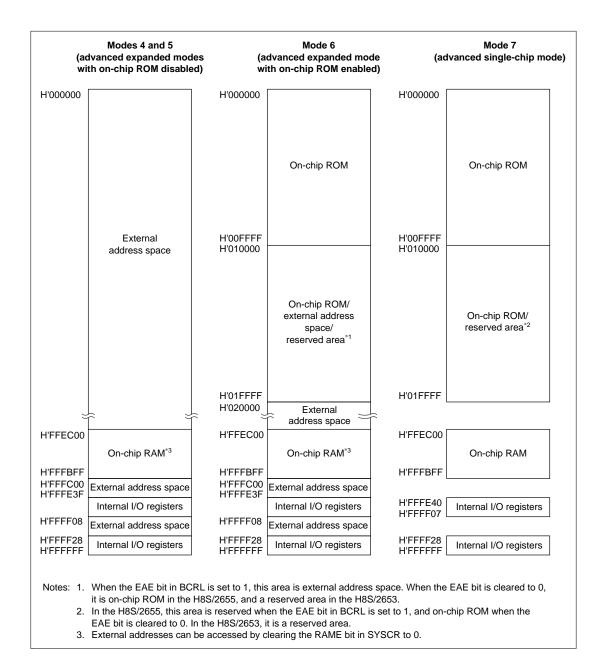
This section shows the address map in each operating mode.

The address space is 64 kbytes in modes 1 to 3 (normal modes) and 16 Mbytes in modes 4 to 7 (advanced modes).

The on-chip ROM size is 64 kbytes, but only 56 kbytes of on-chip ROM can be used in modes 2 and 3 (normal modes).

Address Map in Each Operating Mode

	Mode 1 normal expanded mode h on-chip ROM disabled)		Mode 2 normal expanded mode ith on-chip ROM enabled)	Mode 3 (normal single-chip mode		
H'0000	External address space	H'0000	On-chip ROM	H'0000	On-chip ROM	
		H'DFFF H'E000	External address space	H'DFFF		
H'EC00		H'EC00		H'EC00		
	On-chip RAM*		On-chip RAM*		On-chip RAM	
H'FBFF		H'FBFF		H'FBFF		
H'FC00 H'FE3F	External address space	H'FC00 H'FE3F	External address space			
	Internal I/O registers	3.	Internal I/O registers	H'FE40 H'FF07	Internal I/O registers	
H'FF08	External address space	H'FF08	External address space	111107		
H'FF28 H'FFFF	Internal I/O registers	H'FF28 H'FFFF	Internal I/O registers	H'FF28 H'FFFF	Internal I/O registers	



In modes 4 to 7 the address space is divided into 8 areas. See the Bus Controller section for details.

Section 3 Peripheral Functions

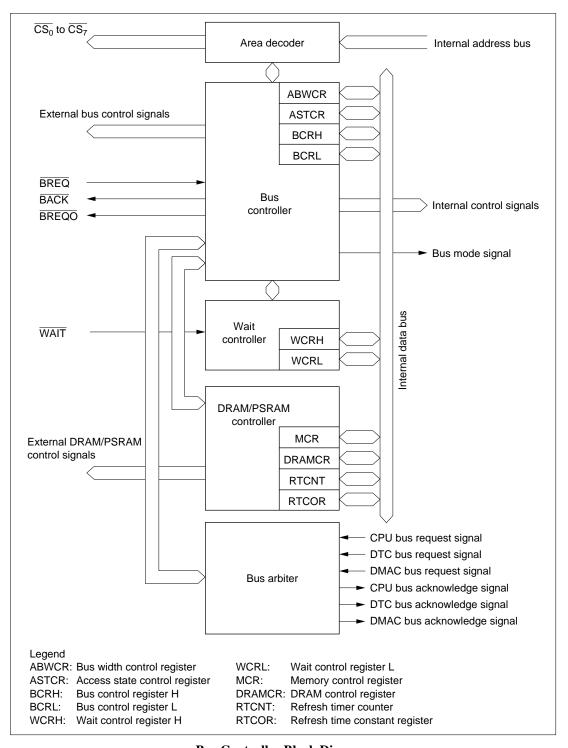
3.1 Bus Controller (BSC)

The bus controller (BSC) manages the external address space divided into eight areas. The bus specifications, such as bus width and number of access states, can be set independently for each area, enabling multiple memories to be connected easily. The bus controller also has a bus arbitration function, and controls the operation of the internal bus masters: the CPU, DMA controller (DMAC), and data transfer controller (DTC).

Features

- Manages external address space in area units
 - In advanced mode, manages the external space as 8 areas of 128-kbytes/2-Mbytes
 - In normal mode, manages the external space as a single area
 - Bus specifications can be set independently for each area
 - DRAM/PSRAM/burst ROM interfaces can be set
- · Basic bus interface
 - Chip select (CS₀ to CS₇) can be output for areas 0 to 7
 - 8-bit access or 16-bit access can be selected for each area
 - 2-state access or 3-state access can be selected for each area
 - Program wait states can be inserted for each area
- DRAM interface
 - DRAM interface can be set for areas 2 to 5 (in advanced mode)
- Pseudo-SRAM direct interface
 - PSRAM interface can be set for areas 2 to 5 (in advanced mode)
- Burst ROM interface
 - Burst ROM interface can be set for area 0
- Idle cycle insertion
- Write buffer functions
 - External write, DMAC single-address mode transfer, and internal access can be executed in parallel
- Bus release
- Bus arbitration function
 - Includes a bus arbiter that arbitrates bus mastership among the CPU, DMAC, and DTC
- Other features

- Refresh counter (refresh timer) can be used as an interval timer
- External bus release function



Bus Controller Block Diagram

3.1.1 Area Partitioning

In advanced mode, the bus controller partitions the 16-Mbyte address space into eight areas, 0 to 7, in 128-kbyte or 2-Mbyte units, and performs bus control for external space in area units. In normal mode, it controls a 64-kbyte access space comprising part of area 0.

Area partitioning is only effective in expanded mode, and has no significance in single-chip mode.

H'000000 H'01FFFF	Area 0 (128 kbytes)	H'000000	Area 0	H'0000
H'020000 H'03FFFF	Area 1 (128 kbytes)	H'1FFFFF	(2 Mbytes)	
H'040000 H'05FFFF	Area 2 (128 kbytes)	H'200000	Area 1 (2 Mbytes)	
H'060000 H'07FFF	Area 3 (128 kbytes)	H'3FFFFF H'400000	(2 Mbytes)	
H'080000 H'09FFFF	Area 4 (128 kbytes)	H'400000	Area 2 (2 Mbytes)	
H'0A0000 H'0BFFFF	Area 5 (128 kbytes)	H'5FFFFF H'600000		L. H'FFEF.
H'0C0000 H'0DFFFF	Area 6 (128 kbytes)		Area 3 (2 Mbytes)	
H'0E0000		H'7FFFF H'800000	Area 4 (2 Mbytes)	
	Area 7	H'9FFFF H'A00000	Area 5 (2 Mbytes)	
	(15 Mbytes)	H'BFFFF H'C00000	Area 6 (2 Mbytes)	
H'FFFFFF		H'DFFFFF H'E00000	Area 7 (2 Mbytes)	
	<u>_</u>		1	
(1) Advanced mode When ASS = 0		(2) Advanced mode When ASS = 1	(2) Normal mode

Overview of Area Partitioning

Bus Specifications

The external address space bus specifications consist of three elements: bus width, number of access states, and number of program wait states. The bus width and number of access states for on-chip memory and internal I/O registers are fixed, and are not affected by the bus controller.

Bus specifications can be set as shown below by means of the bus controller control registers.

Bus Specifications for Each Area (Basic Bus Interface)

ABWCR	ASTCR	WCF	RH, WCRL	Bus Specifications (Basic Bus Interface)				
ABWn	ASTn	Wn1	Wn0	Bus Width	Access States	Program Wait States		
0	0	_	_	16	2	0		
	1	0	0	=	3	0		
			1	=		1		
		1	0	=		2		
			1	_		3		
1	0	_	_	8	2	0		
	1	0	0	_	3	0		
			1	=		1		
		1	0	_		2		
			1	_		3		

Note: In normal mode the bus width is fixed at 8 bits.

Memory Interfaces

The H8S/2655 Series' memory interfaces comprise (1) a basic bus interface that allows direct connection of ROM, SRAM, and so on; (2) a DRAM interface that allows direct connection of DRAM; (3) a PSRAM interface that allows direct connection of PSRAM; and (4) a burst ROM interface that allows direct connection of burst ROM. The interface can be designated independently for each area.

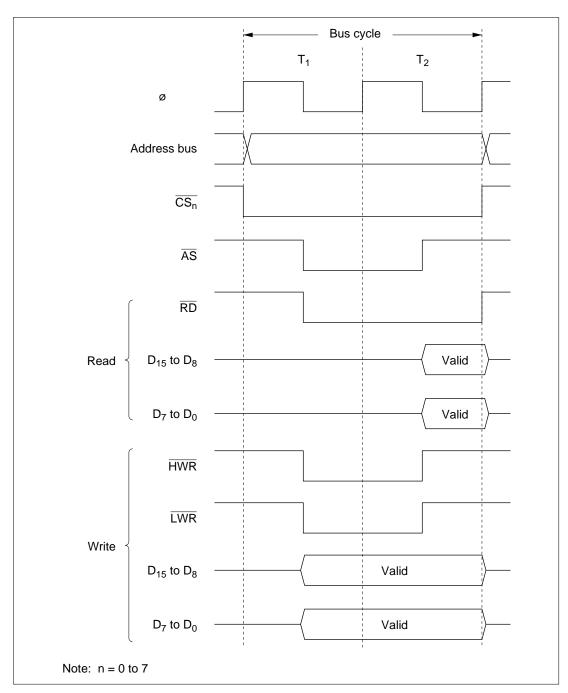
3.1.2 Basic Bus Interface

This interface can be designated for areas 0 to 7. When external address space is accessed, the chip select signal $(\overline{CS}_0$ to $\overline{CS}_7)$ for each area can be output.

In 3-state access space, 0 to 3 program wait states or a pin wait by means of the $\overline{\text{WAIT}}$ pin can be inserted.

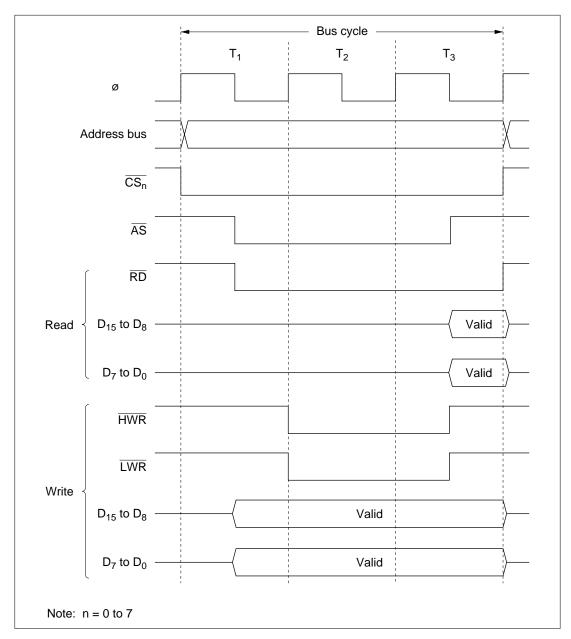
After a reset, all areas are designated as basic bus interface, 3-state access space (the bus width is determined by the MCU operating mode).

Basic Bus Timing



Basic Bus Timing (Word Access to 16-Bit 2-State Access Space)

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Basic Bus Timing (Word Access to 16-Bit 3-State Access Space)

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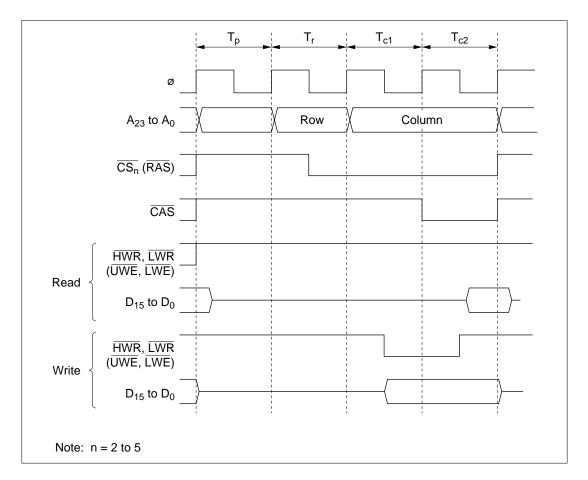
3.1.3 DRAM Interface

In advanced mode, external space areas 2 to 5 can be designated as DRAM space, and DRAM interfacing performed. With the DRAM interface, DRAM can be directly connected to the H8S/2655 Series. Selectable DRAM space settings are: one area (area 2); two areas (areas 2 and 3); and four areas (areas 2 to 5). In an area designated as DRAM space, the $\overline{\text{CS}}$ pin functions as the $\overline{\text{RAS}}$ pin.

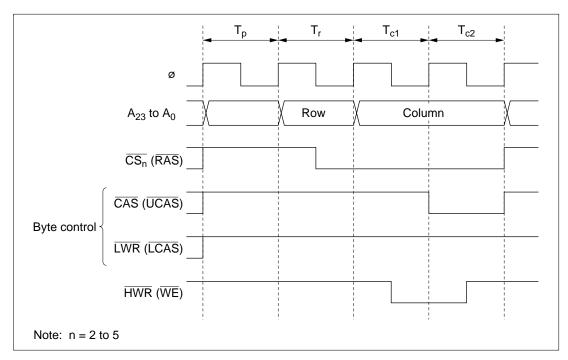
Features

- 2/4/8-Mbyte or 128/256/512-kbyte DRAM space can be set
- Address multiplexing
 - Row address and column address are multiplexed.
 - Selection of 8, 9, or 10 bits as the row address shift size
- Basic timing
 - 4-state basic timing
 - Wait state insertion possible
- DRAM interface
 - Selection of 2-CAS system or 2-WE system for control signals required for byte access, according to the kind of DRAM
- Burst operation
 - Fast page mode
- · Refresh control
 - Selection of CAS-before-RAS refreshing or self-refreshing
 - Can be used as interval timer

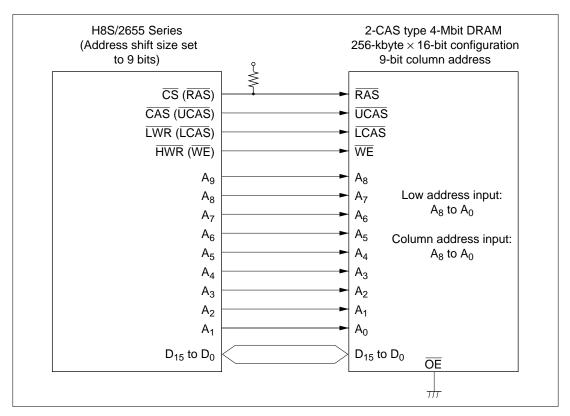
DRAM Basic Timing



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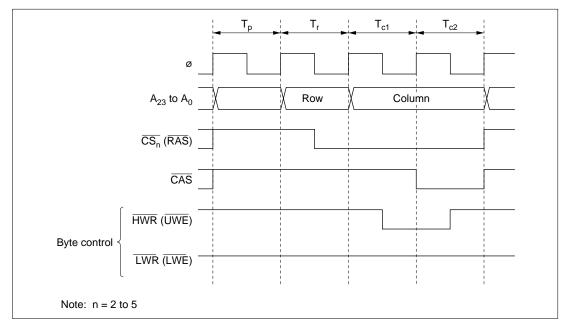


2-CAS System Control Timing (Upper Byte Write Access)

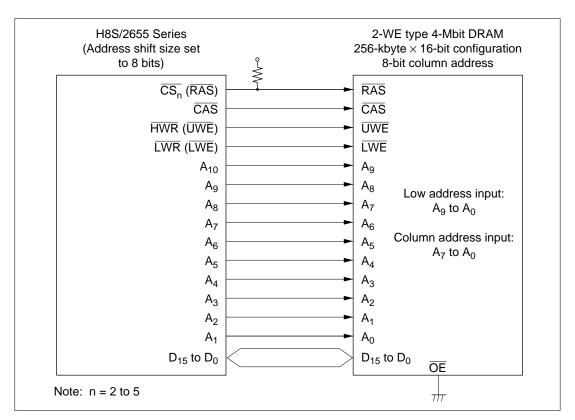


Example of 2-CAS Type DRAM Connection

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2-WE System Control Timing (Upper Byte Write Access)



Example of 2-WE Type DRAM Connection

3.1.4 Pseudo-SRAM Interface

In advanced mode, external space areas 2 to 5 can be designated as pseudo-SRAM (PSRAM) space, and PSRAM interfacing performed. Selectable PSRAM space settings are: one area (area 2); two areas (areas 2 and 3); and four areas (areas 2 to 5).

Features

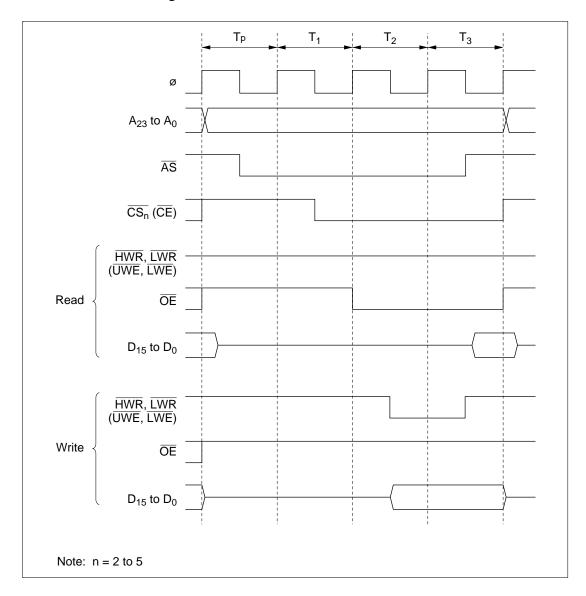
- 2/4/8-Mbyte or 128/256/512-kbyte PSRAM space can be set
- Signal multiplexing
 - Multiplexing of the refresh signal (\overline{RFSH}) and the output enable signal (\overline{OE})
- Basic timing
 - 4-state basic timing
 - Wait state insertion possible
- Burst operation
 - Static column mode

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- Refresh control
 - Selection of auto-refreshing or self-refreshing

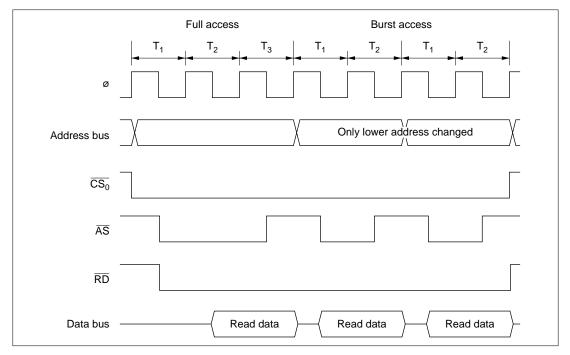
PSRAM Basic Timing



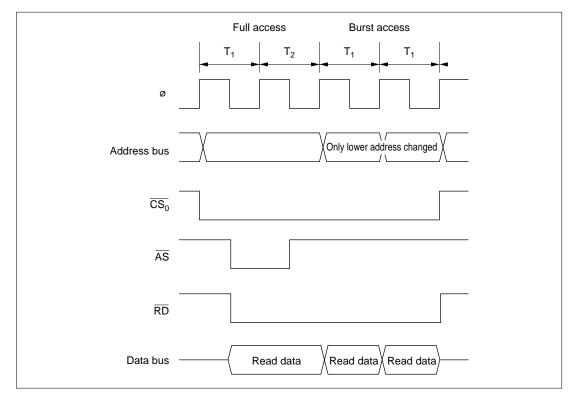
3.1.5 Burst ROM Interface

External space area 0 can be designated as burst ROM space, and burst ROM space interfacing can be performed. The burst ROM space interface enables 16-bit configuration ROM with burst access capability to be accessed at high speed.

Consecutive burst accesses of a maximum or 4 words or 8 words can be performed for CPU instruction fetches only. One or two states can be selected for burst access.



Example of Burst ROM Access Timing (When AST0 = BRSTS1 = 1)



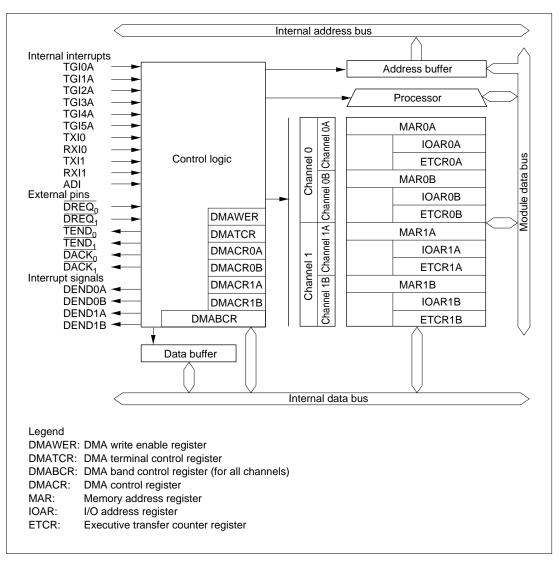
Example of Burst ROM Access Timing (When AST0 = BRSTS1 = 0)

3.2 DMA Controller (DMAC)

The DMA controller (DMAC) can carry out data transfer on up to 4 channels (channels 0A, 0B, 1A, and 1B). Short address transfer can be performed on each channel independently, and full address transfer is possible by using pairs of channels.

Features

- Selection of short address mode or full address mode
- Short address mode
 - Maximum of four channels can be used
 - Selection of dual address mode or single address mode
 - In dual address mode, one of the two addresses, transfer source and transfer destination, is specified as 24 bits and the other as 16 bits
 - In single address mode, transfer source or transfer destination address only is specified as 24 bits
 - In single address mode, transfer can be performed in one bus cycle
 - Selection of sequential mode, idle mode, or repeat mode for dual address mode and single address mode
- Full address mode
 - Maximum of two channels can be used
 - Transfer source and transfer destination address specified as 24 bits
 - Selection of normal mode or block transfer mode
- 16-Mbyte address space can be specified directly
- Byte or word can be set as the transfer unit
- Activation sources: internal interrupt, external request, auto-request (depending on transfer mode)
 - Six 16-bit timer-pulse unit (TPU) compare-match/input capture interrupts
 - Serial communication interface (SCI1, SCI0) transmission complete interrupt, reception complete interrupt
 - A/D converter conversion end interrupt
 - External request
 - Auto-request



DMAC Block Diagram

Transfer Modes

The DMAC has the transfer modes shown in the table below. In short address mode, up to four-channel transfer is possible, with channels A and B operating independently. In full address mode, up to two-channel transfer is possible, with channels A and B combined.

Transfer Mode Table

				Address Register Bit Length		
		Transfer Mode	Transfer Source	Source	Destination	
Short address mode	Dual address mode	 (1) Sequential mode 1-byte or 1-word transfer executed for one transfer requese Memory address incremented/decremented by 1 or 2 1 to 65536 transfers (2)Idle mode 1-byte or 1-word transfer executed for one transfer request Memory address fixed 1 to 65536 transfers (3)Repeat mode 1-byte or 1-word transfer executed for one transfer request Memory address incremented/decremented by 1 or 2 After specified number of transfers (1 to 256), initial state is restored and operation continues 	TPU channel 0 to 5 compare- match/input capture A interrupt SCI transmission complete interrupt SCI reception complete interrupt A/D converter conversion end interrupt External request	24/16	16/24	
	Single address mode	 1-byte or 1-word transfer executed for one transfer request Transfer in 1 bus cycle using DACK pin in place of address specifying I/O Specifiable for modes (1) to (3) 	External request	24/DACK	DACK/24	
Full addr	ess mode	(4)Normal mode Auto-request Transfer request retained internally Transfers continue for the specified number of times (1 to 65536) Selection of burst or cycle steal transfer	Auto-request	24	24	

Transfer Mode	Transfer Source	Source	Destination
 External request 1-byte or 1-word transfer executed for one transfer request 1 to 65536 transfers 	External request		
 (5) Block transfer mode Specified block size transfer executed for one transfer request 1 to 65536 transfers Either source or destination specifiable as block area Block size: 1 to 256 bytes or words 	TPU channel 0 to 5 comparematch/input capture A interrupt SCI transmission complete interrupt SCI reception complete interrupt SCI reception complete interrupt External request A/D converter conversion end interrupt	24	24

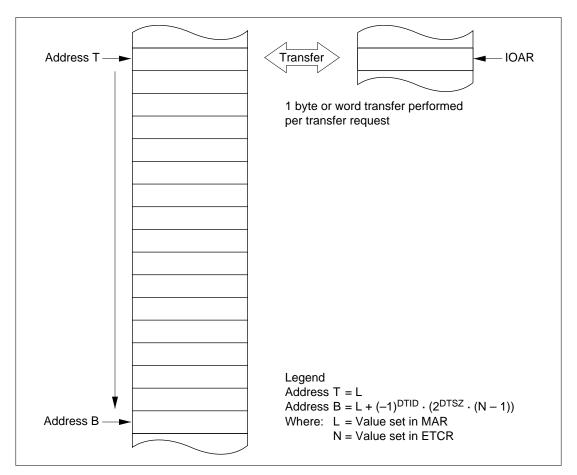
3.2.1 Short Address Mode

There are two kinds of DMAC transfer mode—short address mode and full address mode. Here, short address mode is described.

Short address mode includes (1) sequential mode, (2) idle mode, (3) repeat mode, and (4) single address mode.

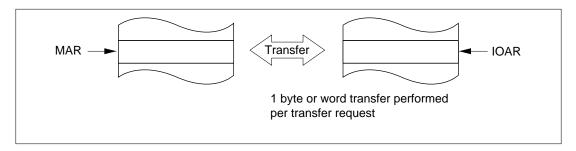
In short address mode, data transfer can be performed on a maximum of four channels.

Operation in Sequential Mode: One byte or word is transferred per transfer request, and a designated number of these transfers are executed. A CPU or DTC interrupt can be requested on completion of the designated number of transfers. One 24-bit address and one 16-bit address are specified. The transfer direction is programmable.



Operation in Sequential Mode

Operation in Idle Mode: One byte or word is transferred per transfer request, and a designated number of these transfers are executed. A CPU or DTC interrupt can be requested on completion of the designated number of transfers. One 24-bit address and one 16-bit address are specified. The transfer source and transfer destination addresses are fixed. The transfer direction is programmable.

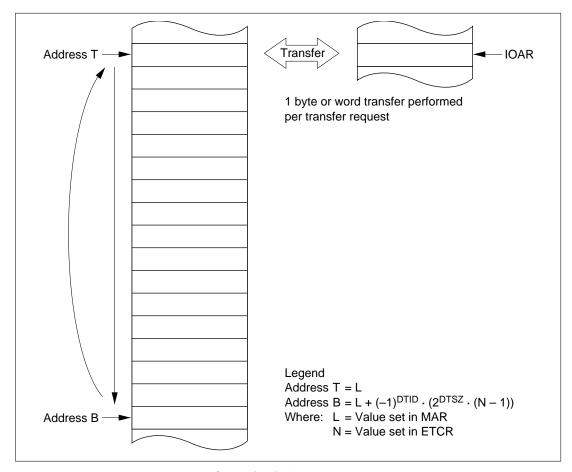


Operation in Idle Mode

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Operation in Repeat Mode: One byte or word is transferred per transfer request, and a designated number of these transfers are executed. On completion of the specified number of transfers, address and transfer counter are automatically restored to their original settings and operation continues. No CPU or DTC interrupt is requested. One 24-bit address and one 16-bit address are specified. The transfer direction is programmable.



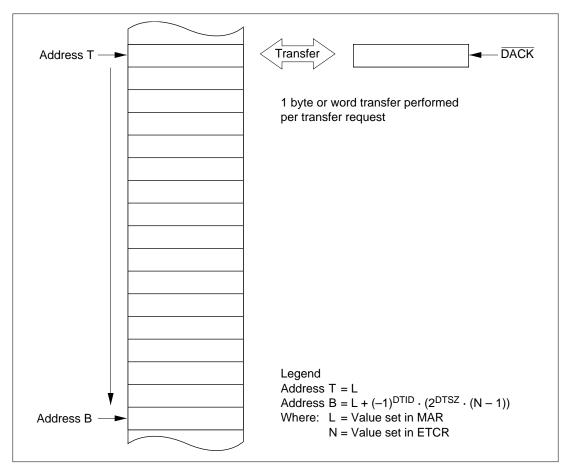
Operation in Repeat Mode

Single Address Mode: One byte or word is transferred per transfer request, and a designated number of these transfers are executed between external memory and an external device. Unlike dual transfer, the source and destination accesses are performed in parallel. Consequently, either the source or the destination is an external device that can be accessed only by a strobe by means of the \overline{DACK} pin. One address is 24 bits, and for the other, the pins are set automatically. The transfer direction is programmable.

Sequential, idle, and repeat modes can also be specified in single address mode.

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Single address mode can only be specified for channel B.



Operation in Single Address Mode (When Sequential Mode is Specified)

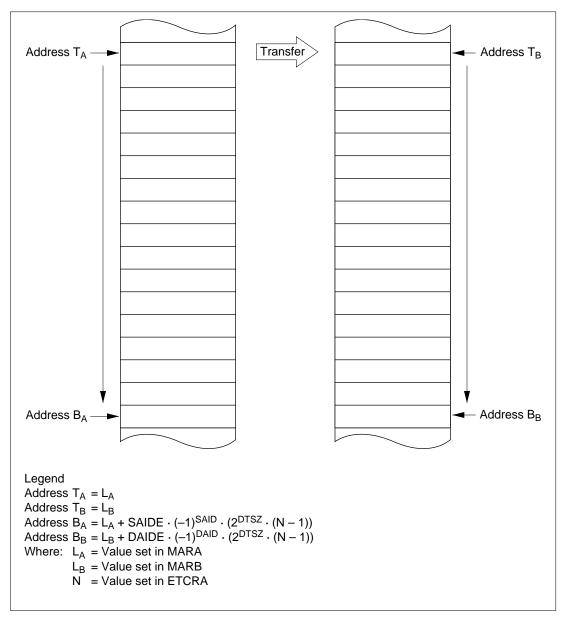
3.2.2 Full Address Mode

There are two kinds of DMAC transfer mode—short address mode and full address mode. Here, full address mode is described.

Full address mode includes (5) normal mode and (6) block transfer mode.

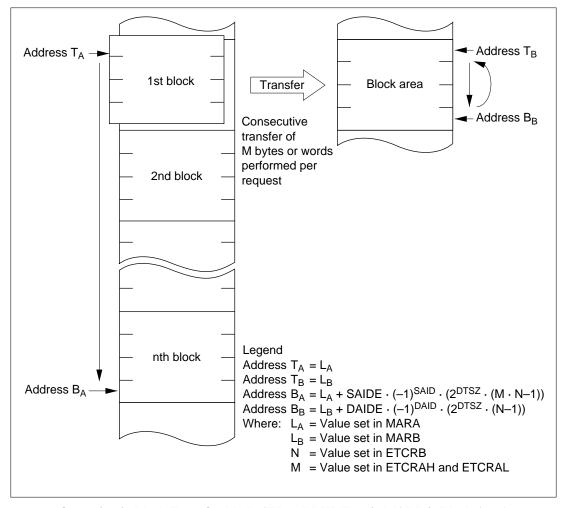
In full address mode, data transfer can be performed on a maximum of two channels, with channels A and B combined.

Normal Mode: One byte or word is transferred per transfer request, and a designated number of these transfers are executed. A CPU or DTC interrupt can be requested on completion of the designated number of transfers. Both addresses are 24-bit addresses. There are two transfer requests (activation sources)—an external request and an auto request.



Operation in Normal Mode

Block Transfer Mode: One block of the specified size is transfer per request, and a designated number of block transfers are executed. At the end of each block transfer, one address is restored to its initial value. When the designated number of blocks have been transferred, a CPU or DTC interrupt can be requested. Both addresses are 24-bit addresses.



Operation in Block Transfer Mode (When BLKDIR = 0: MARB is Block Area)

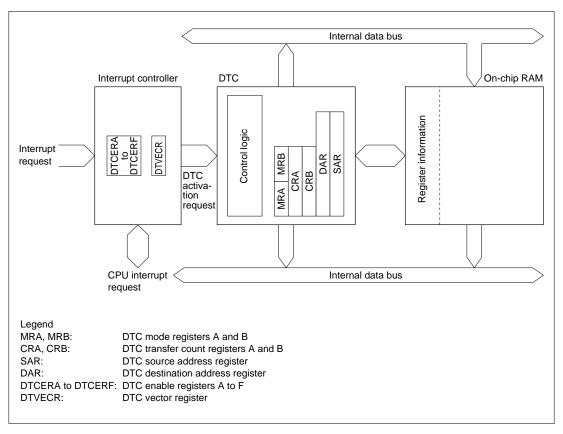
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3.3 Data Transfer Controller (DTC)

The data transfer controller (DTC) is activated by an interrupt or software, and can transfer data without imposing any load on the CPU.

Features

- Transfer possible over any number of channels
 - Transfer information is stored in memory
 - One activation source can trigger a number of data transfers (chain transfer)
- Variety of transfer modes
 - Normal, repeat, and block transfer modes available
 - Incrementing, decrementing, and fixing of source and destination addresses can be selected (destination → destination)
- Direct specification of 16-Mbyte address space possible
- Transfer can be set in byte or word units
- A CPU interrupt can be requested for the interrupt that activated the DTC
 - An interrupt request can be issued to the CPU after one data transfer ends
 - An interrupt request can be issued to the CPU after all specified data transfers have ended
- Can be activated by software

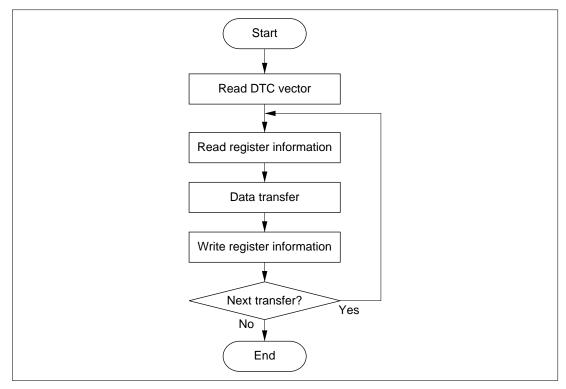


DTC Block Diagram

3.3.1 Data Transfer Operation

When activated, the DTC reads register information previously stored in memory, and transfers data on the basis of that register information. After the data transfer, it writes updated register information back to memory.

Pre-storage of register information in memory makes it possible to transfer data over any required number of channels. The DTC can also execute a number of transfers with a single activation.



Flowchart of DTC Operation

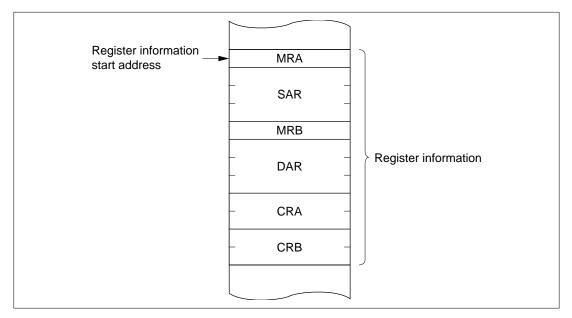
DTC Activation Sources

The DTC operates when activated by an interrupt or by a write to the DTC vector register (DTVECR) by software. An interrupt request can be designated as a CPU interrupt source or a DTC activation source.

When an interrupt has been designated a DTC activation source, existing CPU mask level and interrupt controller priorities have no effect. If there is more than one activation source at the same time, the DTC operates in accordance with the default priorities.

Interrupt Sources and DTC Vector Address

The DTC vector address indicates the start address of the register information in memory. The MRA, SAR, MRB, DAR, CRA, and CRB registers are located in that order from the start address of the register information. Locate the register information in the on-chip RAM (addresses H'FFF800 to H'FFFBFF).



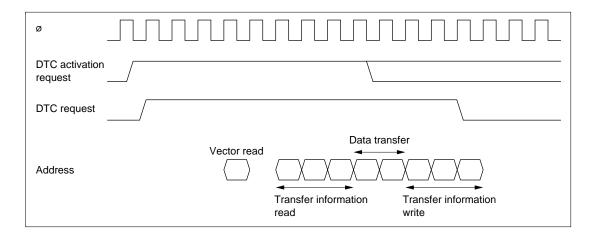
Location of DTC Register Information in Address Space

Interrupt Sources, DTC Vector Addresses, and Corresponding DTCEs

Interrupt Source	Origin of Interrupt Source	Vector Number	Vector Address DTCE	Priorit y
------------------	----------------------------------	------------------	------------------------	--------------

Interrupt Source	Origin of Interrupt Source	Vector Number	Vector Address	DTCE	Priorit y
IRQ0	External pin	16	H'0420	DTCEA7	High
IRQ1		17	H'0422	DTCEA6	lack
IRQ2		18	H'0424	DTCEA5	
IRQ3		19	H'0426	DTCEA4	
IRQ4		20	H'0428	DTCEA3	
IRQ5		21	H'042A	DTCEA2	
IRQ6		22	H'042C	DTCEA1	
IRQ7		23	H'042E	DTCEA0	
ADI (A/D conversion end)	A/D	28	H'0438	DTCEB6	
TGI0A (GR0A compare-match/input capture)	TPU channel	32	H'0440	DTCEB5	
TGI0B (GR0B compare-match/input capture)	0	33	H'0442	DTCEB4	
TGI0C (GR0C compare-match/input capture)		34	H'0444	DTCEB3	
TGI0D (GR0D compare-match/input capture)		35	H'0446	DTCEB2	
TGI1A (GR1A compare-match/input capture)	TPU channel	40	H'0450	DTCEB1	
TGI1B (GR1B compare-match/input capture)	1	41	H'0452	DTCEB0	
TGI2A (GR2A compare-match/input capture)	TPU channel	44	H'0458	DTCEC7	
TGI2B (GR2B compare-match/input capture)	2	45	H'045A	DTCEC6	
TGI3A (GR3A compare-match/input capture)	TPU channel	48	H'0460	DTCEC5	
TGI3B (GR3B compare-match/input capture)	3	49	H'0462	DTCEC4	
TGI3C (GR3C compare-match/input capture)		50	H'0464	DTCEC3	
TGI3D (GR3D compare-match/input capture)	<u> </u>	. 51	H'0466	DTCEC2	
TGI4A (GR4A compare-match/input capture)	TPU channel	56	H'0470	DTCEC1	
TGI4B (GR4B compare-match/input capture)	4	57	H'0472	DTCEC0	
TGI5A (GR5A compare-match/input capture)	TPU channel	60	H'0478	DTCED5	
TGI5B (GR5B compare-match/input capture)		61	H'047A	DTCED4	
CMI0A	8-bit timer channel 0	64	H'0480	DTCED3	
CMI0B		-65	H'0482	DTCED2	
CMI1A	8-bit timer channel 1	68	H'0488	DTCED1	
CMI1B	-DMAC	69	H'048A	DTCED0	
DMTEND0A (DMAC transfer end 0)	DIVIAC	72	H'0490	DTCEE7	
DMTEND0B (DMAC transfer end 1)		73	H'0492	DTCEE6	
DMTEND1A (DMAC transfer end 2)		74	H'0494	DTCEE5	
DMTEND1B (DMAC transfer end 3)	SCI timer	- 75	H'0496	DTCEE4	
RXI0 (reception complete 0)	channel 0	81	H'04A2	DTCEE3	
TXI0 (transmit data empty 0)	_SCI timer	82	H'04A4	DTCEE2	
RXI1 (reception complete 1)	channel 1	85	H'04AA	DTCEE1	
TXI1 (transmit data empty 1)	SCI timer	86	H'04AC	DTCEE0	
RXI2 (reception complete 2)	channel 2	89	H'04B2	DTCEF7	
TXI2 (transmit data empty 2)	Software TACHI	90	H'04B4	DTCEF6	
Write to DTVECR		DTVECR	H'0400+ DTVECR [6:0]<<1		Low

DTC Operation Timing (Normal Mode Example)



Number of DTC Execution States

Mode	Vector Read I	Register Information Read/Write J	Data Read K	Data Write L	Internal Operations M
Normal	1	6	1	1	3
Repeat	1	6	1	1	3
Block transfer	1	6	N	N	3

N: Block size (initial setting of CRAH and CRAL)

Number of States Required in Each Execution State

Access To			On- Chip RAM	On- Chip ROM		hip I/O sters	Exter	nal Dev	/ices	
Bus width			32	16	8	16	8	8	16	16
Access states			1	1	2	2	2	3	2	3
Execution state	Vector read	Sı	_	1	_	_	8	12+4r	n4	6+2m
	Register information read/write	on S _J	1	_	_	_	_	_	_	_
	Byte data read	S_{K}	1	1	2	2	2	3+m	2	3+m
	Word data read	S_K	1	1	4	2	4	6+2m	2	3+m
	Byte data write	S_{L}	1	1	2	2	2	3+m	2	3+m
	Word data write	S_L	1	1	4	2	4	6+2m	2	3+m
	Internal operation	S_{M}	1	1	1	1	1	1	1	1

The number of execution states is calculated from the formula below.

Number of execution states =
$$I \cdot S_I + \Sigma (J \cdot S_J + K \cdot S_K + L \cdot S_L) + M \cdot S_M$$

 Σ indicates the sum of all transfers activated by one activation event (the number designated for chain transfer, plus 1).

3.3.2 Transfer Modes

There are three DTC transfer modes—normal mode, repeat mode, and block transfer mode.

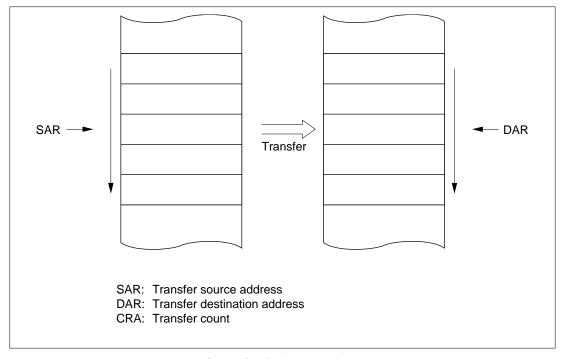
The 24-bit DTC source address register (SAR) designates the DTC transfer source address and the 24-bit destination address register (DAR) designates the transfer destination address. After each transfer, SAR and DAR are independently incremented, decremented, or left fixed.

Address Registers

Transfer Mode	Activation Source	Transfer Source	Transfer Destination
 Normal mode One transfer request transfers one byte or one word Memory addresses are incremented or decremented by 1 or 2 Up to 65,536 transfers possible 	 IRQ TPU TGI 8-bit timer CMI SCI TXI or RXI A/D converter ADI DMAC DEND Software 	24 bits	24 bits
 Repeat mode One transfer request transfers one byte or one word Memory addresses are incremented or decremented by 1 or 2 After the specified number of transfers (1 to 256), the initial state resumes and operation continues 			
 Block transfer mode One transfer request transfers a block of the specified size Block size is from 1 to 256 bytes or words Up to 65,536 transfers possible A block area can be designated at either the source or destination 			

Operation in Normal Mode

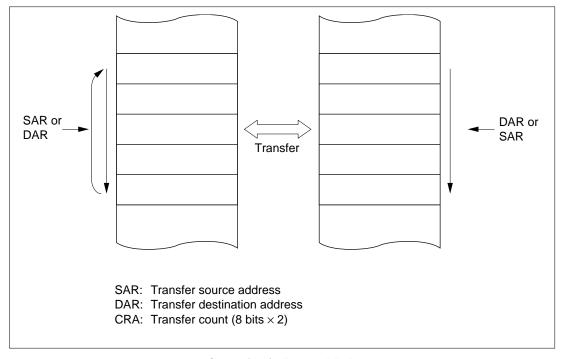
In normal mode, one operation transfers one byte or one word of data. From 1 to 65,536 transfers can be specified. When the specified number of transfers have ended, a CPU interrupt can be requested.



Operation in Normal Mode

Operation in Repeat Mode

In repeat mode, one operation transfers one byte or one word of data. From 1 to 256 transfers can be specified. When the specified number of transfers have ended, the initial settings are restored and transfer is repeated. A CPU interrupt is not requested.

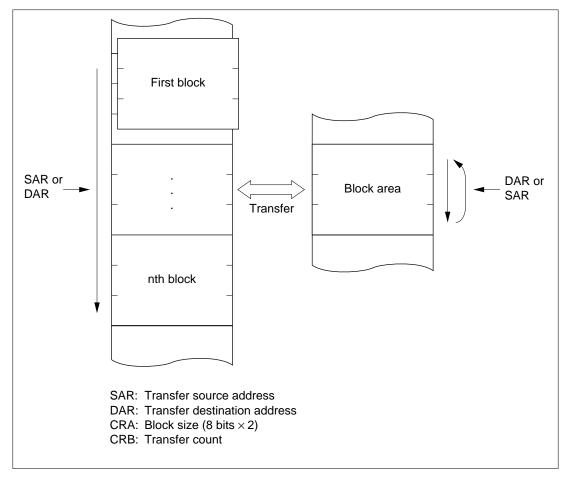


Operation in Repeat Mode

Operation in Block Transfer Mode

In block transfer mode, one operation transfers one block of data. The block size is 1 to 256. When the transfer of one block ends, the initial setting of the address register specified in the block area is restored. The other address register is incremented, decremented, or left fixed.

From 1 to 65,536 transfers can be specified. When the specified number of transfers have ended, a CPU interrupt can be requested.



Operation in Block Transfer Mode

3.4 16-Bit Timer Pulse Unit (TPU)

The 16-bit timer pulse unit (TPU) that comprises six 16-bit timer channels. The TPU can provide up to 16 kinds of pulse input/output.

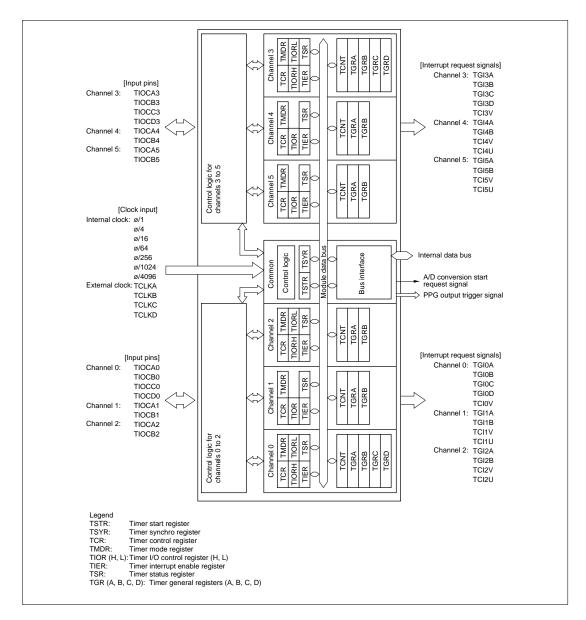
The TPU can perform PWM output, pulse width measurement, and two-phase encoder processing, and can activate the data transfer controller (DTC) and DMA controller (DMAC). It can also generate a programmable pulse generator (PPG) output trigger and A/D converter start trigger.

Features

- Maximum 16 pulse input/outputs
 - A total of 16 timer general registers (TGRs) are provided (four each for channels 0 and 3, and two each for channels 1, 2, 4, and 5), each of which can be set independently as an output compare/input capture register
- Selection of eight counter input clocks for each channel
 - Internal clocks: ø, ø/4, ø/16, ø/64, ø/256, ø/1024, ø/4096
 - External clocks: TCLKA, TCLKB, TCLKC, TCLKD
- The following operations can be set for each channel:
 - Waveform output at compare-match: Selection of 0, 1, or toggle output
 - Input capture function: Selection of rising edge, falling edge, or both edge detection
 - Counter clear operation: Counter clearing possible by compare-match or input capture
 - Synchronous operation:
 - Multiple timer counters (TCNT) can be written to simultaneously
 - Simultaneous clearing by compare-match and input capture possible
 - Simultaneous input/output possible for each register by counter synchronous operation
 - PWM mode:
 - Any PWM output duty can be set
 - Maximum 15-phase PWM output possible by combination with synchronous operation
- Buffer operation settable for channels 0 and 3
 - Input capture register double-buffering possible
 - Automatic rewriting of output compare register possible
- Phase counting mode settable independently for each of channels 1, 2, 4, and 5
 - Two-phase encoder pulse up/down-count possible
- Cascaded operation

- Channel 2 (channel 5) input clock operates as 32-bit counter by setting channel 1 (channel 4) overflow/underflow
- Fast access via internal 16-bit bus
 - Fast access is possible via a 16-bit bus interface
- 26 interrupt sources
 - For channels 0 and 3, four compare-match/input capture dual-function interrupts and one overflow interrupt can be requested independently
 - For channels 1, 2, 4, and 5, two compare-match/input capture dual-function interrupts, one overflow interrupt, and one underflow interrupt can be requested independently
- Automatic transfer of register data
 - Block transfer, one-word transfer, and one-byte transfer possible by data transfer controller (DTC) or DMA controller (DMAC) activation
- Programmable pulse generator (PPG) output trigger can be generated
 - Channel 0 to 3 compare-match/input capture signals can be used as a PPG output trigger
- A/D converter conversion start trigger can be generated
 - Channel 0 to 5 compare-match A/input capture A signals can be used as an A/D converter conversion start trigger

TPU Block Diagram



Interrupt Sources and Data Transfer Controller (DTC) and DMA Controller (DMAC) Activation

TPU Interrupts

	Interrupt		DMAC	DTC	
Channel	Source	Description	Activation	Activation	Priority

0	TGI0A	TGR0A input capture/compare-match	Possible	Possible	High
	TGI0B	TGR0B input capture/compare-match	Not possible	Possible	^
	TGI0C	TGR0C input capture/compare-match	Not possible	Possible	-
	TGI0D	TGR0D input capture/compare-match	Not possible	Possible	-
	TCI0V	TCNT0 overflow	Not possible	Not possible	-
1	TGI1A	TGR1A input capture/compare-match	Possible	Possible	-
	TGI1B	TGR1B input capture/compare-match	Not possible	Possible	-
	TCI1V	TCNT1 overflow	Not possible	Not possible	-
2	TCI1U	TCNT1 underflow	Not possible	Not possible	-
<u>-</u>	TGI2A	TGR2A input capture/compare-match	Possible	Possible	-
	TGI2B	TGR2B input capture/compare-match	Not possible	Possible	-
	TCI2V	TCNT2 overflow	Not possible	Not possible	-
3	TCI2U	TCNT2 underflow	Not possible	Not possible	-
	TGI3A	TGR3A input capture/compare-match	Possible	Possible	-
	TGI3B	TGR3B input capture/compare-match	Not possible	Possible	-
	TGI3C	TGR3C input capture/compare-match	Not possible	Possible	-
	TGI3D	TGR3D input capture/compare-match	Not possible	Possible	-
4	TCI3V	TCNT3 overflow	Not possible	Not possible	-
	TGI4A	TGR4A input capture/compare-match	Possible	Possible	-
	TGI4B	TGR4B input capture/compare-match	Not possible	Possible	-
	TCI4V	TCNT4 overflow	Not possible	Not possible	-
5	TCI4U	TCNT4 underflow	Not possible	Not possible	
	TGI5A	TGR5A input capture/compare-match	Possible	Possible	
	TGI5B	TGR5B input capture/compare-match	Not possible	Possible	-
	TCI5V	TCNT5 overflow	Not possible	Not possible	-
	TCI5U	TCNT5 underflow	Not possible	Not possible	Low

Note: This table shows the initial state immediately after a reset. The relative channel priorities can be changed by the interrupt controller.

Operation

Normal Operation: Each channel has a TCNT and TGR register. TCNT performs up-counting, and is also capable of free-running operation, synchronous counting, and external event counting. Each TGR can be used as an input capture register or output compare register.

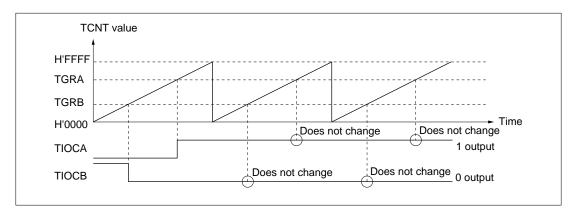
Buffer Operation

- When TGR is an output compare register
- When a compare-match occurs, the value in the buffer register for the relevant channel is transferred to TGR.
- When TGR is an input capture register
- When input capture occurs, the value in TCNT is transfer to TGR and the value previously held in TGR is transferred to the buffer register.

Waveform Output by Compare-Match

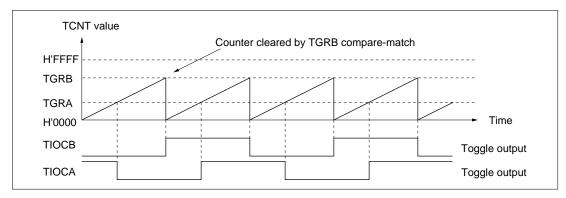
0, 1, or toggle output can be selected.

Example of 0 Output/1 Output Operation: In this example, TCNT has been designated as a free-running counter, and settings have been made so that 0 is output by compare-match A, and 1 is output by compare-match B.



Example of 0 Output/1 Output Operation

Example of Toggle Output: In this example, settings have been made so that TCNT counter clearing is performed by compare-match B, and output is toggled by both by compare-match A and compare-match B.



Example of Toggle Output Operation

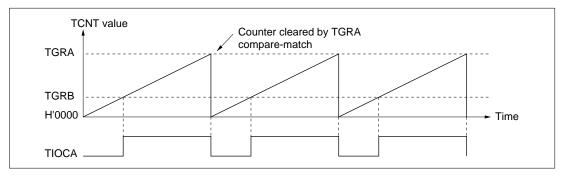
PWM Modes

In PWM mode, PWM waveforms are output from the output pins. There are two PWM modes—PWM mode 1 with a maximum of 8-phase pulse output, and PWM mode 2 with a maximum of 15-phase pulse output.

PWM Mode 1: **PWM output is generated by pairing TGRA with TGRB and TGRC with TGRD.**

In PWM mode 1, a maximum 8-phase PWM output is possible.

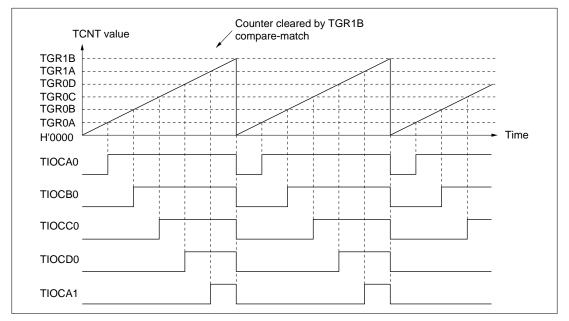
- Example of operation in PWM mode 1
- In this example, TGRA compare-match is set as the TCNT clearing source, 0 is set for the TGRA initial output value and output value, and 1 output is set as the TGRB output value. In this case, the value set in TGRA is the cycle, and the value set in TGRB is the duty.



Operation in PWM Mode 1

PWM Mode 2: PWM output is generated using one TGR register as the cycle register and the others as duty registers. In PWM mode 2, a maximum 15-phase PWM output is possible by combined use with synchronous operation.

- Example of operation in PWM mode 2
- In this example, synchronous operation is designated for channels 0 and 1, TGR1B comparematch is set as the TCNT clearing source, and 0 is set for the initial output value and 1 for the output value of the other TGR registers, to output a 5-phase PWM waveform. In this case, the value set in TGR1B is the cycle, and the value set in the other TGR registers is the duty.



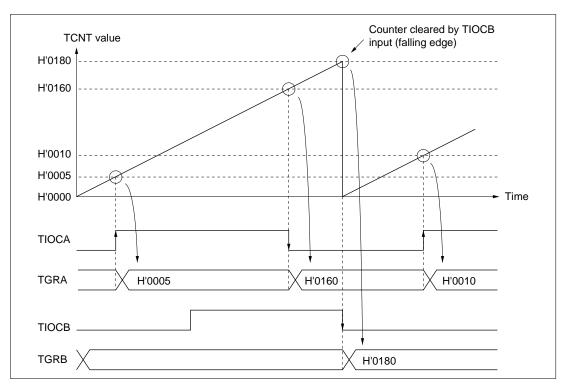
Operation in PWM Mode 2

Input Capture Operation

The TCNT value can be transferred to TGR on detection of the TIOC pin input edge.

Rising edge, falling edge, or both edges can be selected as the input edge.

- Example of input capture operation
- In this example both rising and falling edges have been selected as the TIOCA pin input edge, falling edge has been selected as the TIOCB pin input edge, and counter clearing by TGRB input capture has been designated for TCNT.

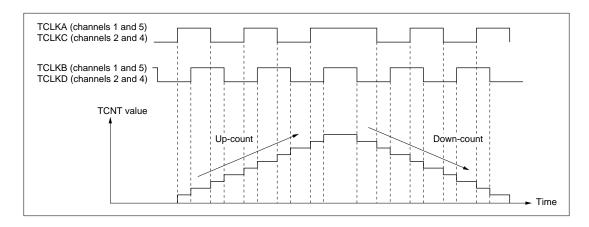


Input Capture Operation

Phase Counting Mode

In phase counting mode, the phase difference between two external clock inputs is detected and TCNT operates as an up/down-counter. There are four modes (phase counting modes 1 to 4) with different setting conditions. These modes can be set for channels 1, 2, 4, and 5.

Example of Operation in Phase Counting Mode (Mode 1 Example)



• Up/Down-Count Conditions in Phase Counting Mode 1 (The up-/down-count conditions are different in phase counting modes 2 to 4.)

TCLKA (Channels 1 and 5) TCLKC (Channels 2 and 4)	TCLKB (Channels 1 and 5) TCLKD (Channels 2 and 4)	Operation
Н		Up-count
L	7_	_
	L	_
1	Н	_
Н	1	Down-count
L		_
<u></u>	Н	_
7_	L	_

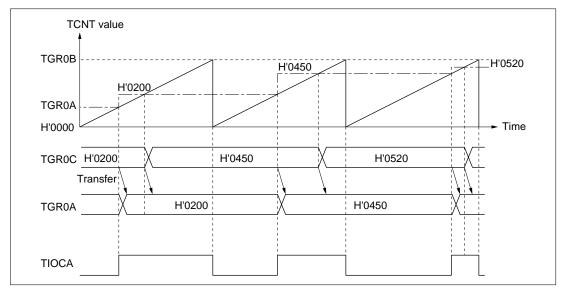
Legend

H: High level
L: Low level
____: Rising edge
____: Falling edge

Buffer Operation

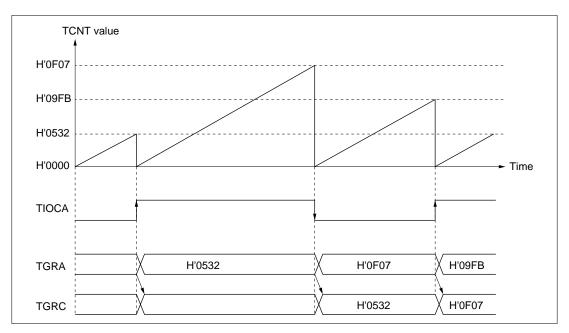
Buffer operation, provided for channels 0 and 3, enables TGRC and TGRD to be used as buffer registers.

- Example of buffer operation (1) (When TGR is an output compare register)
- In this example, PWM mode 1 has been designated for channel 0, and buffer operation has been designated for TGRA and TGRC. The settings used are TCNT clearing by a TGRB compare-match, 1 output at TGRA compare-match, and 0 output at TGRB compare-match. When a compare-match occurs, the output is changed and the value in buffer register TGRC is simultaneously transferred to timer general register TGRA.



Example of Buffer Operation (1) (When TGR Is an Output Compare Register)

- Example of buffer operation (2) (When TGR is an input capture register)
- In this example, TGRA has been designated as an input capture register, and buffer operation has been designated for TGRA and TGRB. Counter clearing by TGRA input capture has been set for TCNT, and both rising and falling edges have been selected as the TIOCA pin input capture input edge. When the TCNT value is stored in TGRA upon occurrence of input capture A, the value previously stored in TGRA is simultaneously transferred to TGRC.

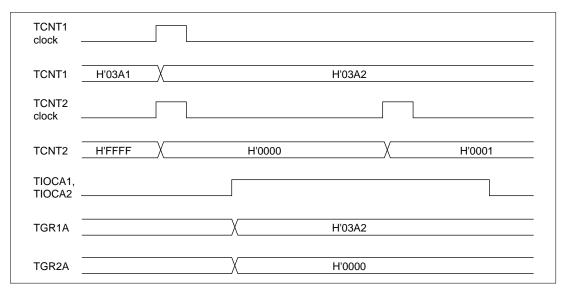


Example of Buffer Operation (2) (When TGR Is an Input Capture Register)

Cascading

In cascaded operation, two 16-bit counters for different channels are used together as a 32-bit counter. Channels 1 and 2, and channels 4 and 5, can be cascaded.

- · Example of cascaded operation
- In this example, counting upon TCNT2 overflow/underflow has been set for TCNT1, TGR1A and TGR2A have been designated as input capture registers, and TIOC pin rising edge detection has been selected. When a rising edge is input to the TIOCA1 and TIOCA2 pins simultaneously, the upper 16 bits of the 32-bit data are transferred to TGR1A, and the lower 16 bits to TGR2A.



Example of Cascaded Operation (32-Bit Input Capture Operation)

Synchronous Operation

When synchronous operation is designated for a channel, TCNT for that channel performs synchronous presetting and clearing. That is, when TCNT for a channel designated for synchronous operation is rewritten, the TCNT counters for the other channels are also rewritten at the same time. When any clearing condition occurs, the TCNT counters for the other channels are also cleared simultaneously.

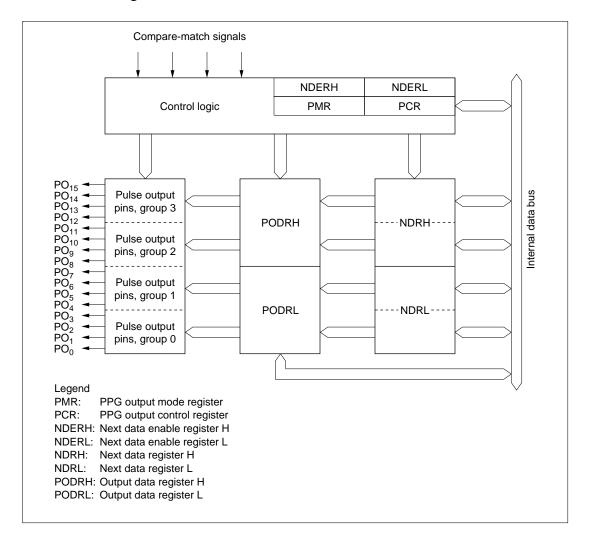
3.5 Programmable Pulse Generator (PPG)

The programmable pulse generator (PPG) can handle up to 16 outputs simultaneously, using a signal from the 16-bit timer-pulse unit (TPU) as input.

Features

- 16-bit output data
 - Maximum 16-bit data can be output, and pulse output can be enabled on a bit-by-bit basis.
- Four output groups
 - Output trigger signals can be selected in 4-bit groups to provide up to four different 4-bit outputs.
- Selectable output trigger signals
 - Output trigger signals can be selected for each group from the compare-match signals of four TPU channels.
- Non-overlap mode
 - A non-overlap margin can be provided between pulse outputs.
- Can operate together with the data transfer controller (DTC) and DMA controller (DMAC)
 - The compare-match signals selected as trigger signals can activate the DTC or DMAC for sequential output of data without CPU intervention.
- Settable inverted output
 - Inverted data can be output for each group.

PPG Block Diagram

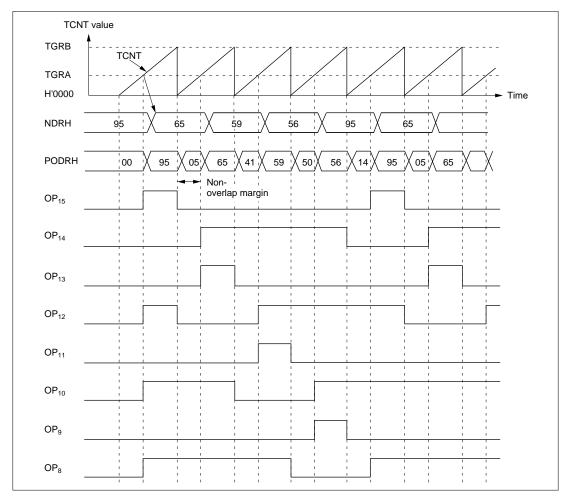


Example of Four-Phase Complementary Non-Overlapping Output

In this example, pulse output is used for four-phase complementary non-overlapping pulse output.

When a TGRB compare-match with occurs, outputs change from 1 to 0. When a TGRA compare-match occurs, outputs change from 0 to 1. Set the non-overlap margin in the TPU TGRA for which the output trigger is selected, and set the cycle in TGRB.

If the DTC or DMAC is set for activation by a TGIA interrupt, pulse output can be performed without imposing a load on the CPU.



Example of Non-Overlapping Pulse Output (Four-Phase Complementary Non-Overlapping)

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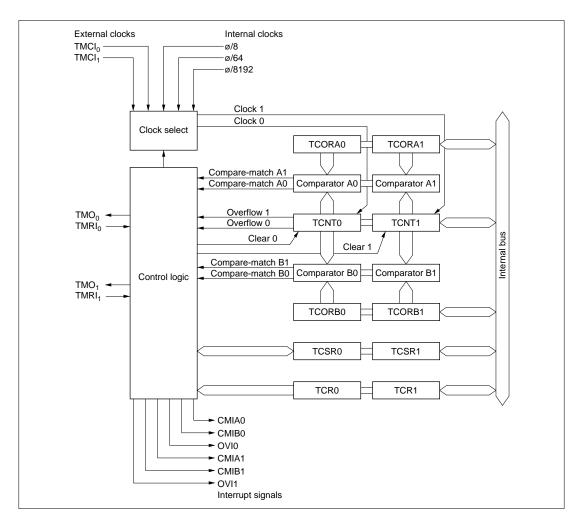
3.6 8-Bit Timer

The H8S/2655 Series includes an 8-bit timer with two channels based on an 8-bit counter. The 8-bit timer can be used for a variety of applications as a multifunctional timer, including pulse output with an arbitrary duty cycle.

Features

- Selection of four input clock sources
 - The clock source can be selected from three internal clock signals ($\emptyset/8$, $\emptyset/64$, or $\emptyset/8192$) or an external clock (external event counting is possible).
- Counter clearing specification
 - The counters can be cleared on compare-match A or B, or by an external reset signal.
- Timer output controlled by combination of two compare-match signals
 - The timer output signal in each channel is controlled by two independent compare-match signals, enabling the timer to generate pulse output or PWM output with an arbitrary duty cycle.
- Three interrupt sources for each channel
 - There are two compare-match sources and one overflow source, capable of independent requests.

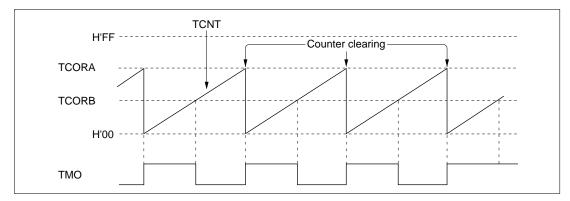
8-Bit Timer Block Diagram



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Example of Pulse Output

TCR is used to set counter clearing by a TCORA compare-match. The cycle is set in TCORA, and the duty in TCORB. The above pulses can be output continuously without software intervention.



Example of Pulse Output

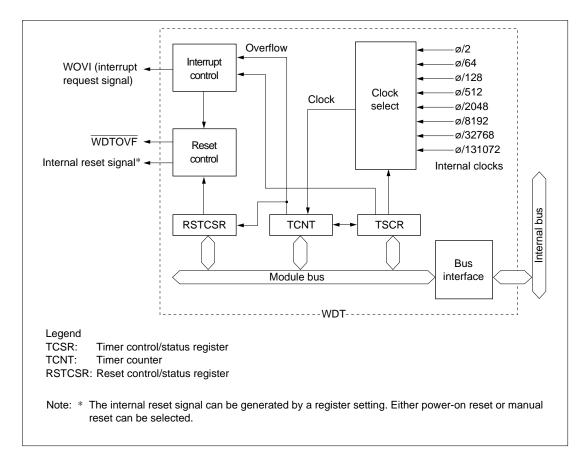
3.7 Watchdog Timer

The H8S/2655 Series can perform system monitoring using its watchdog timer (WDT). When not used as a watchdog timer, this module can be used as an interval timer.

Features

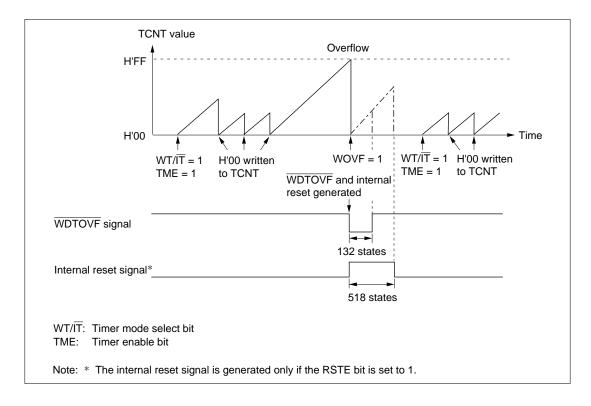
- Selection of eight counter clock sources
 - ø/2, ø/64, ø/128, ø/512, ø/2048, ø/8192, ø/32768, ø/131072
- Can be used as an interval timer
- WDTOVF signal output in watchdog timer mode
 - When the counter overflows, the WDT outputs WDTOVF signal externally. It is possible to select whether or not the entire chip is reset at the same time. Power-on reset or manual reset can be selected as the internal reset.
- Interrupt generation in interval timer mode
 - When the counter overflows, the WDT generates an interval timer interrupt.

Watchdog Timer Block Diagram



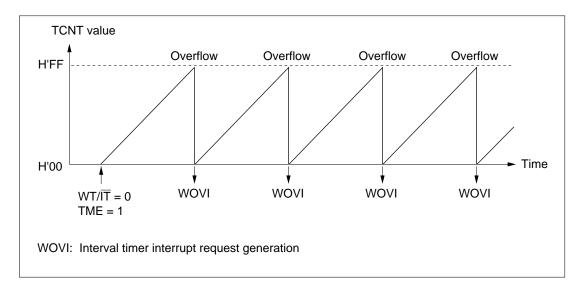
Watchdog Timer Operation

The example below shows this module used as a watchdog timer. The timer counter (TCNT) starts counting up using the specified clock.



Interval Timer Operation

The example below shows this module used as an interval timer. The timer counter (TCNT) starts counting up using the specified clock, and an interval timer request (WOVI) is generated each time TCNT overflows. This function can be used to generate interrupt requests at regular intervals.



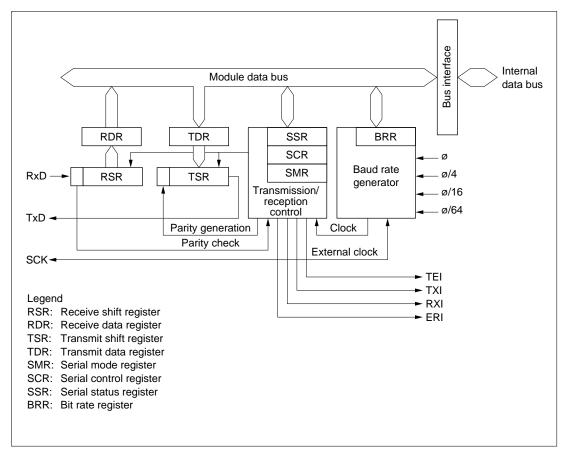
3.8 Serial Communication Interface (SCI)

The H8S/2655 Series is equipped with a three-channel serial communication interface (SCI). All three channels have the same functions, and can handle both asynchronous and synchronous serial communication. A function is also provided for serial communication between processors (multiprocessor communication function).

Features

- Selection of synchronous or asynchronous serial communication mode
- Full-duplex communication capability
- Data register double-buffering enables continuous transmission/reception
- On-chip dedicated baud rate generator allows any bit rate to be selected
- Selection of internal clock from baud rate generator or external clock input (SCK pin) as serial clock source
- Detection of three receive errors
 - Overrun errors, framing errors, and parity errors can be detected
- Break detection
- Four interrupt sources
 - Four interrupt sources—transmit data empty, transmission end, receive data full, and receive error—that can issue requests independently:
 - The transmit data empty interrupt and receive data full interrupt can activate the DMA controller (DMAC) or data transfer controller (DTC) to execute data transfer
- Built-in multiprocessor communication function

SCI Block Diagram



SCI Block Diagram (One Channel)

SCI Interrupt Sources

Channel	Interrupt Source	Description	DTC Activation	DMAC Activation	Priority
0	ERI	Interrupt due to receive error (ORER, FER, or PER)	Not possible	Not possible	High
	RXI	Interrupt due to receive data full (RDRF)	Possible	Possible	-
	TXI	Interrupt due to transmit data empty (TDRE)	Possible	Possible	
	TEI	Interrupt due to transmission end (TEND)	Not possible	Not possible	-
1	ERI	Interrupt due to receive error (ORER, FER, or PER)	Not possible	Not possible	
1	RXI	Interrupt due to receive data full (RDRF)	Possible	Possible	-
	TXI	Interrupt due to transmit data empty (TDRE)	Possible	Possible	
	TEI	Interrupt due to transmission end (TEND)	Not possible	Not possible	
2	ERI	Interrupt due to receive error (ORER, FER, or PER)	Not possible	Not possible	
	RXI	Interrupt due to receive data full (RDRF)	Possible	Not possible	
	TXI	Interrupt due to transmit data empty (TDRE)	Possible	Not possible	
	TEI	Interrupt due to transmission end (TEND)	Not possible	Not possible	Low

Note: * This table shows the initial state immediately after a reset. Relative priorities among channels can be changed by means of ICR and IPR.

3.8.1 SCI Asynchronous Mode

There are two SCI operating modes—asynchronous mode and synchronous mode. Asynchronous mode is described here.

Asynchronous mode is a serial communication mode in which synchronization is achieved character by character basis, using a start bit and one or two stop bits.

Features

• Twelve serial data transfer formats

Data length: 7 or 8 bits
Stop bit length: 1 or 2 bits
Parity: Even/odd/none
Multiprocessor bit: 1 or 0

- Selection of internal baud rate generator or external clock from SCK pin as clock source
- Transmit/receive clock can be output from SCK pin
- Break detection capability
 - Break can be detected by reading the RxD pin level directly in case of a framing error
- Multiprocessor communication capability

Transfer Format and Frame Length in Asynchronous Communication

	SMR S	ettings		Serial Transmit/Receive Format and Frame Length
CHR	PE	MP	STOP	1 2 3 4 5 6 7 8 9 10 11 12
0	0	0	0	S 8-bit data STOP
0	0	0	1	S 8-bit data STOP STOP
0	1	0	0	S 8-bit data P STOP
0	1	0	1	S 8-bit data P STOP STOP
1	0	0	0	S 7-bit data STOP
1	0	0	1	S 7-bit data STOP STOP
1	1	0	0	S 7-bit data P STOP
1	1	0	1	S 7-bit data P STOP STOP
0	_	1	0	S 8-bit data MPB STOP
0	_	1	1	S 8-bit data MPB STOP STOP
1	_	1	0	S 7-bit data MPB STOP
1	_	1	1	S 7-bit data MPB STOP STOP

Legend

S: Start bit
STOP: Stop bit
P: Parity bit
MPB: Multiprocessor bit

Multiprocessor Communication Function

A multiprocessor format, in which a multiprocessor bit is added to the transfer data, can be used for serial communication, enabling data transfer to be performed among a number of processors.

The transmitting station first sends the ID of the receiving station with which it wants to perform serial communication as data with a 1 MPB (multiprocessor bit) added. It then sends transmit data as data with a 0 MPB added.

Receiving stations skip data until data with a 1 MPB is received. Each receiving station then compares that data with its own ID. The station whose ID matches then continues with reception, and accepts data. Stations whose ID does not match continue to skip the data until data with a 1 MPB is sent again.

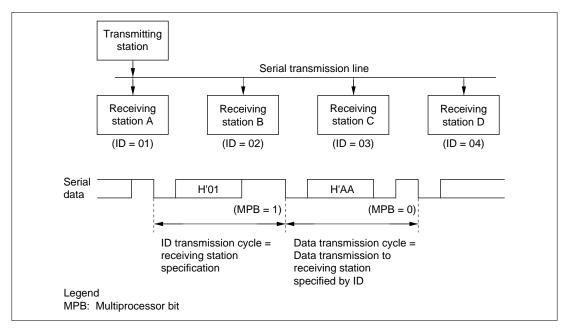
3.8.2 SCI Synchronous Communication

There are two SCI operating modes—asynchronous mode and synchronous mode. Synchronous mode is described here.

In synchronous mode, data is transmitted or received in synchronization with clock pulses, making it suitable for high-speed serial communication.

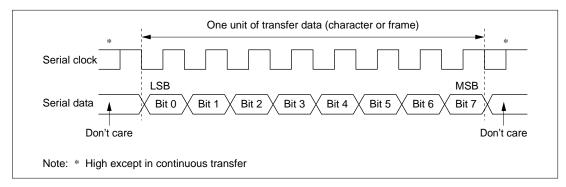
- Data length: 8 bits per character
- Overrun error detection
- Selection of internal baud rate generator or external clock from SCK pin as transmit/receive clock source
- LSB-first system
- Communication is possible with chips provided with a synchronous mode, such as the H8 Series, HD64180, and HD6301

When the internal baud rate generator is selected, the SCK pin is automatically set to output mode, and outputs eight synchronization clock pulses.



Example of Inter-Processor Communication Using Multiprocessor Format (Transmission of Data H'AA to Receiving Station A)

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Data Format in Synchronous Communication

Sample BRR Settings for Various Bit Rates (Synchronous Mode)

2 N 70	n	4 N	_	8		10		16		20
		N						. •		20
70			n	N	n	N	n	N	n	N
	_	_	_	_	_	_	_	_	_	_
12	4 2	249	3	124	_	_	3	249	_	_
24	9 2	124	2	249	_	_	3	124	_	_
12	4 1	249	2	124	_	_	2	249	_	_
19	9 1	99	1	199	1	249	2	99	2	124
99	0	199	1	99	1	124	1	199	1	249
49	0	99	0	199	0	249	1	99	1	124
19	0	39	0	79	0	99	0	159	0	199
9	0	19	0	39	0	49	0	79	0	99
4	0	9	0	19	0	24	0	39	0	49
1	0	3	0	7	0	9	0	15	0	19
0*	0	1	0	3	0	4	0	7	0	9
	0	0*	0	1	_	_	0	3	0	4
		_	_	_	0	0*	_	_	0	1
		_					_	_	0	0*
	12 24 12 19 99 49 19 9	124 2 249 2 124 1 199 1 99 0 49 0 19 0 9 0 4 0 1 0 0* 0	124 2 249 249 2 124 124 1 249 199 1 99 99 0 199 49 0 99 19 0 39 9 0 19 4 0 9 1 0 3 0* 0 1	124 2 249 3 249 2 124 2 124 1 249 2 199 1 99 1 99 0 199 1 49 0 99 0 19 0 39 0 9 0 19 0 4 0 9 0 1 0 3 0 0* 0 1 0	124 2 249 3 124 249 2 124 2 249 124 1 249 2 124 199 1 99 1 199 99 0 199 1 99 49 0 99 0 199 19 0 39 0 79 9 0 19 0 39 4 0 9 0 19 1 0 3 0 7 0* 0 1 0 3	124 2 249 3 124 — 249 2 124 2 249 — 124 1 249 2 124 — 199 1 99 1 199 1 99 0 199 1 99 1 49 0 99 0 199 0 19 0 39 0 79 0 9 0 19 0 39 0 4 0 9 0 19 0 1 0 3 0 7 0 0* 0 1 0 3 0 0 0* 0 1 —	124 2 249 3 124 — — 249 2 124 2 249 — — 124 1 249 2 124 — — 199 1 199 1 249 99 0 199 1 124 49 0 99 0 199 0 249 19 0 39 0 79 0 99 9 0 19 0 39 0 49 4 0 9 0 19 0 24 1 0 3 0 7 0 9 0* 0 1 0 3 0 4 0 0* 0 1 — —	124 2 249 3 124 — — 3 249 2 124 2 249 — — 3 124 1 249 2 124 — — 2 199 1 99 1 199 1 249 2 99 0 199 1 124 1 49 0 99 0 199 0 249 1 19 0 39 0 79 0 99 0 9 0 19 0 39 0 49 0 4 0 9 0 19 0 24 0 1 0 3 0 7 0 9 0 0* 0 1 0 3 0 4 0 0 0* 0 1 — — 0	124 2 249 3 124 — — 3 249 249 2 124 2 249 — — 3 124 124 1 249 2 124 — — 2 249 199 1 99 1 199 1 249 2 99 99 0 199 1 124 1 199 49 0 99 0 199 0 249 1 99 19 0 39 0 79 0 99 0 159 9 0 19 0 39 0 49 0 79 4 0 9 0 19 0 24 0 39 1 0 3 0 7 0 9 0 15 0* 0 1 0 3 0 4 0 7 0 0* 0 1 — —	124 2 249 3 124 — — 3 249 — 249 2 124 2 249 — — 3 124 — 124 1 249 2 124 — — 2 249 — 199 1 99 1 124 1 199 1 49 0 99 0 199 0 249 1 199 1 49 0 99 0 199 0 249 1 99 1 19 0 39 0 79 0 99 0 159 0 9 0 19 0 39 0 49 0 79 0 4 0 9 0 19 0 24 0 39 0 1 0 3 0 7 0 9 0 15 0 0* 0 1 0 3 0

Note: As far as possible, the setting should be made so that the error is no more than 1%.

The BRR setting is found from the following formula:

$$N = \frac{\varnothing}{8 \times 2^{2n-1} \times B} \times 10^6 - 1$$

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Legend

Blank: Cannot be set.

—: Can be set, but there will be a degree of error.

* Continuous transfer is not possible.

N: Baud rate generator setting $(0 \le N \le 255)$

ø: Operating frequency (MHz)

B: Bit rate (bit/s)

n: Baud rate generator input clock (n = 0 to 3)

See the table below for the relation between n and the clock.

n	Clock
0	Ø
1	ø/4
2	ø/16
3	ø/64

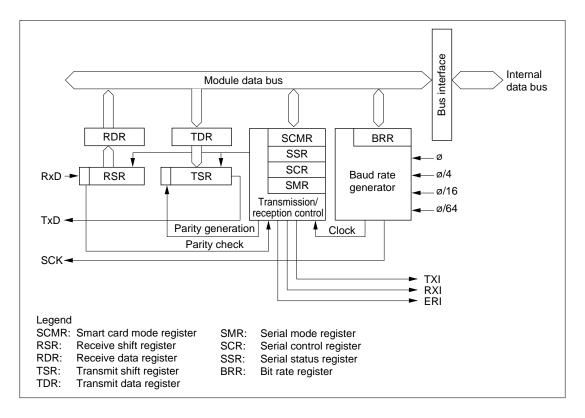
3.9 Smart Card Interface

The SCI supports a smart card interface as an IC card interface serial communication function conforming to ISO/IEC7816-3 (Identification Card).

Features

- Asynchronous mode
 - Data length: 8 bits
 - Parity bit generation and checking
 - Transmission of error signal (parity error) in receive mode
 - Error signal detection and automatic data retransmission in transmit mode
 - Direct convention and inverse convention both supported
- Internal baud rate generator allows any bit rate to be selected
- Three interrupt sources
 - Three interrupt sources—transmit data empty, receive data full, and transmit/receive error—that can issue requests independently
 - The transmit data empty interrupt and receive data full interrupt can activate the DMA controller (DMAC) or data transfer controller (DTC) to execute data transfer

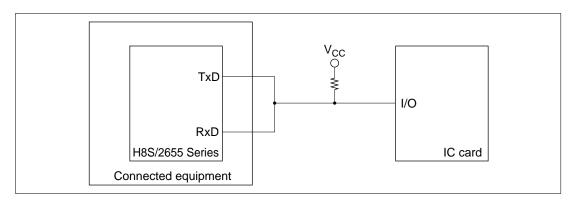
Smart Card Interface Block Diagram



Outline of Operation

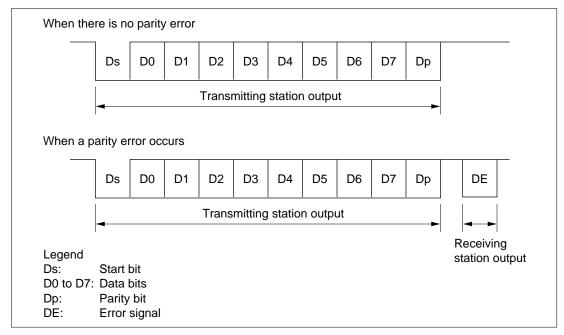
- Only asynchronous communication is supported, with one frame consisting of 8-bit data plus a parity bit.
- In transmission, a guard time of at least 2 etu (Elementary Time Unit: the time for transfer of one bit) is left between the end of the parity bit and the start of the next frame.
- If a parity error is detected during reception, a low error signal level is output for a 1 etu period 10.5 etu after the start bit..
- If the error signal is sampled during transmission, the same data is transmitted automatically after the elapse of 2 etu or longer.

Schematic Connection Diagram



Schematic Diagram of Smart Card Interface Pin Connections

Data Format



Smart Card Interface Data Format

3.10 A/D Converter

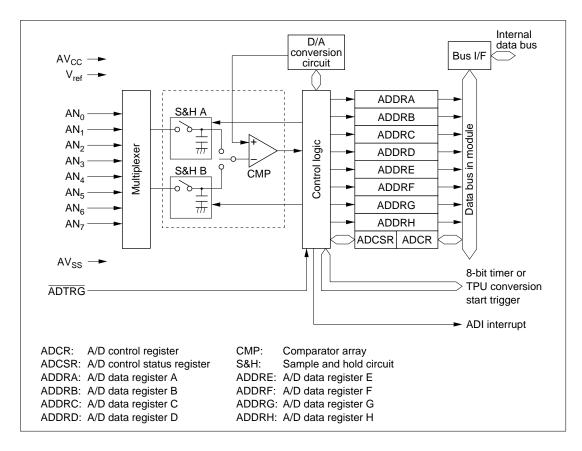
The H8S/2655 Series has an on-chip A/D converter with 10-bit precision. Analog signals can be input on up to eight channels by the program.

Features

A/D converter features are listed below.

- 10-bit resolution
- Eight input channels
- Settable analog conversion voltage range
 - Conversion of analog voltages from 0 V to V_{ref}, with the reference voltage pin (V_{ref}) as the analog reference voltage
- High-speed conversion
 - Minimum conversion time:
 - 2.3 µs per channel (at 20 MHz operation)
 - 1.0 μs per channel in continuous conversion
- Variety of conversion modes
 - Selection of select mode or group mode
 - Selection of single mode or scan mode
 - Buffer operation possible
 - Simultaneous two-channel sampling possible
- Three kinds of conversion start
 - Selection of software or timer conversion start trigger (TPU or 8-bit timer), or ADTRG pin
- Eight data registers
 - Conversion results held in a data register for each channel
- Sample and hold function
- A/D conversion end interrupt generation
 - A/D conversion end interrupt (ADI) request can be generated at the end of A/D conversion

A/D Converter Block Diagram



Input Channel Setting

Eight-channel analog input is performed by means of the group mode bit (GRP) and channel select bits (CH2 to CH0) in ADCSR.

Bit 2	Bit 1	Bit 0	Description	
CH2	CH1	CH0	Select mode (GRP = 0)	Group mode (GRP = 1)
0 0		0	AN ₀ (initial value)	AN ₀
		1	AN ₁	AN ₀ to AN ₁
	1	0	AN ₂	AN ₀ to AN ₂
		1	AN ₃	AN ₀ to AN ₃
1	0	0	AN ₄	AN ₀ to AN ₄
		1	AN ₅	AN ₀ to AN ₅
	1	0	AN ₆	AN ₀ to AN ₆
		1	AN ₇	AN ₀ to AN ₇

Operation

The successive comparison method is used for A/D conversion, with a 10-bit resolution. There are four operating modes—select or group, and single or scan—which can be combined with buffer operation or simultaneous sampling operation.

Select Single Mode: Select single mode is selected when A/D conversion is to be performed on a single channel only.

A/D conversion is started when the ADST bit is set to 1, according to the specified conversion start condition.

On completion of conversion, the ADF flag is set to 1. If the ADIE bit is set to 1 at this time, an ADI interrupt request is generated.

Select Scan Mode: Select scan mode is selected when A/D conversion is to be performed repeatedly on a single channel.

Once the ADST bit is set to 1 according to the specified conversion start condition, A/D conversion is performed repeatedly on the selected channel until the ADST bit is cleared to 0 by software.

An ADI interrupt request can be generated on completion of the first conversion operation.

Group Single Mode: Group single mode is selected when A/D conversion is to be performed on a number of channels.

When the ADST bit is set to 1 according to the specified conversion start condition, A/D conversion starts and is performed on all the selected input channels.

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HITACHI

An ADI interrupt request can be generated on completion of the first conversion operation for all the selected input channels.

Group Scan Mode: Group scan mode is selected when A/D conversion is to be performed repeatedly on a number of channels.

Once the ADST bit is set to 1 according to the specified conversion start condition, A/D conversion is performed repeatedly on the selected channel until the ADST bit is cleared to 0 by software.

An ADI interrupt request can be generated on completion of the first conversion operation for all the selected input channels.

Buffer Operation: In buffer operation, when conversion ends, the conversion result is stored in an ADDR register and the previously stored conversion result is simultaneously transferred to another ADDR register.

Three kinds of buffer operation are possible: a two-stage operation, AIN0 \rightarrow ADDRA \rightarrow ADDRB; dual two-stage operations, AIN0 \rightarrow ADDRA \rightarrow ADDRC and AIN1 \rightarrow ADDRB \rightarrow ADDRD; and a four-stage operation, AIN0 \rightarrow ADDRA \rightarrow ADDRB \rightarrow ADDRC \rightarrow ADDRD.

Simultaneous Sampling: In simultaneous sampling, the analog inputs on two channels are sampled simultaneously, and continuous conversion is performed. The channels involved in simultaneous sampling are shown in the following table.

• Channels Used in Simultaneous Sampling

Channel	Setting	Sampled Channels
CH2	CH1	GRP = 1
0	0	AN_0 , AN_1
	1	AN_0 , $AN_1 \rightarrow AN_2$, AN_3
1	0	$AN_0, AN_1 \to AN_2, AN_3 \to AN_4, AN_5$
	1	$AN_0, AN_1 \to AN_2, AN_3 \to AN_4, AN_5 \to AN_6, AN_7$

3.11 D/A Converter

The H8S/2655 Series has an on-chip D/A converter with 8-bit precision. Analog signals can be output on up to two channels by the program.

Features

D/A converter features are listed below

- Eight-bit resolution
- Two output channels
- Maximum conversion time of 10 µs (with 20 pF load capacitance)
- Output voltage of 0 V to V_{ref}
- D/A output hold function in software standby mode

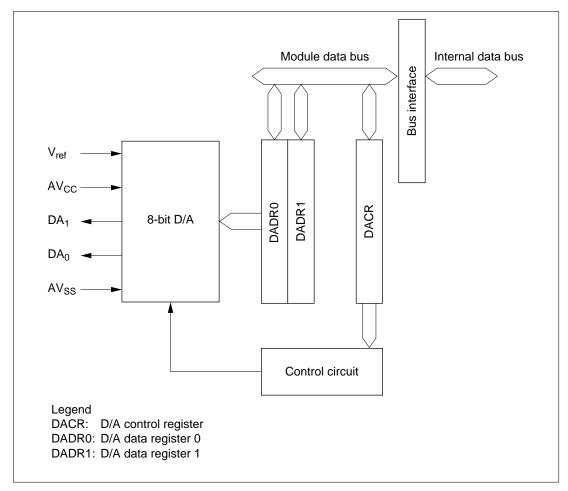
Operation

D/A converter operation is enabled by setting the D/A output enable bit to 1. While this bit is set to 1, DADR contents are constantly converted and output to the corresponding pin.

The output value is:

$$\frac{\text{DADR contents}}{256} \times V_{\text{ref}}$$

D/A Converter Block Diagram



Block Diagram of D/A Converter

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3.12 I/O Ports

The H8S/2655 Series has twelve I/O ports (ports 1, 2, 3, 5, 6, and A to G), and one input-only port (port 4). The ports also function as bus control pins and on-chip supporting module I/O pins.

Each port includes a data direction register (DDR) that controls input/output, a data register (DR) that stores output data, and a port register (PORT) used to read the pin states.

In addition to DDR and DR, ports A to E also have a MOS input pull-up control register (PCR) to control the on/off state of MOS pull-up.

Port Functions in Each Operating Mode

Port Functions

Port	Description	Pins	Mode 1	Mode 2	Mode 3	Mode 4	Mode 5	Mode 6	Mode 7
Port 1	8-bit I/O port	P1 ₇ /PO ₁₅ /TIOCB ₂ / TLCKD P1 ₆ /PO ₁₄ /TIOCA ₂ P1 ₅ /PO ₁₃ /TIOCB ₁ / TLCKC P1 ₄ /PO ₁₂ /TIOCA ₁ P1 ₃ /PO ₁₁ /TIOCD ₀ / TLCKB P1 ₂ /PO ₁₀ /TIOCO ₀ / TLCKA P1 ₁ /PO ₉ /TIOCB ₀ / DACK ₁ P1 ₀ /PO ₈ /TIOCA ₀ / DACK ₀	(DACK ₀ TCLKD,	and DAC TIOCA ₀ ,	K ₁), TPU TIOCB ₀ , '	I/O pins (TIOCC ₀ , T	Toller ou TCLKA, T FIOCD ₀ , T output pir	CLKB, TO	
Port 2	8-bit I/O port Schmitt-triggered input	P2 ₇ /PO ₇ /TIOCB ₅ / TMO ₁ P2 ₆ /PO ₆ /TIOCA ₅ / TMO ₀ P2 ₅ /PO ₅ /TIOCB ₄ / TMCI ₁ P2 ₄ /PO ₄ /TIOCA ₄ / TMRI ₁ P2 ₃ /PO ₃ /TIOCD ₃ / TMCI ₀ P2 ₂ /PO ₂ /TIOCC ₃ / TMRI ₀ P2 ₁ /PO ₁ /TIOCB ₃ P2 ₀ /PO ₀ /TIOCA ₃	TIOCC ₃ , timer (ch	TIOCD ₃ , nannels 0	TIOCA ₄ , and 1) I/0	TIOCB ₄ , 0 pins (TM	pins (TIO TIOCA ₅ , T I/RI ₀ , TMC out pins (P	ΓΙΟČΒ ₅), Cl ₀ , TΜΟ ₀	8-bit

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Port	Description	Pins	Mode 1	Mode 2	Mode 3	Mode 4	Mode 5	Mode 6	Mode 7
Port 3	6-bit I/O port Opendrain output capability	$P3_{5}/SCK_{1}$ $P3_{4}/SCK_{0}$ $P3_{3}/RxD_{1}$ $P3_{2}/RxD_{0}$ $P3_{1}/TxD_{1}$ $P3_{0}/TxD_{0}$		port mult xxD ₀ , SCh	•	,		and 1) I/	O pins
Port 4	8-bit I/O port	P4 ₇ /AN ₇ /DA ₁ P4 ₆ /AN ₆ /DA ₀ P4 ₅ /AN ₅ P4 ₄ /AN ₄ P4 ₃ /AN ₃ P4 ₂ /AN ₂ P4 ₁ /AN ₁ P4 ₀ /AN ₀	8-bit input port multiplexed as A/D converter analog inputs (AN $_7$ to AN $_0$) and D/A converter analog outputs (DA $_1$ and DA $_0$)						
Port 5	4-bit I/O port	P5 ₃ /ADTRG P5 ₂ /SCK ₂ P5 ₁ /RxD ₂ P5 ₀ /TxD ₂		port mult CK_2) and				I/O pins (DTRG)	TxD ₂ ,
Port 6	8-bit I/O port Schmitt- triggered input (P6 ₄ to P6 ₇)	$\begin{array}{c} P6_{7}/\overline{\text{IRQ}}_{3}/\overline{\text{CS}}_{7} \\ P6_{6}/\overline{\text{IRQ}}_{2}/\overline{\text{CS}}_{6} \\ P6_{5}/\overline{\text{IRQ}}_{1} \\ P6_{4}/\overline{\text{IRQ}}_{0} \\ P6_{3}/\overline{\text{TEND}}_{1} \\ P6_{2}/\overline{\text{DREQ}}_{1} \\ P6_{1}/\overline{\text{TEND}}_{0}/\overline{\text{CS}}_{5} \\ P6_{0}/\overline{\text{DREQ}}_{0}/\overline{\text{CS}}_{4} \end{array}$	as DMA pins (DF DREQ ₁ ,	port mult controlle \overrightarrow{REQ}_0 , \overrightarrow{TE} \overrightarrow{TEND}_1) a input pin \overrightarrow{IRQ}_3)	r I/O ND ₀ , and	as DMA pins (\overline{DF} \overline{DREQ}_1 , control to \overline{CS}_7),	port multi controlle REQ ₀ , TE TEND ₁), output pin and inter is (IRQ ₀ t	r I/O ND ₀ , bus s (CS ₄ rrupt	8-bit I/O port multi-plexed as interrupt input pins (IRQ ₀ to IRQ ₃)

Port A	8-bit I/O port Built-in MOS input pull-up Open-drain output capability Schmitt-triggered input (PA ₄ to PA ₇)	$\begin{array}{c} PA_7 \! / A_{23} \! / \overline{IRQ}_7 \\ PA_6 \! / A_{22} \! / \overline{IRQ}_6 \\ PA_5 \! / A_{21} \! / \overline{IRQ}_5 \end{array}$	Multiplexed as I/O port and interrupt input pins (IRQ ₇ to IRQ ₄)	When DDR = 0 (after reset): multiplexed as input port and interrupt input pins (\overline{IRQ}_7 to \overline{IRQ}_5) When DDR = 1: address output	When DDR = 0 (after reset): multi-plexed as input port and interrup t input pins (\overline{IRQ}_7) to \overline{IRQ}_4) When DDR = 1: addres s output	Multi- plexed as I/O port and interrupt input pins (IRQ ₇ to IRQ ₄)
	•	PA ₃ /A ₁₉ to PA ₀ /A ₁₆	I/O port	Address output	When DDR = 0 (after reset): input port When DDR = 1: addres s output	I/O port

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Port	Description	Pins	Mode 1	Mode 2	Mode 3	Mode 4	Mode 5	Mode 6	Mode 7
Port B	8-bit I/O port Built-in MOS input pull-up	PB ₇ /A ₁₅ to PB ₀ /A ₈	Address output	When DDR = 0 (after reset): input port When DDR = 1: address	I/O port	Address	output	When DDR = 0 (after reset): input port When DDR = 1: address	I/O port
Port C	8-bit I/O port Built-in MOS input pull-up	PC ₇ /A ₇ to PC ₀ /A ₀	Address output	output When DDR = 0 (after reset): input port When DDR = 1: address output	I/O port	Address	output	output When DDR = 0 (after reset): input port When DDR = 1: address output	I/O port
Port D	8-bit I/O port Built-in MOS input pull-up	PD ₇ /D ₁₅ to PD ₀ /D ₈	Data bus input/output		I/O port	Data bus input/output		I/O port	
Port E	8-bit I/O port Built-in MOS input pull-up	PE ₇ /D ₇ to PE ₀ /D ₀	In 8-bit but mode: I/O port In 16-bit It mode: data bus input/outp	ous	I/O port	In 8-bit bus mode: I/O port In 16-bit bus mode: data bus input/output			I/O port
Port F	8-bit I/O port	PF ₇ /ø	When DDR = 0: input port When DDR = 1 (after reset): Ø output		When DDR = 0 (after reset): input port When DDR = 1: ø output	When DDR = 0: input port When DDR = 1 (after reset): Ø output		When DDR = 0 (after reset): input port When DDR = 1: ø output	
	•	PF ₆ /AS PF ₅ /RD PF ₄ /HWR PF ₃ /LWR	AS, RD, I LWR outp		I/O port	AS, RD, output	HWR, L	WR	I/O port

•	PF ₂ WAIT/BREQ O	When WAITE = 0 and BREQOE = 0 (after reset): I/O port When WAITE = 1 and BREQOE = 0: WAIT input When WAITE = 0 and BREQOE = 1: BREQO input	When WAITE = 0 and BREQOE = 0 (after reset): I/O port When WAITE = 1 and BREQOE = 0: WAIT input When WAITE = 0 and BREQOE = 1: BREQO output	
•	PF₁/BACK PF₀/BREQ	When BRLE = 0 (after reset): I/O port When BRLE = 1: BREQ input, BACK output	When BRLE = 0 (after reset): I/O port When BRLE = 1: BREQ input, BACK output	

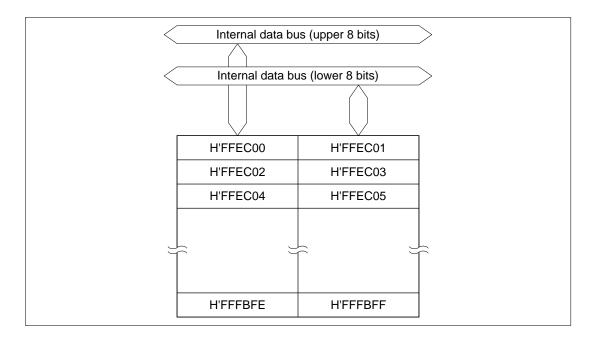
Port	Description	Pins	Mode 1	Mode 2	Mode 3	Mode 4	Mode 5	Mode 6	Mode 7
Port G	• 5-bit I/O port	PG ₄ /CS ₀	When DD input port When DD (after reso output	0R = 1	I/O port	When DD When DD CS ₀ outpu	R = 1 (afte	•	I/O port
	•	PG_3/\overline{CS}_1 PG_2/\overline{CS}_2 PG_1/\overline{CS}_3	I/O port			reset): inp	R = 1: CS		
	•	PG ₀ /CAS/OE				output PSRAM s output	ace set: C	ŌĒ	

3.13 RAM

The H8S/2655 Series has 4 kbytes of on-chip high-speed static RAM. The on-chip RAM is connected to the CPU by a 16-bit data bus, enabling both byte data and word data to be accessed in one state. This makes it possible to perform fast word data transfer.

The on-chip RAM can be enabled or disabled by means of the RAM enable bit (RAME) in the system control register (SYSCR).

Block Diagram of RAM

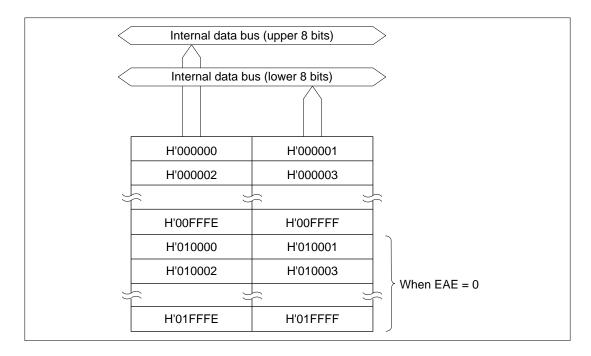


3.14 ROM (**PROM**)

The H8S/2655 Series has 128 kbytes of on-chip ROM (PROM or mask ROM), and the H8S/2653 has 64 kbytes. The ROM is connected to the CPU by a 16-bit data bus, enabling both byte data and word data to be accessed in one state. This makes possible rapid instruction fetches and high-speed processing.

In normal mode, a maximum of 56 kbytes of ROM can be used.

Block Diagram of ROM



PROM Programming

H8S/2655 Series PROM version suspend their microcomputer functions when placed in PROM mode, enabling the on-chip PROM to be programmed. This programming can be done with a PROM programmer set up in the same way as for the HN27C101 EPROM ($V_{PP}=12.5~V$). Use of a 120-pin/32-pin socket adapter enables programming with a commercial PROM programmer. The address range is H'00000 to H'1FFFF. However, page programming is not supported.

Section 4 Power-Down State

4.1 Power-Down State

In addition to the normal program execution state, the H8S/2655 Series has a power-down state in which operation of the CPU and oscillator is halted and power consumption is reduced. The CPU, on-chip peripheral functions, etc., are controlled individually, enabling low-power operation to be achieved. The power-down state includes medium-speed mode, sleep mode, module stop mode, software standby mode, and hardware standby mode.

Medium-Speed Mode: When one or both of the SCK1 and SCK0 bits in the system clock control register (SCKCR) are set to 1, medium-speed mode is entered as soon as the current bus cycle ends. In medium-speed mode, the bus masters—the CPU, DMAC, and DTC—operate on the operating clock (\emptyset /2, \emptyset /4, \emptyset /8, \emptyset /16, or \emptyset /32) specified by the SCK0 and SCK1 bits. However, onchip peripheral functions other than the bus masters operate on the high-speed clock (\emptyset).

In medium-speed mode, a bus access is executed in the specified number of states with respect to the bus master operating clock. For example, if $\emptyset/4$ is selected as the operating clock, on-chip memory is accessed in four states, and internal I/O registers in eight states.

Medium-speed mode is cleared by clearing both the SCK1 and the SCK0 bit to 0. High-speed mode is restored at the end of the current bus cycle.

Sleep Mode: If a SLEEP instruction is executed when the SSBY bit in the system standby register (SBYCR) is cleared to 0, the CPU enters sleep mode. In sleep mode, CPU operation stops but the contents of the CPU's internal registers are retained. Other peripheral functions do not stop.

Sleep mode is cleared by a reset or any interrupt, and the CPU returns to the normal program execution state via the exception handling state.

Module Stop Mode: Module stop mode can be set for individual on-chip peripheral functions.

When the MSTP bit corresponding to a particular peripheral function in the module stop control register (MSTPCR) is set to 1, operation of the specified module stops at the end of the bus cycle and a transition is made to module stop mode. The CPU continues operating independently.

When the corresponding MSTP bit is cleared to 0, module stop mode is cleared and the module starts operating at the end of the bus cycle.

In module stop mode, the internal states of modules other than the SCI are retained.

After a reset, all modules except the DMAC and DTC are in module stop mode.

Software Standby Mode: If a SLEEP instruction is executed when the SSBY bit in SBYCR is set to 1, software standby mode is entered. In this mode, the CPU, on-chip peripheral functions, and oscillator all stop. However, the contents of the CPU's internal registers, RAM data, and the states of on-chip peripheral functions other than the SCI and I/O ports are retained.

Software standby mode is cleared by a reset or an external interrupt. After the elapse of the oscillation stabilization time, the program execution mode is restored via the exception handling state.

As the oscillator is stopped in this mode, power consumption is extremely low.

Hardware Standby Mode: When the STBY pin is driven low, a transition is made to hardware standby mode from any state.

In hardware standby mode, all functions enter the reset state and stop operation, resulting in extremely low power consumption. As long as the prescribed voltage is supplied, on-chip RAM data is retained. I/O ports are set to the high-impedance state.

Hardware standby mode is cleared by means of the STBY pin and the RES pin. When the STBY pin is driven high while the RES pin is low, the reset state is entered and clock oscillation is started. When the RES pin is subsequently driven high, the program execution state is restored via reset exception handling.

In this mode, as in software standby mode, power consumption is extremely low since the oscillator is stopped.

Operating States

Operating	Transition	Clearing		C	PU	Мо	dules	
Mode	Condition	Condition	Oscillator		Registers		Registers	I/O Ports
High- speed mode	Control reg	ister	Functions	High speed	Functions	High speed	Functions	High speed
Medium- speed mode	Control reg	ister	Functions	Medium speed	Functions	High/ medium speed*1	Functions	High speed
Sleep mode	Instruction	Interrupt	Functions	Halted	Retained	High speed	Functions	High speed
Module stop mode	Control reg	ister	Functions	High/ medium speed	Functions	Halted	Retained/ reset *2	Retained
Software standby mode	Instruction	External interrupt	Halted	Halted	Retained	Halted	Retained/ reset*2	Retained
Hardware standby mode	Pin		Halted	Halted	Undefined	Halted	Reset	High impedance

Notes: 1. The bus master operates on the medium-speed clock, and other on-chip peripheral functions operate on the high-speed clock.

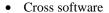
^{2.} The SCI is reset, and other on-chip peripheral functions retain their state.

Section 5 Development Environment

5.1 Development Environment

A comprehensive development environment is provided for the H8S/2655 Series, including cross software, emulators, and HI Series OS.

Lineup



- C compiler
- Assembler
- Simulator/debugger
- Graphical user interface
 - Integration manager
- Emulators
 - E7000
 - E7000PC
- HI Series OS*
 - HI8-2600
- Third-party supplied products
 - Cross software
 - Emulators

Note: * The HI Series OS is Hitachi realtime OS compliant with µITRON specifications.

The product described in this publication is based on the ITRON specifications and was developed under the guidance of Dr. Ken Sakamura of The University of Tokyo.

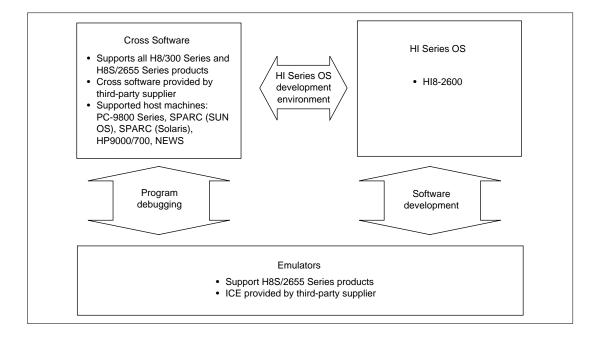
ITRON is an acronym of "Industrial TRON."

TRON is an acronym of "The Real Time Operating System Nucleus."

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µITRON is an acronym of "Micro Industrial TRON."



5.2 Cross Software

Various kinds of cross software—including a C compiler, assembler, and simulator/debugger—are available for the H8S/2655 Series to improve development efficiency.

C Compiler

The C programming language allows system operation and control structures to be written in a concise form. The C compiler converts a program written in C into machine language.

- Comprehensive support of H8S/2655, H8/300, H8/300L, and H8/300H Series
- Complies with ANSI (American National Standard Institute) C standard proposals
- Object program size reduced by an average 30% or more through optimization processing
- · Extended functions for product embedding
 - Intrinsic function facility: H8S/2655 system control instructions can be written in function call format
 - Memory access function: Efficient memory access provided by indirect memory addressing mode and short absolute addressing mode
 - Assembler embedding: Assembly language code can be mixed in with C code
 - Interrupt function creation: Interrupt vector setting and interrupt handling routines can be written in C
- Generates object designed with ROM programming in mind
- Supports debugging information output in optimization

Assembler

Assembly language is suited to hardware-dependent processing that requires fast execution. The assembler converts a program written in assembly language into machine language.

- · Common assembly language for all H Series
- Structured assembler for easy maintenance
 - Supports IF, FOR, WHILE, REPEAT statements, etc.
- Execution instructions and control instructions compliant with IEEE standards
- Efficient macro function

Simulator/Debugger

CPU operation can be simulated by software, and operation of a completed program can be tested without using the actual device.

- Operates on host computer without a target system
- Various powerful debugging functions
 - PC, data, register, sequential, and other break conditions can be set
 - A 1023-instruction trace buffer is supported, and traces can be executed in instruction or subroutine units
 - Stub and function call functions are supported, for easy program debugging in the initial development stage
 - Support of C0 and C1 coverage functions, single-line assembly, disassembly, state retention, command history, command macros, etc.

Lineup

Host Machine	PC-9800 Series	IBM PC FLORA (DOS/V)	SPARC (Sun OS)	SPARC (Solaris)	HP9000/700 HITACHI 9000	NEWS
Assembler system	PS008ASM1-F3	PS008ASM1-IF3	PS008ASM1-SPC	HSS008ASCS1SM	PS008ASM1-H7D	PS008ASM1- NRF
C compiler	PS008CPC100FE 3	PS008CPC100IF3	PS008CSP100	HSS008CLCS1SM	PS008CHP7100	PS008CNR10 0
Simulator/ debugger	PS008SIM1-F3	PS008SIM1-IF3	PS008SIM1-SPC	HSS008SDCS1SM	PS008SIM1-H7D	PS008SIM1- NRF
HI Series OS	HSS008ITPN1SFB (for object contact)	HSS008ITIN1SFB (for object contact (FD))	HSS008ITCN1SFE (for object contact (HSS008ITCN1SME (for object contact ((FD)) 3	HSS008ITHN1SM B (for object contact (MT)) HSS008ITHN1STB (for object contact (DAT))	_
	HSS008ITPN1SFS (for source contact)	HSS008ITIN1SFS (for source contact (FD))		(FD)) S	HSS008ITHN1SM S (for source contact (MT)) HSS008ITHN1STS (for source contact (DAT))	_
Multitask debugger	HSS008I7CN1SF (E7000) HSS008I7IN1SF (E7000PC)					

5.3 Graphical User Interface

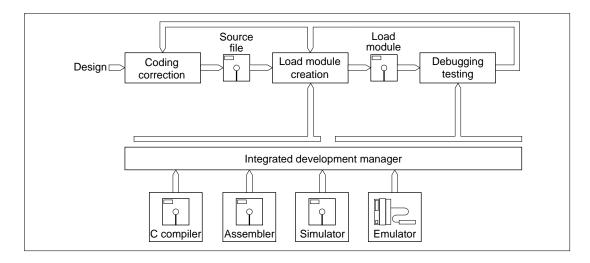
The H8S/2655 Series is provided with a user-friendly, multiwindow-based development environment.

Integrated Development Manager (EWS)*

- · Provides coupling of tools from editor to debugger, for more efficient work
 - Starts editor in case of an assembly/compilation error (and positions cursor at error line)
 - Automatically executes assembly/compilation, object module linkage, and loading into debugger
- · Common, simplified debugger user interface
 - Supports C Source Level debugging in multiwindow environment
 - Simple operation using menus and buttons
 - Unified interface for simulator and emulator
 - Multitask debugger for HI Series OS
- On-line help function
 - Detailed explanations of tool functions displayed in window
 - Error message explanations displayed in window

Note: * Under development.

Conceptual Diagram of Integrated Development Manager



E7000 PC GUI (PC)

- Supports C source level debugging in multiwindow environment
- Simple operation using menus and buttons
- Graphical display of trace information

Lineup

Host Machine	PC-9800 Series	IBM PC FLORA (DOS/V)	SPARC (Sun OS)	SPARC (Solaris)	HP9000/700 HITACHI 9000	NEWS
Integrated development manager or E7000 GUI	-	HS2655G7IW1SF*	_	HS2655IDCM1SM*	HS2655ID7M1ST*	_

Note: * Under development.

5.4 Socket Adapters

A socket adapter is a pin conversion adapter for writing or verifying a program in the on-chip PROM of a microcomputer using a general-purpose PROM writer. A different adapter is available for each product and package type, so be sure to choose the correct adapter for the product concerned.

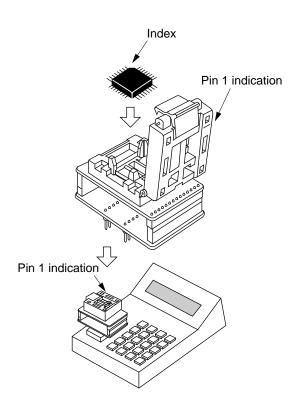
Applicable EPROM Writer Table

				EPROM V	Vriter
Product	Package	Socket Adapter	256	101	101F
H8S/2655, H8S/2653	TFP-120	HS2655ESNS1H*		Yes	
	FP-128	HS2655ESHS1H*		Yes	_

Note: * Under development.

Method of Use (In Case of HS2655ESNS1H)

Refer to the socket adapter instruction manual for details.



5.5 Emulators

Various emulators are available for the H8S/2655 Series, including the E7000 incorporating a wide variety of functions, and the E7000PC for development on MS-WINDOWS.

E7000/E7000PC

These comprise an emulation station common to all models, plus an emulation board and interface cable which can be selected for an individual microcomputer.

An efficient debugging environment is provided by a graphical user interface on a workstation (E7000) or with MS-WINDOWS on an IBM PC or other personal computer (E7000PC).

- · Realtime emulation
 - Capable of realtime emulation supporting a maximum bus cycle speed of up to 50/ns
- Multiwindow environment
 - Runs on standard workstation GUI (graphical user interface) X-Window (E7000)
 - Runs on standard personal computer GUI (graphical user interface) MS-WINDOWS (E7000PC)
 - Provides a multiwindow debugging environment enabling source program amendment while operating the emulator
 - Comprehensive help functions, for manual-less operation
 - Efficient mouse operations through use of pull-down menus and buttons for frequently used commands
- · Fast downloading
 - 1-Mbyte/second downloading
- C source level debugging (graphical user interface)
 - Displays C source program and point at which execution halted
 - Breakpoints can be set or cleared at specified lines in the source program
 - Displays contents of specified variables in the source program
- Powerful debugging functions
 - Break functions: Six independent hardware breaks or 255 software breaks can be set
 - Trace functions: Large trace capacity of 32k cycles; trigger points can be set independently of breakpoints

- Graphical display of trace contents
 - Trace information displayed in graph form, facilitating analysis of program and signal operation

• Other features

- Performance analysis, enabling execution time or number of times executed to be recorded for up to four subroutines
- Parallel mode, enabling memory size to be referenced or changed without halting the program
- Support of memory disassembly, single-line assembly, coverage function (C0), etc.

User System Interface Cable

Product	Package	Cable
H8S/2655, H8S/2653	TFP-120	HS2655ECN71H*
	FP-128	HS2655ECH71H*

Note: * Under development.

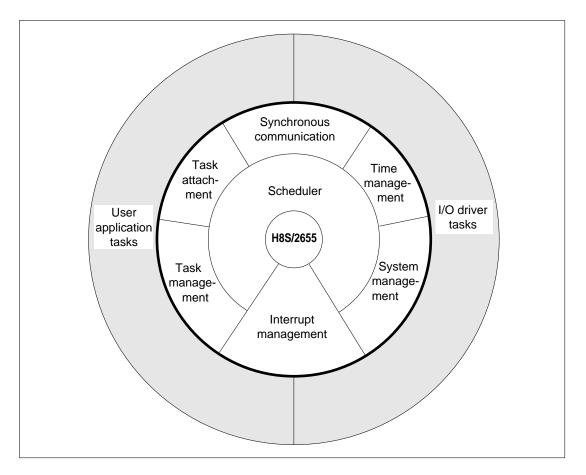
5.6 HI Series OS

The HI Series OS is available for the H8S/2655 Series, enabling the program structure to be made explicit and a custom-made operating system to be constructed.

Features

- Emphasis on realtime characteristics
 - Microsecond order response to interrupts from off-chip
- Compact OS
- Improved development efficiency thought modular software approach
- ROM implementation capability
 - Design presupposes product embedding
- Module structure using building-block system
 - System can be constructed by selecting modules according to the scale of the application system
- Improved development efficiency and reliability
 - System call C interface library provided
 - Shorter development time, lower development costs, improved serviceability and reliability
- Multitask debugger supported on E7000 emulator
 - A multitask debugger that operates on the E7000 emulator is available, for efficient application debugging

HI8-2600 System Configuration Diagram

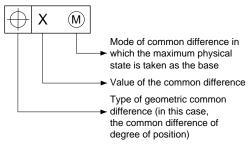


Appendix

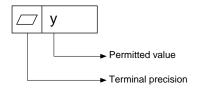
Package

Package Outline Dimensions (Unit: mm)

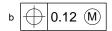
Indication of Geometric Common Difference



Indication of Terminal Precision "y"

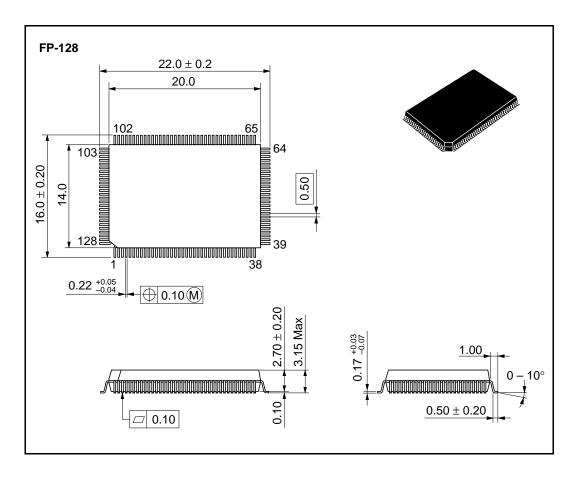


Example



When the terminal width b is the maximum dimension, it indicates that a divergence from the true position of the center position of up to 0.12 mm is permitted.

If b is smaller than the maximum dimension, the common difference corresponding to b can be extended.



TFP-120

